

Logical Foundations of CS – CS5303

Exercises on Predicate Logic 2

Exercise 1 *Translate the following sentences into predicate logic formulas.*

1. *Not every book is interesting.*
2. *Some books are interesting.*
3. *Some books are interesting and some are easy to read.*
4. *Every interesting book is easy to read.*
5. *Whenever a book is well written or easy to understand, it is expensive.*
6. *Not all expensive books are easy to understand.*
7. *A well written book is not necessarily easy to understand.*
8. *For a book to be easy to understand it is necessary that it is well written.*
9. *For every expensive book, there is a book that is not expensive but well written.*
10. *Not every red bee is large.*
11. *Every red bee likes every large flower.*
12. *Not everyone likes everyone else.*
13. *Not every nice person is happy.*
14. *Every nice person likes every happy person.*
15. *There is someone whom no one likes.*
16. *For any even number there is an odd number which is greater than that.*

Then, put the formulas into clausal form (i.e., prenex and then skolem, and CNF).

Exercise 2 *Formalize the following reasonings in predicate logic. Then prove that these reasonings are correct.*

1. *John, a student in this class, is 16 years old. Everyone who is 16 years old can get a driver's liscence. Therefore, someone in this class can get a driver's liscence.*
2. *Somebody in this class enjoys hiking. Every person who enjoys hiking also likes biking. Therefore, there is a person in this class who likes biking.*

3. *Every student in this class owns a personal computer. Everyone who owns a personal computer can use the Internet. Therefore, John, a student in this class, can use the Internet.*
4. *Everyone in this class owns a personal computer. Someone in this class has never used the Internet. Therefore, someone who owns a personal computer has never used the Internet.*

Exercise 3 *Determine whether each of the following arguments is valid. If an argument is correct, what rule of inference is being used? If it is not, what fallacy occurs?*

1. *If n is a real number with $n > 1$, then $n^2 > 1$. Suppose that $n^2 \leq 1$. Then $n \leq 1$.*
2. *If n is a real number with $n > 1$, then $n^2 > 1$. Suppose that $n^2 > 1$. Then $n > 1$.*

Solutions and other:

Check: <http://andromeda.cs.odu.edu:28080/InferenceCheck/index.html>

Translations:

1. Not every book is interesting: $\neg\forall x(\text{book}(x) \rightarrow \text{interesting}(x))$
2. Some books are interesting: $\exists x, (\text{book}(x) \wedge \text{interesting}(x))$
3. Some books are interesting and some are easy to read:
 $(\exists x, (\text{book}(x) \wedge \text{interesting}(x))) \wedge (\exists x, (\text{book}(x) \wedge \text{easy} - \text{to} - \text{read}(x)))$
4. Every interesting book is easy to read:
 $\forall x, ((\text{book}(x) \wedge \text{interesting}(x)) \rightarrow \text{easy} - \text{to} - \text{read}(x))$
5. Whenever a book is well written or easy to understand, it is expensive:
 $\forall x, ((\text{book}(x) \wedge (\text{well} - \text{written}(x) \vee \text{easy} - \text{to} - \text{understand}(x))) \rightarrow \text{expensive}(x))$
6. Not all expensive books are easy to understand:
 $\neg\forall x ((\text{book}(x) \wedge \text{expensive}(x)) \rightarrow \text{easy} - \text{to} - \text{understand}(x))$
7. A well written book is not necessarily easy to understand:
 $\neg\forall x, ((\text{book}(x) \wedge \text{well} - \text{written}(x)) \rightarrow \text{easy} - \text{to} - \text{understand}(x))$
8. For a book to be easy to understand it is necessary that it is well written:
 $\forall x, (\text{book}(x) \rightarrow (\text{easy} - \text{to} - \text{understand}(x) \rightarrow \text{well} - \text{written}(x)))$
 $\forall x, ((\text{book}(x) \wedge \text{easy} - \text{to} - \text{understand}(x)) \rightarrow \text{well} - \text{written}(x))$
9. For every expensive book, there is a book that is not expensive but well written:
 $\forall x, ((\text{book}(x) \wedge \text{expensive}(x)) \rightarrow (\exists y, (\text{book}(y) \wedge \neg\text{expensive}(y) \wedge \text{well} - \text{written}(y))))$
10. Not every red bee is large: $\neg\forall x((\text{red}(x) \wedge \text{bee}(x)) \rightarrow \text{large}(x))$
11. Every red bee likes every large flower:
 $\forall x, ((\text{red}(x) \wedge \text{bee}(x)) \rightarrow (\forall y, ((\text{flower}(y) \wedge \text{large}(y)) \rightarrow \text{likes}(x, y))))$
12. Not everyone likes everyone else:
 $\neg\forall x, (\text{person}(x) \rightarrow (\forall y, (\text{person}(y) \rightarrow \text{likes}(x, y))))$
13. Not every nice person is happy:
 $\neg\forall x, ((\text{person}(x) \wedge \text{nice}(x)) \rightarrow \text{happy}(x))$
14. Every nice person likes every happy person:
 $\forall x, ((\text{person}(x) \wedge \text{nice}(x)) \rightarrow (\forall y, ((\text{person}(y) \wedge \text{happy}(y)) \rightarrow \text{likes}(x, y))))$
15. There is someone whom no one likes:
 $\exists x, (\text{person}(x) \wedge (\forall y, (\text{person}(y) \rightarrow \neg\text{likes}(y, x))))$
16. For any even number there is an odd number which is greater than that. $\forall x, ((\text{number}(x) \wedge \text{even}(x)) \rightarrow (\exists y, (\text{number}(y) \wedge \text{odd}(y) \wedge \text{greater}(y, x))))$

