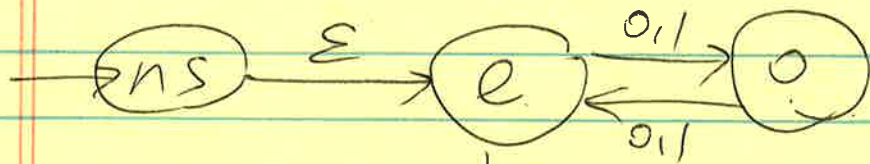
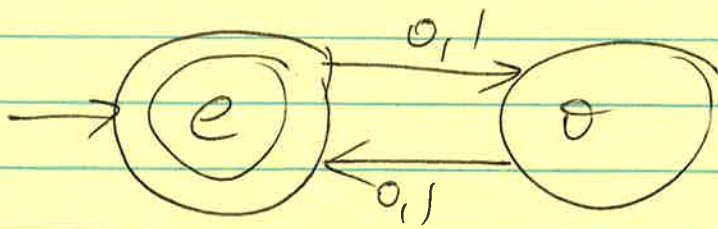
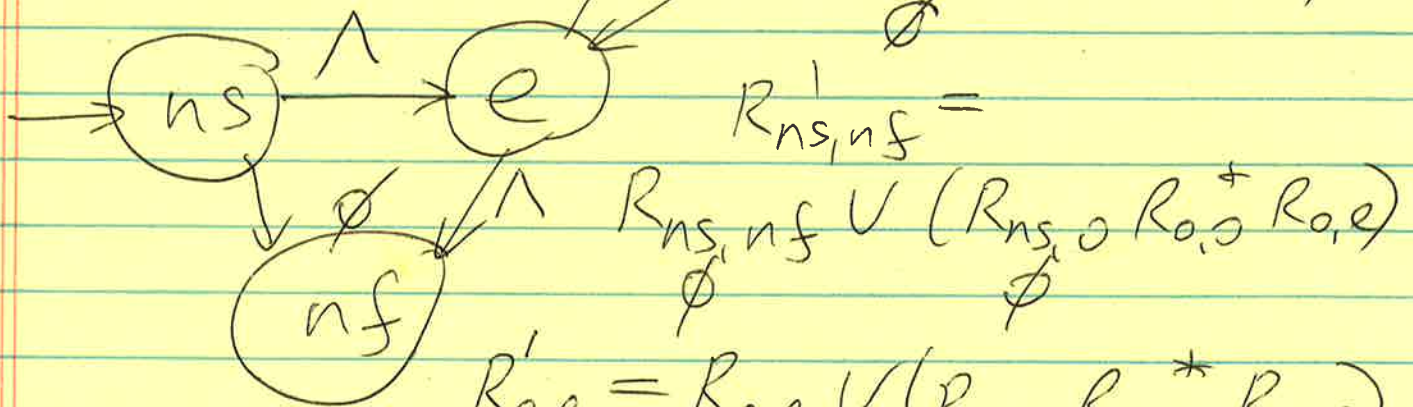


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 (1-)

$$R'_{ij} = R_{ij} \cup (R_{ik} R_{kk}^* R_{kj})$$



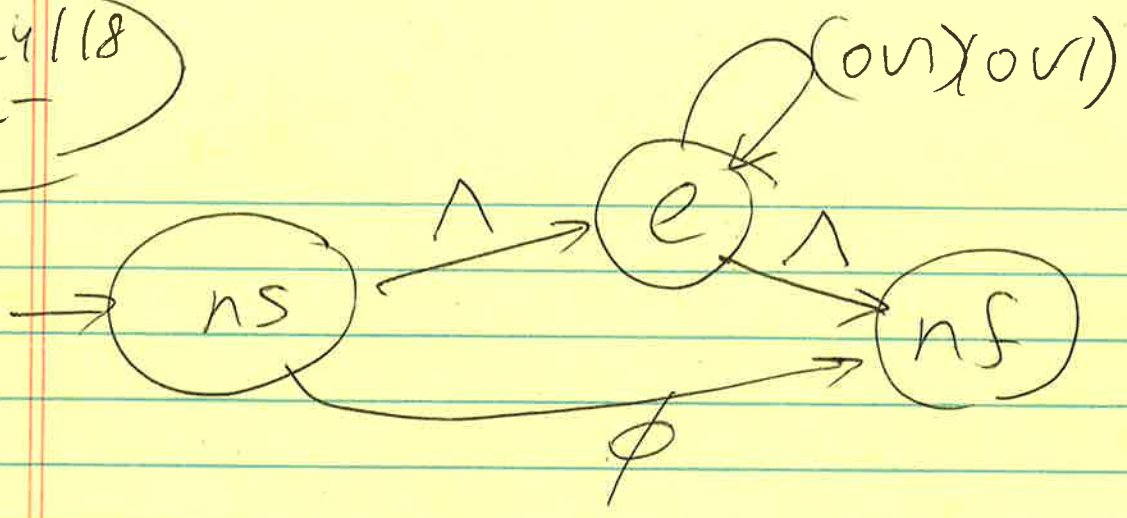
$$R'_{ns,e} = R_{ns,e} \cup (R_{ns,o} R_{o,o}^* R_{o,e})$$



