

Logical Foundations of CS – CS5303

Quiz 1 50 points / 50 minutes

Important note: Please make sure that you justify your answers, and that your answers are readable. All non-justified answers will be graded half the points, all unreadable answers will be graded 0. Besides 3 points will be given for the clarity of your answers, and their presentation.

1 Sets (Total: 9 points)

Exercise 1 (5 points) Let A, B, C be three subsets of E . Show that:

$$A \cap \overline{B} = A \cap \overline{C} \Leftrightarrow A \cap B = A \cap C$$

Solution: given in class.

Exercise 2 (4 points) Let A be a subset of E , and let $(B_i)_{i \in I}$ a family of subsets of E . Show that:

$$A \cup \left(\bigcap_{i \in I} B_i \right) = \bigcap_{i \in I} (A \cup B_i)$$

Solution: given in class.

2 Functions (Total: 12 points)

Exercise 3 (4 points) Let $f : A \rightarrow B$ be an application. Show that:

$$f \text{ injective} \Rightarrow \forall X \subseteq A, f^{-1}(f(X)) = X$$

Solution: Suppose that $f : A \rightarrow B$ is an injective function. This means, by definition, that:

$$\forall x_1, x_2 \in A, (f(x_1) = f(x_2)) \Rightarrow x_1 = x_2$$

Let us show that $\forall X \subseteq A, f^{-1}(f(X)) = X$.

To prove this, we are going to consider such an X , subset of A , and prove successively that: $f^{-1}(f(X)) \subseteq X$ and $X \subseteq f^{-1}(f(X))$. And in order to prove each inclusion $Y \subseteq Z$, we will each time consider an element of Y and show that it is in Z .

1. $f^{-1}(f(X)) \subseteq X$? Let $x \in f^{-1}(f(X))$. By definition of $f^{-1}(Y)$, $y \in f^{-1}(Y)$ means that $\exists z \in Y$ such that $z = f(y)$. As a result:

$$x \in f^{-1}(f(X)) \stackrel{\text{def}}{\Leftrightarrow} \exists y \in f(X), f(x) = y$$

At this point we know that $f(x) \in f(X)$. Now we need to finish the proof by showing that if $f(x) \in f(X)$ then necessarily $x \in X$.

In order to prove this, we use the contrapositive rule¹, and we assume that $x \notin X$.

If $x \notin X$, and $f(x) = y \in f(X)$, then it means that $\exists x' \in X$, such that $f(x') = y$. As a result, $\exists x, x'$, such that $f(x) = f(x') = y$ and $x \neq x'$. This contradicts the property of injectivity. As a result, since f is injective, it means that $x = x'$ or $f(x) \notin f(X)$.

If $x = x'$, then $x \in X$, and we are done with the proof. If $f(x) \notin f(X)$, it contradicts our intermediate deduction that $f(x) \in f(X)$, and therefore, by using the contrapositive rule, the condition is contradicted: i.e., $\neg(x \notin X)$ or $\neg(f \text{ is injective})$. Since f is injective, it means that $\neg(x \notin X)$ holds, i.e., $x \in X$, and we are done again with the proof.

2. $X \subseteq f^{-1}(f(X))$?

Let $x \in X$. Let us prove that $x \in f^{-1}(f(X))$. Since $x \in X$, there exists $y \in f(X)$ such that $y = f(x)$. As a result, $x = f^{-1}(y)$ and therefore $x \in f^{-1}(\{y\})$, and more generally, $x \in f^{-1}(f(X))$.

Exercise 4 (8 points) *Let us consider two applications $f : A \rightarrow B$ and $g : B \rightarrow C$. Show that:*

1. $g \circ f$ injective $\Rightarrow f$ injective
2. $g \circ f$ surjective $\Rightarrow g$ surjective

Solution:

1. *Suppose that $g \circ f$ is injective, and that f is not injective.*

Then it means that there are at least x_1 and $x_2 \in A$, such that $x_1 \neq x_2$ and $f(x_1) = f(x_2) = y$. As a result, $g(f(x_1)) = g(f(x_2)) = g(y)$. This contradicts the fact that $g \circ f$ is injective.

2. *Suppose that $g \circ f$ is surjective and that g is not surjective.*

Then it means that there exists at least one element of C , say z^ , such that $\forall y \in B$, $f(y) \neq z^*$. Now since $g \circ f$ is supposed to be surjective too, it means that $\forall z \in C$, $\exists x \in A$, such that $z = g(f(x))$. In particular, for $z^* \in C$, $\exists x^* \in A$, such that $z^* = g(f(x^*))$. Let $y^* = f(x^*)$. By definition of f , $y^* \in B$. Therefore, $\exists y^* \in B$, such that $z^* = g(f(y^*))$, which contradicts g not being surjective. We conclude that under the assumption that $g \circ f$ is surjective, g is necessarily surjective too.*

3 Cardinality (Total: 6 points)

Exercise 5 (6 points) *Show that the set $\mathbb{N} \times \mathbb{N}$ is countable.*

¹The objective here is to prove that $(x \notin X \text{ and } f \text{ injective}) \Rightarrow f(x) \notin f(X)$. And then, since we know (we just proved it), that $f(x) \in f(X)$, then it contradicts our assumption that $(x \notin X \text{ and } f \text{ is injective})$. As soon as our assumption is contradicted, it means that $\neg(x \notin X)$ or $\neg(f \text{ is injective})$. But since we know, by hypothesis in this problem, that f is injective, then we deduce that $\neg(x \notin X)$, which is $x \in X$.

