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SOLUTION TO
PROBLEM 11
FALL 2024

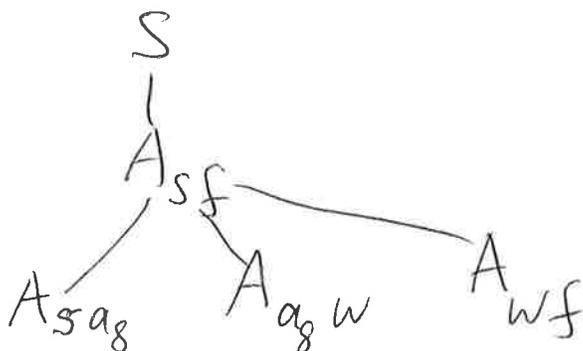
Acceptance of the word RSR has the following form (see solution to Problem 6):

read			R		S			R	
state	s	w	a_1	w	a_7	a_8	w	w	f
stack		\$	\$	R	\$		N		
				\$					

There is only one final state f in this PDA, so we start with rule



In the tracing of acceptance, there are two intermediate states with empty stack, so we use transitivity rule:



② Let us deal with the three transitions A_{sa_8} , A_{a_8w} , and A_{w_1s} one by one.

A_{sa_8} . This transition has the following form:

		R		S	
S	w	a_1	w	a_7	a_8
	\$	\$	R	\$	
			\$		

The first symbol we push is \$, so let us write down the rules used to push \$ and then to pop \$:



By using the general pattern, we get the rule $A_{sa_8} \rightarrow \epsilon A_{wa_7} \epsilon$, i.e., the rule $A_{sa_8} \rightarrow A_{wa_7}$. Let us write

it down:

$$A_{sa_8} \rightarrow A_{wa_7}$$

③

We took care of $\$$ and of the two transitions:

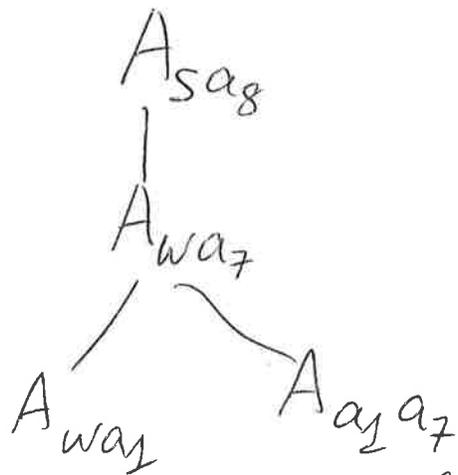
S	w	R a_1	w	S a_7	a_8
	$\\$	$\\$	R $\\$	$\\$	

Now, we need to describe the transition from w to a_7 .

Here, we again have an intermediate state (a_1) with an empty stack, so we again need to use transitivity rule:

$$A_{wa_7} \rightarrow A_{wa_1} A_{a_1 a_7}$$

Now the tree gets the following form:



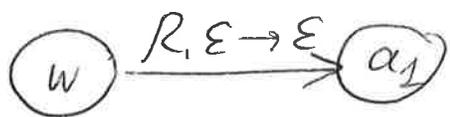
Let us consider the transitions A_{wa_1} and $A_{a_1 a_7}$ one by one

(4)

A_{w a₁}

	R
w	a ₁
⊥	⊥

For going from w to a₁, we use the rule R, ε → ε (actually, in Problem 6, we use the rule R, ⊥ → ⊥ but in reality nothing changes in the stack, so it is the same as using the rule R, ε → ε);



We have only one rule, so we add a fictitious rule



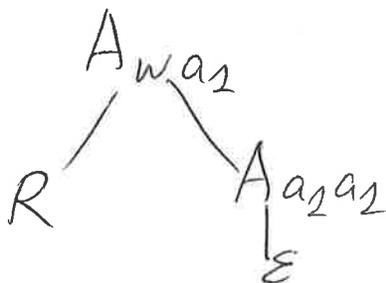
Now we can form a rule

$$A_{w a_1} \rightarrow R A_{a_1 a_1} \epsilon, \text{ i.e. } A_{w a_1} \rightarrow R A_{a_1 a_1}$$

