

Solution to Problem 15

Task. Prove that the language consisting of all the words that have equal number of +’s and –’s and exactly two times as many opening parentheses is not context-free.

Solution. Let us prove that this language is not context-free.

The proof will be by contradiction. Let us assume that this language is context-free. Then, by pumping lemma, there exists an integer p such that every word w from the language L whose length is at least p can be represented as $w = uvxyz$, where:

- $\text{len}(vxy) \leq p$;
- $\text{len}(vy) > 0$, and
- for all natural numbers i , the word $uv^i xy^i z$ also belongs to the language L .

Let us take the word

$$w = +^p -^p (^{2p} = + \dots + - \dots - (\dots (,$$

where + and – are repeated p times each, and (is repeated $2p$ times. The length of this word – i.e., the number of symbols in this word – is equal to $p + p + 2p = 4p$. Clearly, $4p \geq p$, so, according to the Pumping Lemma, this word can be described as $uvxyz$ with the above properties.

Where can the central part vxy of this word be? We know that the length $\text{len}(vxy)$ of this part cannot exceed p . Thus, it cannot contain all three types of symbols: +’s, –’s, and (’s – since then it would have to include all p letters – plus additional – and (symbols, so its length would have been larger than p . So, there are only 5 cases remaining for the location of the part vxy :

1. it can be in the +’s;
2. it can be in +’s and –’s;
3. it can be in –’s;
4. it can be in –’s and (’s; and
5. it can be in (’s.

Let us consider these cases one by one.

Case 1. If vxy is in $+$'s, this means that the parts v and y contain only $+$'s. Thus, when we pump, i.e., when we go from the original word $uvxyz$ to the word $uv^2xy^2z = uvvxyyz$, we add $+$'s – but we do not add any $-$'s or $()$'s. In the original word $w = +^p -^p ({}^{2p}$, there were exactly as many $+$'s as $-$'s. When we add more $+$'s, the balance is disrupted. Since the language L only contains the words which have the exact same number of $+$'s and $-$'s, the word $uvvxyyz$ cannot belong to the language L .

Case 2. If vxy is in $+$'s and in $-$'s, this means that the parts v and y contain only $+$'s and $-$'s. Thus, when we pump, i.e., when we go from the original word $uvxyz$ to the word $uv^2xy^2z = uvvxyyz$, we add $+$'s and $-$'s – but we do not add any $()$'s. In the original word $w = +^p -^p ({}^{2p}$, there were exactly two times as many $()$'s as $+$'s and as $-$'s. When we add more $+$'s and $-$'s, the balance is disrupted. Since the language L only contains the words which have exactly two times as many $()$'s than $+$'s, the word $uvvxyyz$ cannot belong to the language L .

Case 3. If vxy is in $-$'s, this means that the parts v and y contain only $-$'s. Thus, when we pump, i.e., when we go from the original word $uvxyz$ to the word $uv^2xy^2z = uvvxyyz$, we add $-$'s – but we do not add any $+$'s or $()$'s. In the original word $w = +^p -^p ({}^{2p}$, there were exactly as many $-$'s as $+$'s. When we add more $-$'s, the balance is disrupted. Since the language L only contains the words which have exactly as many $-$'s as $+$'s, the word $uvvxyyz$ cannot belong to the language L .

Case 4. If vxy is in $-$'s and $()$'s, this means that the parts v and y contain only $-$'s and $()$'s. Thus, when we pump, i.e., when we go from the original word $uvxyz$ to the word $uv^2xy^2z = uvvxyyz$, we add $-$'s and $()$'s – but we do not add any $+$'s. In the original word $w = +^p -^p ({}^{2p}$, there were exactly as many $+$'s as $-$'s and exactly two times more $()$'s than $+$'s. When we add more $-$'s and $()$'s, the balance is disrupted. Since the language L only contains the words which have the exact same number of $+$'s and $-$'s and two times more $()$'s, the word $uvvxyyz$ cannot belong to the language L .

Case 5. If vxy is in $()$'s, this means that the parts v and y contain only $()$'s. Thus, when we pump, i.e., when we go from the original word $uvxyz$ to the word $uv^2xy^2z = uvvxyyz$, we add $()$'s – but we do not add any $+$'s or $-$'s. In the original word $w = +^p -^p ({}^{2p}$, there were exactly two times as many $()$'s as $+$'s. When we add more $()$'s, the balance is disrupted. Since the language L only contains the words which have exactly two times more $()$'s than $+$'s, the word $uvvxyyz$ cannot belong to the language L .

So, in all five cases, we get a contradiction. This means that the original assumption – that the language L is context-free – is wrong. Thus, the language L is not context-free.