Solution: To show that a set E is countable, you just have to show that there exist a bijective function between E and \mathbb{N} . So here, we are going to try and define a bijective function between $\mathbb{N} \times \mathbb{N}$ and \mathbb{N} .

Let $f : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ be defined as follows:

$$\forall (x, y) \in \mathbb{N} \times \mathbb{N}, (x, y) \mapsto y + (0 + 1 + 2 + 3 + \dots + (x + y)) = y + \sum_{i=0}^{x+y} i$$

Now, we have to prove that this function is a bijective function. There are several ways to prove that f is a bijective function, among which:

- show that it is injective and surjective
- show that $\forall n \in \mathbb{N}, |f^{-1}(\{n\})| = 1$.

We choose the second option, and we prove it by showing that $\forall n \in \mathbb{N}$, there exists a unique element (x, y) of $\mathbb{N} \times \mathbb{N}$ such that $n = f(x, y)$.

$$y = n - (0 + 1 + \dots + a) \quad \text{and} \quad x = a - y$$

where a is the unique element of \mathbb{N} such that:

$$(0 + 1 + \dots + a) \leq n < (0 + 1 + \dots + a + (a + 1))$$

making thus y and x unique.

4 Relations (Total: 20 points)

Exercise 6 (20 points) Let E be a finite set: $E = \{e_1, \dots, e_n\}$, et let \mathcal{R} be a binary relation over E . We can represent \mathcal{R} using a $n \times n$ matrix, say $\mathcal{M}_{\mathcal{R}}$, whose elements belong to $\{0, 1\}$, and are defined as follows:

$$m_{i,j} = \begin{cases} 1 & \text{if } e_i \mathcal{R} e_j \\ 0 & \text{otherwise} \end{cases}$$

1. Describe $\mathcal{M}_{\mathcal{R}}$ in case:

- (a) \mathcal{R} is a symmetrical relation
- (b) \mathcal{R} is a reflexive relation

2. Suppose you have two relations over E , say \mathcal{R} and \mathcal{R}' , how can you compute: $\mathcal{M}_{\mathcal{R}^{-1}}$, $\mathcal{M}_{\overline{\mathcal{R}}}$.

Solution:

1.a. If \mathcal{R} is a symmetric relation, it means that as soon as $(e, e') \in \mathcal{R}$, then $(e', e) \in \mathcal{R}$ too. Let $\mathcal{M}_{\mathcal{R}}(i, j)$ be the coefficient of the $\mathcal{M}_{\mathcal{R}}$ representing whether $(e, e') \in \mathcal{R}$ or not.

Then, if $\mathcal{M}_{\mathcal{R}}(i, j) = 1$, it means that $(e, e') \in \mathcal{R}$, and therefore $(e', e) \in \mathcal{R}$ which translate in the matrix as $\mathcal{M}_{\mathcal{R}}(j, i) = 1$. Reversely, if $\mathcal{M}_{\mathcal{R}}(i, j) = 0$, it means that $(e, e') \notin \mathcal{R}$, and therefore $(e', e) \notin \mathcal{R}$ (otherwise we would have $(e, e') \in \mathcal{R}$ too, by symmetry), which translate in the matrix as $\mathcal{M}_{\mathcal{R}}(j, i) = 0$.

As a result, $\forall i, j, \mathcal{M}_{\mathcal{R}}(i, j) = \mathcal{M}_{\mathcal{R}}(j, i)$, which means that $\mathcal{M}_{\mathcal{R}}$ is symmetric.

1.b. If \mathcal{R} is a reflexive relation, it means that $\forall e \in E, (e, e) \in \mathcal{R}$. This translates in the matrix as $\forall i, \mathcal{M}_{\mathcal{R}}(i, i) = 1$. Therefore $\mathcal{M}_{\mathcal{R}}$ is a matrix whose diagonal contains only 1s.

2. Given $\mathcal{M}_{\mathcal{R}}$, we need to define a transformation to get $\mathcal{M}_{\mathcal{R}^{-1}}$ and $\mathcal{M}_{\overline{\mathcal{R}}}$. Let us go back to the definitions of \mathcal{R}^{-1} and $\overline{\mathcal{R}}$.

\mathcal{R}^{-1} : $(e, e') \in \mathcal{R}$ iff $(e', e) \in \mathcal{R}^{-1}$. As a result, $\mathcal{M}_{\mathcal{R}^{-1}} = \mathcal{M}_{\mathcal{R}}^T$.

$\overline{\mathcal{R}}$: $(e, e') \in \mathcal{R}$ iff $(e, e') \notin \overline{\mathcal{R}}$. As a result, $\mathcal{M}_{\overline{\mathcal{R}}} = \text{One} - \mathcal{M}_{\mathcal{R}}$, where One is a matrix of the same size as $\mathcal{M}_{\mathcal{R}}$, but containing only 1s.