

A Fully Lexicographic Extension of Min or Max Operation Cannot Be Associative

Olga Kosheleva and Vladik Kreinovich
University of Texas at El Paso
500 W. University
El Paso, TX 79968, USA
olgak@utep.edu, vladik@utep.edu

Abstract

In many applications of fuzzy logic, to estimate the degree of confidence in a statement $A \& B$, we take the minimum $\min(a, b)$ of the expert's degrees of confidence in the two statements A and B . When $a < b$, then an increase in b does not change this estimate, while from the commonsense viewpoint, our degree of confidence in $A \& B$ should increase. To take this commonsense idea into account, Ildar Batyrshin and colleagues proposed to extend the original order in the interval $[0, 1]$ to a lexicographic order on a larger set. This idea works for expressions of the type $A \& B$, so maybe we can extend it to more general expressions? In this paper, we show that such an extension, while theoretically possible, would violate another commonsense requirement – associativity of the “and”-operation. A similar negative result is proven for lexicographic extensions of the maximum operation – that estimates the expert's degree of confidence in a statement $A \vee B$.

1 Formulation of the Problem

Min and max operations: reminder. In fuzzy logic (see, e.g., [5, 8, 9, 10, 11, 12, 15]), the expert's degree of confidence in a statement is described by a number from the interval $[0, 1]$.

Often, we know the expert's degree of certainty a and b of statements A and B , and, based on these two values, we need estimate the expert's degree of confidence in a composite statement $A \& B$. The corresponding estimate can be denoted by $a \& b$. In many applications, we have $a \& b = \min(a, b)$.

Similarly, as an estimate $a \vee b$ for the expert's degree of confidence in a composite statement $A \vee B$, often, the max operation $a \vee b = \max(a, b)$ is used.

Need to describe a subtle difference. According to the usual min-operation, $a \& a = a$ and $a \& 1 = a$, so the value $a \& b$ remains the same when $b = a$ and

when $b = 1$. However, intuitively, if we increase our degree of confidence in a statement B , the degree of confidence in a composite statement $A \& B$ should increase. Thus, we expect that $a \& a < a \& 1$. In other words, instead of a single common value $a \in [0, 1]$, we should have different values $a \& b$ corresponding to different $b \in [a, 1]$.

In other words, we need to extend the interval $[0, 1]$ to a larger set, and extend the original order to the new set, so that:

- what was smaller remains smaller, but
- what was equal may not remain equal anymore.

How can we describe this subtle difference. How can we compare expressions $a \& b$? Since the “and”-operation is naturally commutative, we can, without losing generality, order a and b in increasing order, i.e., we can always assume that $a \leq b$.

How can we compare expressions $a_1 \& b_1$ and $a_2 \& b_2$ in which $a_1 \leq b_1$ and $a_2 \leq b_2$? If $a_1 < a_2$, then for the min-operation, we have

$$a_1 \& b_1 = a_1 < a_2 = a_2 \& b_2.$$

Since we want to retain the previous order, we thus conclude that $a_1 \& b_1 < a_2 \& b_2$ in the desired extension as well.

If $a_2 < a_1$, then similarly, we should have $a_2 \& b_2 < a_1 \& b_1$.

What if $a_1 = a_2$? In this case, for the min-operation, we get equality, but this is exactly the equality that want to clarify, so we say that $a_1 \& b_1 < a_2 \& b_2$ if $a_1 = a_2$ and $b_1 < b_2$.

This order on expressions $a \& b$ can be naturally extended to values $a \in [0, 1]$, since each such value can be described as $a \& 1$.

So, we arrive at the following *lexicographic order*: when $a_1 \leq b_1$ and $a_2 \leq b_2$, then $a_1 \& b_1 \leq a_2 \& b_2$ if and only if:

- either $a_1 < a_2$,
- or $a_1 = a_2$ and $b_1 \leq b_2$.

Such an order was first proposed in [1, 2, 3, 4, 14]. It was successfully used in applications to geosciences; see, e.g., [14].