$$R'_{ee} = R_{ee} \cup (R_{e,o} R_{o,o}^* R_{o,e})$$

$$R'_{e,ns} = R_{e,ns} \cup (R_{e,o} R_{o,o}^* R_{o,ns})$$

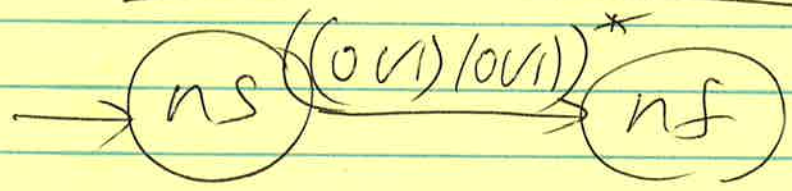
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$$R'_{ns,nf} = R_{ns,nf} \cup (R_{ns,e} R_{e,e}^* R_{e,nf})$$

$\emptyset \qquad \qquad \qquad \Lambda((ou)(ovi)(ovi))^* \Lambda$

$$\boxed{((ou)(ovi)(ovi))^*}$$



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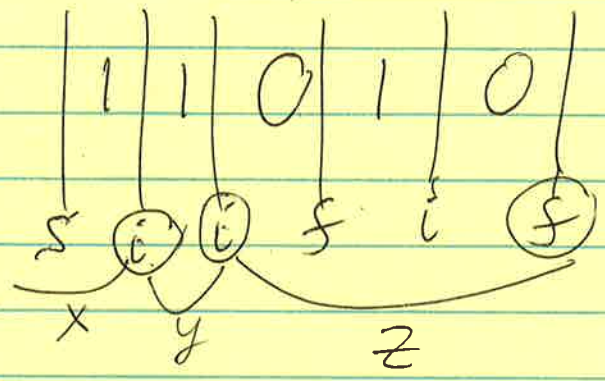
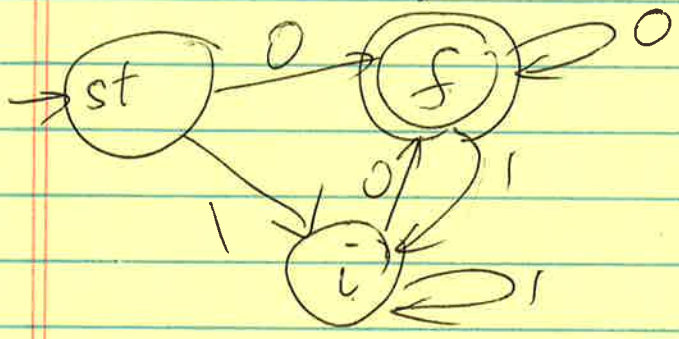
Pumping Lemma.

$\forall \text{reg. } L \exists p \forall w \in L (\text{len}(w) \geq p \rightarrow$
 $\exists x, y, z (w = xyz \ \& \ \text{len}(y) > 0 \ \& \ \text{len}(xy) \leq p$
 $\ \& \ \forall i (xy^i p \in L)))$

For every regular language L there exists a natural number p such that every word w from L whose length is at least p can be represented as $w = xyz$, where y is non-empty, length of xy is smaller than or equal to p , and for every i , the word $x \underbrace{y \dots y}_i z$ is in L .

