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N - # of states q_0, \dots, q_{N-1} ^{initial}

M - # of symbols s_0, \dots, s_{M-1}
state[n][m] q_n moves if it sees symbol s_m
final[n] - boolean

```
public static void main (String [] args) {  
    int N; i  
    S.o.p ("Please enter # of states");  
    N = reader.next Int ();  
    int M;  
    ...  
  
    int [] [] state = new int [N] [M];  
    for (int n=0; n < N; n++) {  
        for (int m=0; m < M; m++) {  
            S.o.p ("what state do you get in"  
                "if you are in state" + n +  
                "and you see symbol" + m);  
            state [n] [m] = reader.next Int ();  
        }  
    }  
}
```

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```
boolean done = false;
while (!done) {
    int length;
    S.o.p. ask
    ...
    int[] word = ...,
    ask ...

    int currentState = 0;
    for (int i = 0; i < length; i++) {
        currentState =
            state[currentState][word[i]];
        if (final[currentState])
            { S.o.p. ("Accepted"); }
        else { S.o.p. ("Rejected"); }
        S.o.p. ("Do you want to continue?");
        if ( ...
    }
}
```

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Pumping lemma.

$$\forall \text{ reg. } L \exists p \forall w \left(\text{len}(w) \geq p \rightarrow \exists x, y, z \right. \\ \left. (w = xyz \ \& \ \text{len}(y) > 0 \ \& \ \text{len}(xy) \leq p \ \& \right. \\ \left. \forall i (xy^i z \in L)) \right)$$

Th. $L = \{a^n b^n : n = 0, 1, 2, \dots\} =$

Theorem $\{\Lambda, ab, aabb, aaa bbb, \dots\}$

is not regular.

Proof: by contradiction. Let's assume

that L is regular, and let's derive a contradiction from this assumption.

Since L is regular, by pumping lemma there exist an integer p such that every word w from L whose length is $\geq p$ can be represented as xyz , where $\text{len}(y) > 0$, $\text{len}(xy) \leq p$, and for all i , $xy^i z \in L$.

Let's take $w = a^p b^p = \underbrace{a \dots a}_{p \text{ times}} \underbrace{b \dots b}_{p \text{ times}}$.

$\text{len}(xy) \leq p$, $w = xyz$ starts with xy

So y is in a 's, in $xyyz$ - we add a 's but not b 's & we started with word that has same # of a 's and b 's, so $xyyz$ has more a 's than b 's, so $xyyz \notin L$. But by pumping lemma $xyyz \in L$. We get a contradiction.

So our assumption - that L is regular - is false. Thus, L is not regular.

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$L = \{ \text{all the words that have same number of a's and b's} = \{ \epsilon, ab, ba, aabb, abba, abab, \dots \}$

$L = \{ a^{n+1} b^n \}$

$\{ a^n b^n \}$

$a^p b^p$

~~aaab~~

~~abab~~

a...a b...b

$\text{len}() = 2p \geq p$

~~$a^p b^p$~~

~~$2p \geq p$~~

~~$\text{len}(xy) \leq p$~~

~~$a...a$~~ ~~$b...b$~~
 ~~$p-1$~~ ~~p~~

$\{ a^p b^{p+1} \}$ ϵ

$\{ a^n b^{2n} \}$

~~$a^1 b^2$~~

~~abb~~

$a^p b^{2p}$

$L = \{ \boxed{WW} \} = \{ \text{catcat, dogdog, bagbag, } \dots \}$

~~$W = \{ a^p b^p a^p b^p \}$~~ $= xyz$

a...a b...b a...a b...b

$a^{p+z} b^p$ $a^p b^p$ xyz

$\text{len}(xy) \leq p$