Natural question. The idea of a lexicographic order works well for expressions of the type $a \& b$. Can we extend this idea to more general expressions?

In this paper, we show that while such an extension is possible, it is not what we look for: e.g., the corresponding operation will *not* be associative – while we want associativity $a \& (b \& c) = (a \& b) \& c$, since, from the common sense viewpoint, $A \& (B \& C)$ means exactly the same as $(A \& B) \& C$: that all three statement A , B , and C are true.

A similar result is also proven for a similar lexicographic extension of the max-operation.

2 Main Result: Case of Min Operation

Definition 1. Let (S, \leq) be a partially ordered set with the largest element 1 that contains two elements a and b for which $a < b < 1$. Let $\&$ be a commutative operation on the set S for which $a \& 1 = a$ for all a . We say that the order \leq is lexicographic if for all $a_1 \leq b_1$ and $a_2 \leq b_2$, we have $a_1 \& b_1 \leq a_2 \& b_2$ if and only if:

- either $a_1 < a_2$,
- or $a_1 = a_2$ and $b_1 \leq b_2$.

Proposition 1. When the order is lexicographic, the operation $\&$ is not associative.

Proof. Let us consider the elements $a < b < 1$ whose existence is guaranteed by the definition of lexicographic order. Then, by this definition, for $a_1 = a_2 = a$, $b_1 = b$, and $b_2 = 1$, we get

$$a \& b < a \& 1. \quad (1)$$

From $a < b$ and from the fact that $a \& 1 = a$, we conclude that

$$a \& 1 = a < b. \quad (2)$$

Now, for $a_1 = a \& b$, $a_2 = a \& 1$, $b_1 = 1$, and $b_2 = b$:

- we have $a_1 \leq b_1$ – since 1 is the largest element,
- we have $a_2 \leq b_2$ – by formula (2), and
- we have $a_1 < a_2$ – by formula (1).

So, since the order is lexicographic, we can conclude that $a_1 \& b_1 < a_2 \& b_2$, i.e., that

$$(a \& b) \& 1 < (a \& 1) \& b, \quad (3)$$

while by associativity and commutativity, we would have $(a \& b) \& 1 = (a \& 1) \& b$. Thus, the operation $\&$ is not associative.

The proposition is proven.

3 Main Result: Case of Max Operation

Definition 2. Let (S, \leq) be a partially ordered set with the smallest element 0 that contains two elements a and b for which $0 < a < b$. Let \vee be a commutative operation on the set S for which $a \vee 0 = a$ for all a . We say that the order \leq is lexicographic if for all $a_1 \leq b_1$ and $a_2 \leq b_2$, we have $a_1 \vee b_1 \leq a_2 \vee b_2$ if and only if:

- either $b_1 < b_2$,

- or $b_1 = b_2$ and $a_1 \leq a_2$.

Proposition 2. *When the order is lexicographic, the operation \vee is not associative.*

Proof. Let us consider the elements $0 < a < b$ whose existence is guaranteed by the definition of lexicographic order. Then, by this definition, for $a_1 = 0 < a_2 = a$, and $b_1 = b_2 = b$, we get

$$0 \vee b < a \vee b. \quad (4)$$

From $a < b$ and from the fact that $0 \vee b = b$, we conclude that

$$a < 0 \vee b = b. \quad (5)$$

Now, for $a_1 = a$, $a_2 = 0$, $b_1 = 0 \vee b$, and $b_2 = a \vee b$:

- we have $a_1 \leq b_1$ – by formula (5),
- we have $a_2 \leq b_2$ – since 0 is the smallest element, and
- we have $b_1 < b_2$ – by formula (4).

So, since the order is lexicographic, we can conclude that $a_1 \vee b_1 < a_2 \vee b_2$, i.e., that

$$a \vee (0 \vee b) < 0 \vee (a \vee b), \quad (6)$$

while by associativity and commutativity, we would have $a \vee (0 \vee b) = 0 \vee (a \vee b)$. Thus, the operation \vee is not associative.

The proposition is proven.

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