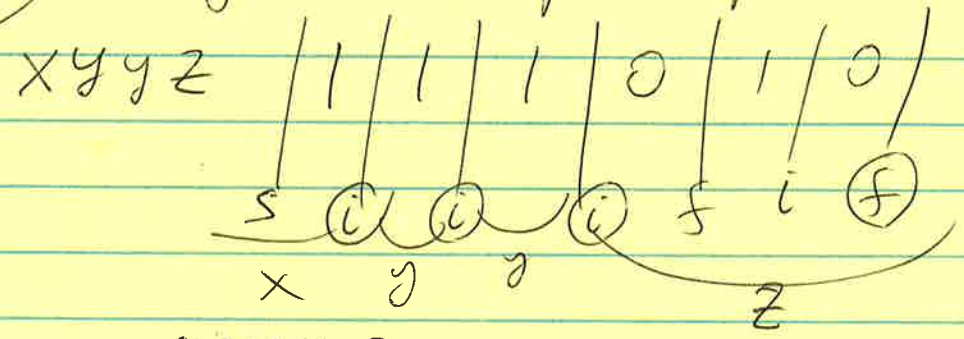
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Proof: $p = \#$ of states in
 we look for first time ~~when state is repeated~~ FA
 x - everything before the first
 repeating pair of states
 y - everything between the first
 & second state repeating states
 z - everything after



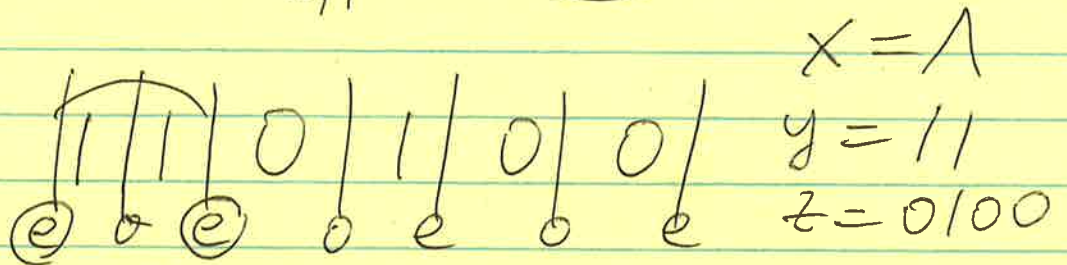
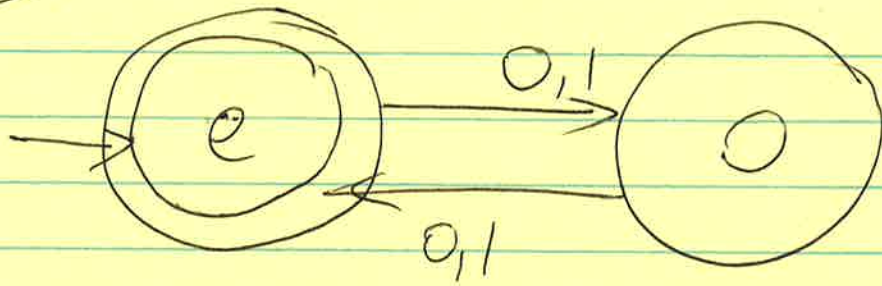
$x=1$
 $y=11$
 $z=010$