Reasonings:

1. Let:

- $C(x)$ be: x is a student in this class,
- $S(x)$ be: x is 16 years old,
- and $D(x)$ be: x can get a driver's licence.

The premises are $C(\text{John})$, $S(\text{John})$ and $\forall x, (S(x) \rightarrow D(x))$. Using universal instantiation and the last premise, $S(\text{John}) \rightarrow D(\text{John})$ follows. Applying modus ponens to this conclusion and the second premise, $D(\text{John})$ follows. Using conjunction and the first premise, $C(\text{John}) \wedge D(\text{John})$ follows. Finally, using existential generalization, the desired conclusion, $\exists x, (C(x) \wedge D(x))$ follows.

2. Let:

- $C(x)$ be: x is in this class,
- $H(x)$ be: x enjoys hiking,
- and $B(x)$ be: x likes biking.

The premises are $\exists x, (C(x) \wedge H(x))$ and $\forall x, (H(x) \rightarrow B(x))$. Using existential instantiation and the first premise $C(d) \wedge H(d)$ for some person d . From this by simplification $H(d)$ is obtained. For that person d , by universal instantiation and the second premise, $H(d) \rightarrow B(d)$. Hence by modus ponens $B(d)$ follows. Also by simplification from $C(d) \wedge H(d)$, $C(d)$ is obtained. Using the conjunction on the last two conclusions, $C(d) \wedge B(d)$ is obtained. Then applying existential generalization to this $\exists x, (C(x) \wedge B(x))$ follows.

3. Let:

- $C(x)$ be: x is a student in this class,
- $P(x)$ be: x owns a personal computer,
- and $I(x)$ be: x can use the Internet.

The premises are $\forall x, (C(x) \rightarrow P(x))$, $\forall x, (P(x) \rightarrow I(x))$, and $C(\text{John})$. Using universal instantiation and the first premise, $C(\text{John}) \rightarrow P(\text{John})$ follows. Applying modus ponens to this conclusion and the last premise, $P(\text{John})$ follows. Using universal instantiation and the second premise, $P(\text{John}) \rightarrow I(\text{John})$ follows. From the last two conclusions by modus ponens $I(\text{John})$ follows. Hence from this conclusion and the last premise by conjunction $C(\text{John}) \wedge I(\text{John})$ follows.

4. Let:

- $C(x)$ be: x is a student in this class,
- $P(x)$ be: x owns a personal computer,
- and $I(x)$ be: x has used the Internet.

The premises are $\forall x, (C(x) \rightarrow P(x))$ and $\exists x, (C(x) \wedge \neg I(x))$. Using existential instantiation and the second premise $C(d) \wedge \neg I(d)$ for some person d . Hence by simplification $C(d)$. For that person d , by universal instantiation and the first premise $C(d) \rightarrow P(d)$. From this and the previous conclusion, by modus ponens $P(d)$. Also from $C(d) \wedge \neg I(d)$ by simplification $\neg I(d)$. Hence by conjunction $P(d) \wedge \neg I(d)$. Hence by existential generalization $\exists x, (C(x) \wedge \neg I(x))$.

Check Reasoning:

1. If n is a real number with $n > 1$, then $n^2 > 1$. Suppose that $n^2 \leq 1$. Then $n \leq 1$.
Valid argument using modus tollens:

$$\neg(n^2 > 1) \rightarrow \neg(n > 1) \equiv (n^2 \leq 1) \rightarrow (n \leq 1)$$

2. If n is a real number with $n > 1$, then $n^2 > 1$. Suppose that $n^2 > 1$. Then $n > 1$.
Fallacy of affirming the conclusion.