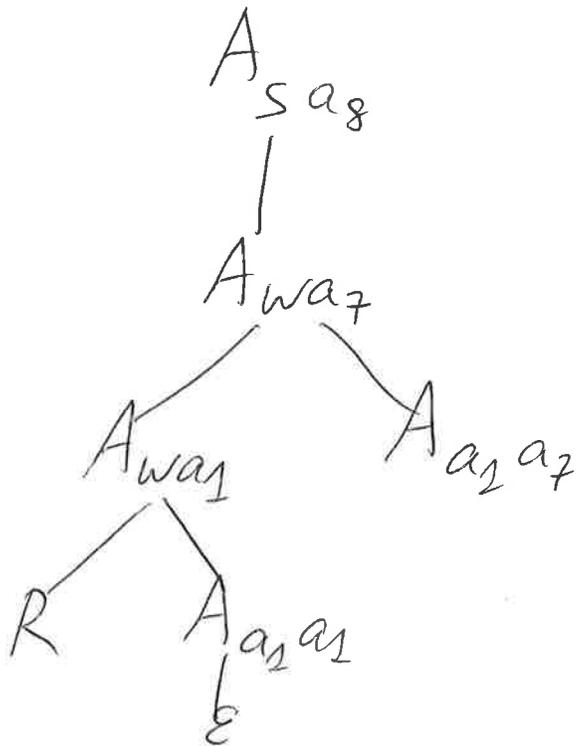
	R
w	a ₁

what remains is a transition from a₁ with empty stack to a₁ with empty stack, so we

use the rule A_{a₁ a₁} → ε. Now the transition from A_{w a₁} takes the form



⑤ and the transition A_{sa_8} takes the form



$A_{a_1a_7}$:

		S
a_1	w	a_7
	R	
\$	\$	\$

we push R,
then we pop R



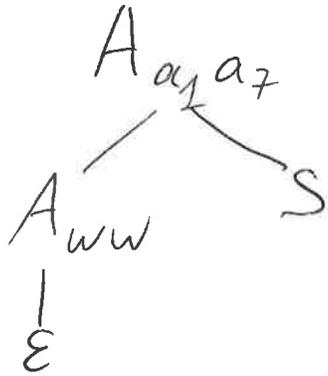
$A_{a_1a_7} \rightarrow \epsilon A_{ww} S$, i.e. $A_{a_1a_7} \rightarrow A_{ww} S$

	R	S
a_1	w	a_7
	R	
\$	\$	\$

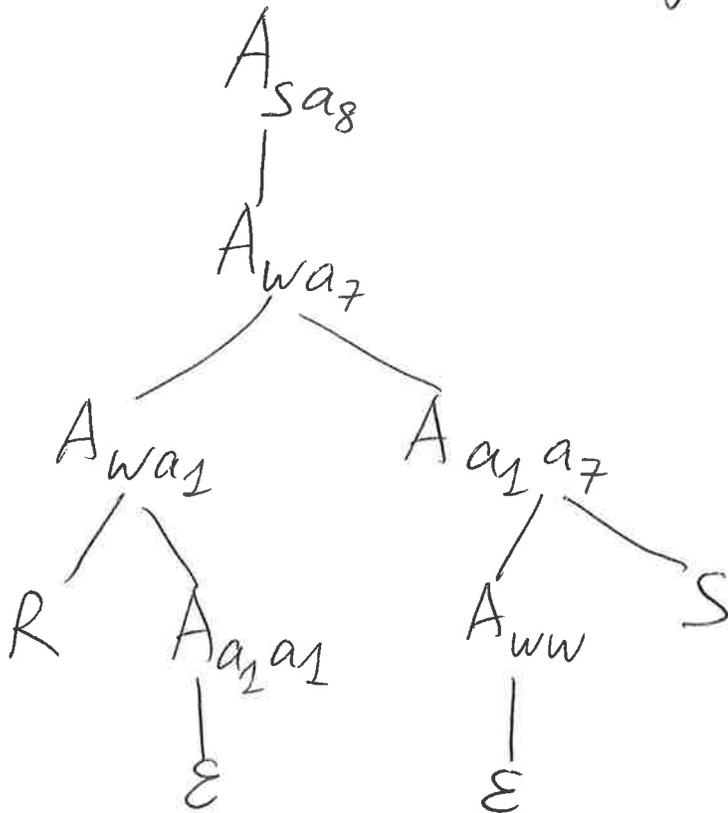
We took care of R and
of transitions, so we
add $A_{ww} \rightarrow \epsilon$

⑥

So, for $A_{a_1 a_7}$ we have:



and the transition from A_{sa_8} takes the following form:



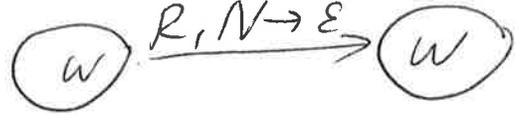
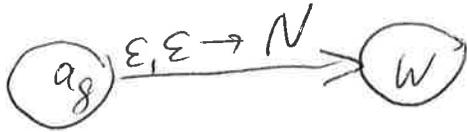
(7)

A_{agw} :

		R
a_g	w	w
	N	

we push N , then pop it.