Pigeonhole principle



$xyyz$
 $xyyz$
 $xyyz$
 $xyyyyz$
 $...$

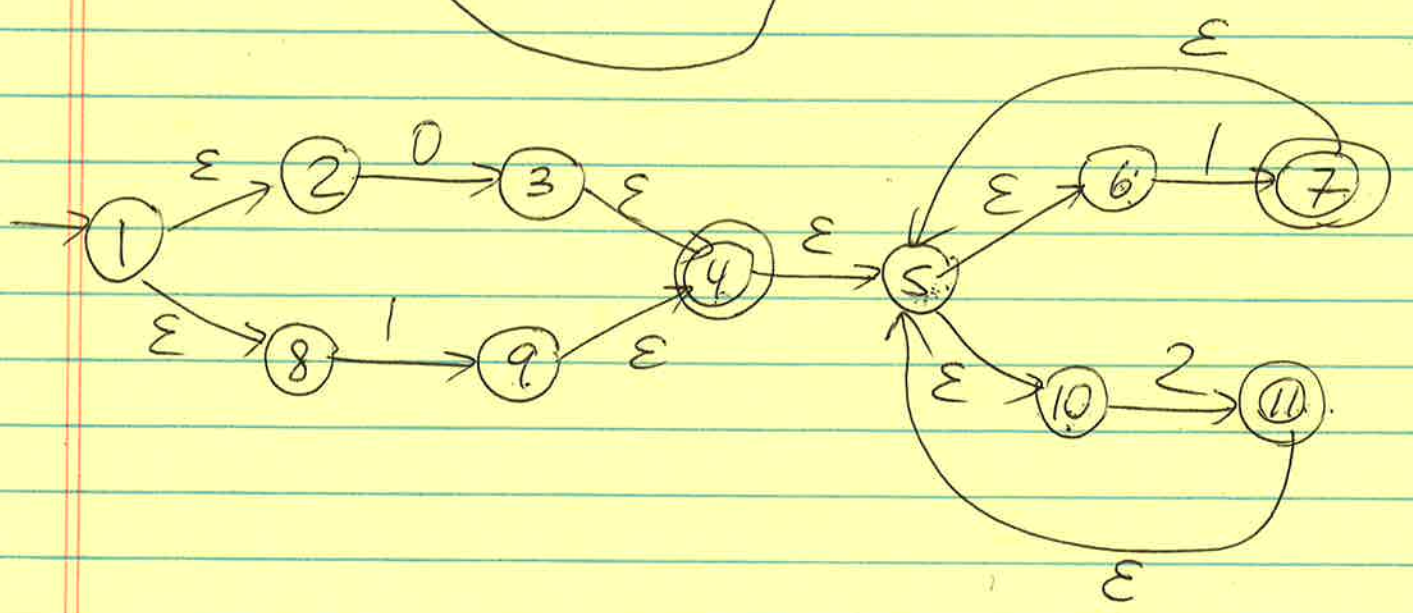
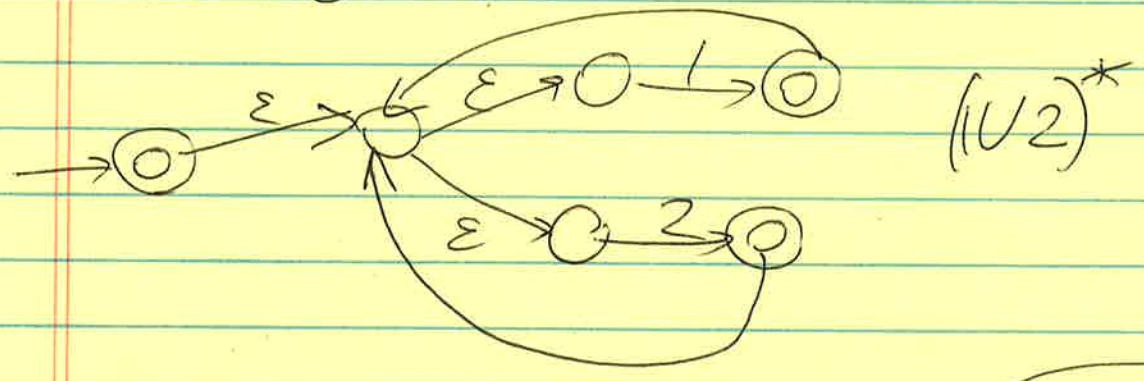
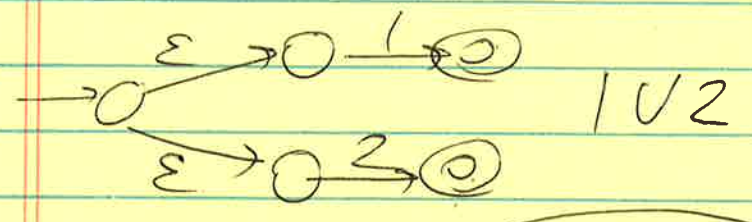
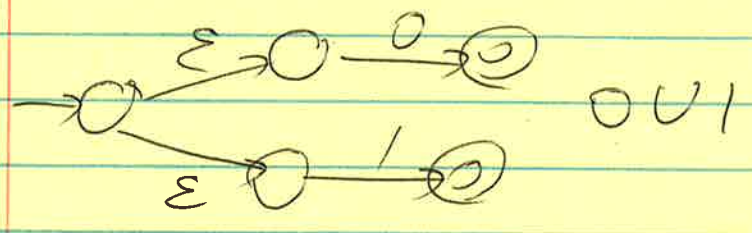
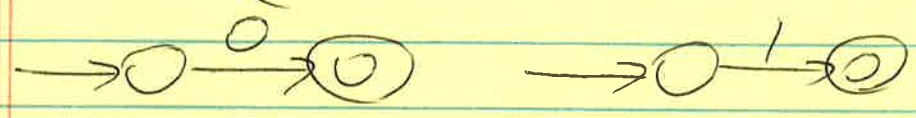
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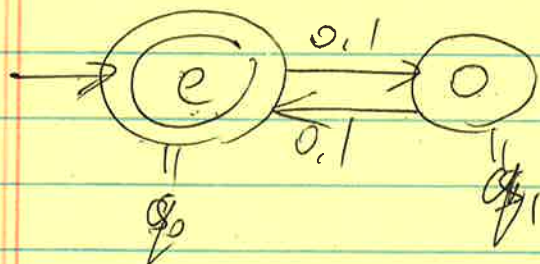
