

# Once We Know that a Polynomial Mapping Is Rectifiable, We Can Algorithmically Find a Rectification

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**Problem.** It is known that some polynomial mappings  $\varphi : \mathbb{C}^k \rightarrow \mathbb{C}^n$  are *rectifiable* in the sense that there exists a polynomial mapping  $\alpha : \mathbb{C}^n \rightarrow \mathbb{C}^n$  whose inverse is also polynomial and for which  $\alpha(\varphi(z_1, \dots, z_k)) = (z_1, \dots, z_k, 0, \dots, 0)$  for all  $z_1, \dots, z_k$ . In many cases, the existence of such a rectification is proven indirectly, without an explicit construction of the mapping  $\alpha$ .

**Our first result.** In this talk, we use Tarski-Seidenberg algorithm (for deciding the first order theory of real numbers) to design an algorithm that:

- given a polynomial mapping  $\varphi : \mathbb{C}^k \rightarrow \mathbb{C}^n$  which is known to be rectifiable,
- returns a polynomial mapping  $\alpha : \mathbb{C}^n \rightarrow \mathbb{C}^n$  that rectifies  $\varphi$ .

**Our second result.** The above general algorithm is not practical for large  $n$ , since its computation time grows faster than  $2^{2^n}$ . To make computations more practically useful, for several important case, we have also designed a much faster alternative algorithm.