

# Data Structures and Algorithms – CS2402

## Exercises on Algorithms Analysis (2)

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**Exercise 1** Show that  $\log_c(ab) = \log_c(a) + \log_c(b)$ .

*Solution:* Let  $k = \log_c(ab)$ , then by definition of the logarithm,  $ab = c^k$ .  
But now, you can also express  $a$  and  $b$  in terms of a power of  $c$ , separately. Indeed, if  $k_1 = \log_c(a)$  and  $k_2 = \log_c(b)$ , then  $a = c^{k_1}$  and  $b = c^{k_2}$ .  
As a result:  $a.b = c^{k_1}.c^{k_2}$ . Since we also know that  $ab = c^k$ , it comes that  $c^{k_1}.c^{k_2} = c^k$ , which is  $c^k = c^{k_1+k_2}$ .  
We deduce from this that  $k = k_1 + k_2$ , and therefore  $\log_c(ab) = \log_c(a) + \log_c(b)$ .  $\square$

**Exercise 2** Show that  $\log_b(a^n) = n \log_b(a)$ .

*Solution:* Let  $k = \log_b(a)$ . Then, by definition of the logarithm,  $a = b^k$ .  $a^n = (b^k)^n = b^{nk}$ . Therefore,  $\log_b a^n = nk = n \cdot \log_b(a)$ .  $\square$

**Exercise 3** Show that  $\log_b(a) = \frac{\log_c(a)}{\log_c(b)}$ .

*Solution:* Let  $a = b^k$ ,  $a = c^l$  and  $b = c^m$ . Then,  $a = (c^m)^k = c^{l}$ .  
As a result,  $l = m.k$ , and therefore:  $\log_c(a) = \log_b(a) \log_c(b)$ .  $\square$

**Exercise 4** Show that  $\log_b(a) = \frac{1}{\log_a(b)}$ .

*Solution:* Let  $k = \log_b(a)$ , then  $a = b^k$  which means that  $b = a^{\frac{1}{k}}$ . As a result,  $\log_a(b) = \frac{1}{k}$ , and  $\log_b(a) = \frac{1}{\log_a(b)}$ .  $\square$

**Exercise 5** Show that  $a^{\log_b(n)} = n^{\log_b(a)}$ .

*Solution:* Let  $k = \log_b(n)$  and  $l = \log_b(a)$ . Therefore,  $n = b^k$  and  $a = b^l$ .  
 $b^{l.k} = (b^l)^k = (b^k)^l$  which means that  $n^l = a^k$ .  $\square$

**Exercise 6**

<pre>sum=0; for (int i = 0; i &lt; n^2; i* = 5)     sum++;</pre>
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The running time function of this fragment of code is as follows:

<i>sum=0;</i>	<i>1 step</i>
<i>int i=0;</i>	<i>1 step</i>
<i>per valid loop:</i>	
<i>i &lt; n<sup>2</sup></i>	<i>1 step</i>
<i>i* = 5</i>	<i>1 step</i>
<i>sum ++ :</i>	<i>1 step</i>
	<i>= 3 steps per valid loop</i>
	<i>× number of valid loops = log<sub>5</sub>(n<sup>2</sup>)</i>
	<i>= 3 × log<sub>5</sub>(n<sup>2</sup>)</i>
<i>exit condition: i &lt; n<sup>2</sup> when i ≥ n<sup>2</sup></i>	<i>1 step</i>
<hr/> <i>Total steps</i>	<hr/> <i>= 3 + 3 × log<sub>5</sub>(n<sup>2</sup>)</i>

Therefore the big-Oh of this function is:  $\mathcal{O}(\log(n))$ , because  $3 + 3 \times \log_5(n^2) = 3 + 6 \times \log_5(n)$ .

**Exercise 7**

$$n^2 \log(n) \succ n^2 \succ n \log(n) \succ n \succ \log(n)$$