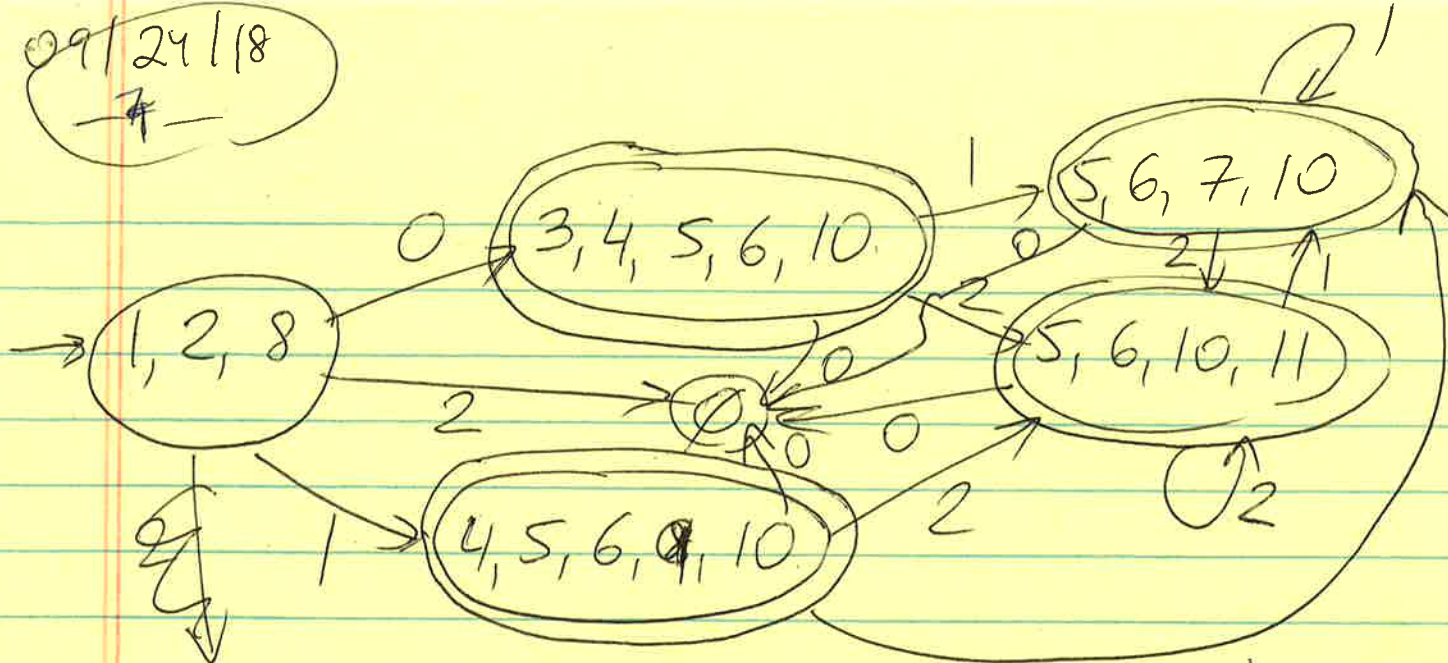
- 1) find x, y, z
- 2) show that $xyyz$ is accepted by this automaton

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$(0V1)(1V2)^*$



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$N=2, \quad q_0, q_1$
 $M=2, \quad s_0, s_1$

state [0][0] = 1

state [1][0] = 0

state [0][1] = 1

state [1][1] = 0

final [0] = true
 final [1] = false