The rules are:



So we get $A_{agw} \rightarrow \epsilon A_{ww} R$, i.e.

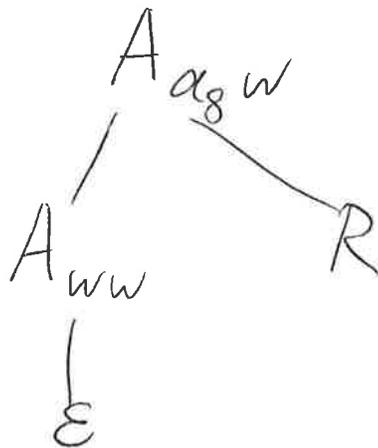
$$A_{agw} \rightarrow A_{ww} R.$$

		R
a_g	w	w
	N	

So we took care of N and of both transitions, so now we go from w with empty stack to itself; the rule is $A_{ww} \rightarrow \epsilon$.

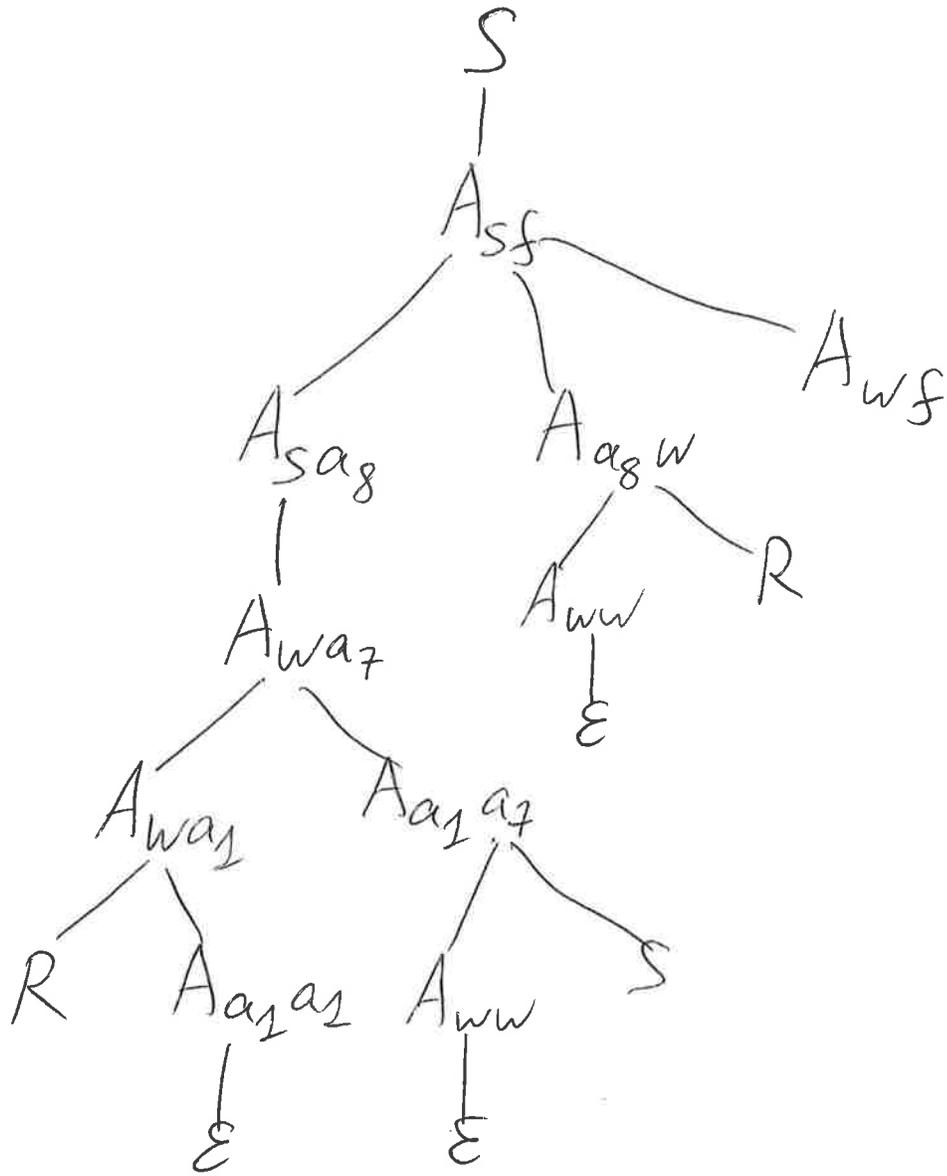
$$A_{ww} \rightarrow \epsilon.$$

So, the subtree coming from A_{agw} has the form



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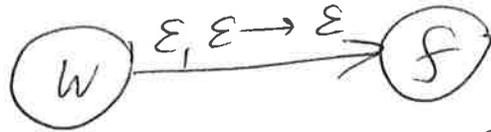
The general derivation so far has the following form:



⑨ Finally, let us deal with the transition A_{wf}

w	f

we have a single rule



So we need to supplement it with a fictitious rule



that leads to $A_{wf} \rightarrow \epsilon A_{ff} \epsilon$, i.e.

