How to Simulate a Turing Machine?

What do we need to describe a Turing machine? For simplicity, let us consider Turing machines for acceptance and rejection, not for computations. So what do we need in this case?

- First, we need to know how many states the head will have. Let us denote this number by $N$.
- For simplicity of simulation, let us not worry about fancy names for the states, let us denote these states by $q_0, q_1, \ldots, q_{N-1}$.
- In the simulating program, we will simply represent each state by the corresponding index:
  - the state $q_0$ will be represented as 0,
  - the state $q_1$ will be represented as 1, etc.
- We need special states: start, accept, and reject. Let us assume:
  - that $q_0$ is the start state,
  - that $q_{N-2}$ is the accept state, and
  - that $q_{N-1}$ is the reject state.
- We also need to know what symbols are allowed. First, we need to know the number $M$ of the symbols.
- For simplicity, let us denote the symbols simply as $s_0, s_1, \ldots, s_{M-1}$.
- In the simulating program, we will simply represent each symbol by the corresponding index:
  - the symbol $s_0$ will be represented as 0,
  - the symbol $s_1$ will be represented as 1, etc.
- We need a special symbol “blank”. Let us assume that this symbol is $s_0$.

At each moment of time, we also need to describe which symbol is in each cell. This can be describe by an integer array $\text{tape}[n]$, so that:

- the value $\text{tape}[0]$ describes the contents of the first cell 0; for example, if in this cell, we have a symbol $s_2$, then we take $\text{tape}[0] = 2$;
• the value \texttt{tape[1]} describes the contents of the second cell 1; for example, if in this cell, we have a symbol \(s_5\), then we take \texttt{tape[1]} = 5; etc.

At each moment of time, we also need to know:

• in what state the head is; we will denote this state by \texttt{head}, and

• at what cell the head is; we will denote this location by a variable \texttt{location}. In the beginning, the head point to the first cell 0, so in the beginning, we have \texttt{location} = 0.

We also need to describe how the configuration changes. As we have mentioned, when the head in state \(q_n\) sees a symbols \(s_m\), it can do three things:

• it can change its state;

• it can replace the original symbol with some other symbol; and also

• it can move one step to the left (L) or to the right (R) or stay in place.

This behavior can be described by three 2-D arrays:

• An integer array \texttt{state[n][m]} that describes to what state the head of the Turing machine changes if it was in the state \(q_n\) and it sees the symbol \(s_m\).

  If we have reached an accept or reject state, i.e., \(n = N - 2\) or \(n = N - 1\), the Turing machine stops, so we only need to describe the values \texttt{state[n][m]} for \(n < N - 2\). In other words, we can define this array as

  \[
  \text{int[][] state} = \text{int[N - 2][M]};
  \]

• An integer array \texttt{symbol[n][m]} that describes what symbol will be placed on the tape when the head in the state \(q_n\) sees the symbol \(s_m\) (it may be the same symbol as before, or it may be some other symbol written by the Turing machine).

• Finally, a character array \texttt{lr[n][m]} that describes, for each state \(q_n\) and for each symbol \(s_m\), whether the Turing machine:

  - moves to the left (L),
  - moves to the right (R), or
  - stays in place (this will be described by a blank symbol).

\textbf{Example 1.} Le us assume that the Turing machine has the following rule:

\[
\text{back, a} \rightarrow \text{test, b, L}
\]

This rule means that when the head in the state “back” see a symbol “a”, then:

• the state of the head changes to “test”;

• the symbol changes to “b”, and
• the head moves one step to the left.

How will rule be described in terms of the arrays? Let us assume that:

• the state “back” is No. 3 in our ordering of states, i.e., is the state $q_3$,
• the state “test” is $q_5$,
• the symbol “a” is the symbol $s_2$, and
• the symbol “b” is the symbol $s_6$.

Then:

• the value $\text{state}[3][2]$ is equal to 5, meaning that if the head is in the state $q_3$ (i.e., in the state “back”) and it sees the symbol $s_2$ (i.e., the symbol “a”), then the state changes to $q_5$ (i.e., to the state “test”);
• the value $\text{symbol}[3][2]$ is equal to 6, meaning that if the head is in the state $q_3$ (i.e., “back”) and it sees the symbol $s_2$ (i.e., “a”), then the corresponding symbol on the tape is changed to $s_6$ (i.e., to the symbol “b”); and
• the value $\text{lr}[3][2]$ is equal to ‘L’, meaning that if the head is in the state $q_3$ (i.e., “back”) and it sees the symbol $s_2$ (i.e., “a”), then the head moves one step to the left (which we describe by the symbol L).

Example 2. With the same numbering of states and symbols, if we have a slightly different rule

    back, a $\rightarrow$ test, L

in which the symbol is not changed (i.e., we keep the same symbol “a”), then:

• the values $\text{state}[3][2]$ and $\text{lr}[3][2]$ are the same as in Example 1, but
• the value $\text{symbol}[3][2]$ changes: now this value is 2, meaning that we keep the same symbol $s_2$ (i.e., “a”).

Example 3. If the rule is as follows:

    back, a $\rightarrow$ L

then the state of the head also does not change. So, in contrast to the two previous examples, we will now have the value $\text{state}[3][2]$ equal to 3 – meaning that we keep the same state $q_3$ (i.e., the state “back”).

How to simulate the Turing machine: idea. In the beginning, the head is near cell 0 and in state $s_0$. So, we must have $\text{location} = 0$ and $\text{head} = 0$:

```c
int location = 0;
int head = 0;
```
To simulate what happens next, we need a loop. As we have mentioned earlier, there are two main types of loops:

- for-loops, when we know how many iterations we need, and
- while-loops, when we do not know beforehand how many iterations we will need.

Here, clearly, in general, we do not know how many iterations we will need, so it should be a while-loop. We stop when we reach either the accept state \( N - 2 \) or the reject state \( N - 1 \). If have any state \( i < N - 2 \), we continue, so the header of the while-loop must take the following form:

\[
\text{while}(\text{head} < N - 2)\
\ldots
\]

What should we have inside the loop? The head is at the cell \texttt{location}, the symbol it sees is the symbol \texttt{tape[location]}; so we can write

\[
\text{currentSymbol} = \text{tape[location]};
\]

Depending on the state \texttt{head} and on the symbol \texttt{currentSymbol}, we change the symbol \texttt{tape[location]} to \texttt{symbol[head][currentSymbol]}:

\[
\text{tape[location]} = \text{symbol[head][currentSymbol]};
\]

We also change the state and the location:

\[
\text{newHead} = \text{state[head][currentSymbol]};
\]
\[
\text{if}(\text{lr[head][currentSymbol]} == \text{'R'}) \{\text{location}++;\}
\]
\[
\text{else if}(\text{lr[head][currentSymbol]} == \text{'L'}) \{\text{location}--;\}
\]
\[
\text{head} = \text{newHead};
\]

At the end, if we reach the state \( N - 2 \), we accept, else we reject.

So, we get the following pseudo-code:

```
public static Boolean check(){
    int location = 0;
    int head = 0;
    int currentSymbol;
    while(head < N - 2){
        currentSymbol = tape[location];
        tape[location] = symbol[head][currentSymbol];
        newHead = state[head][currentSymbol];
        if(lr[head][currentSymbol] == 'R') {location++;}
        else if(lr[head][currentSymbol] == 'L') {location--;}
        head = newHead;
        if(head == N - 2){return true;}
        else{return false;}
    }
}
```