$A_{wf} \rightarrow A_{ff}$. The remaining transition is from f to f so

we use the rule

w	f

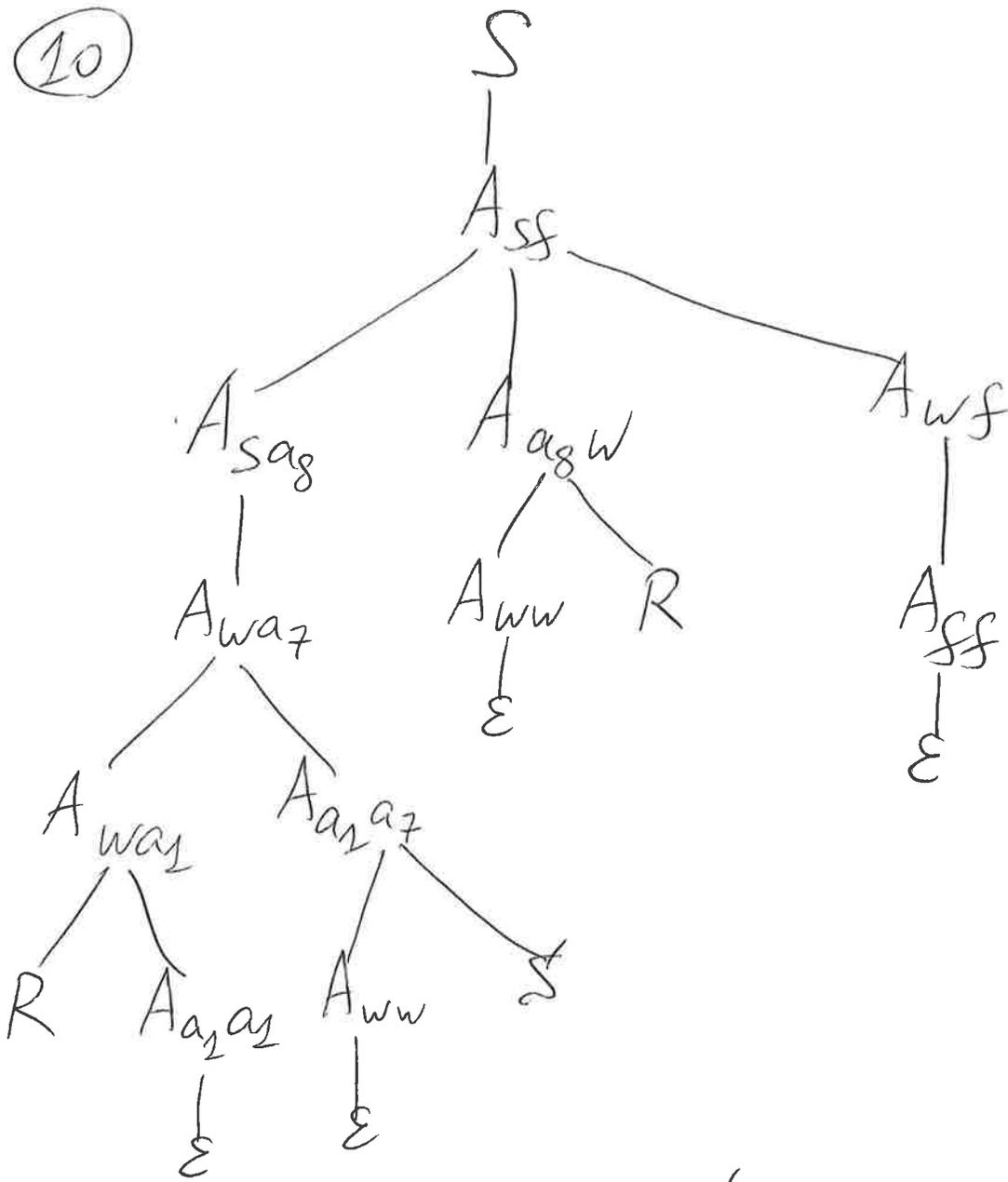
$A_{ff} \rightarrow \epsilon$

So, the tree coming from A_{wf} has the form



and the resulting generation has the form

(10)



We have indeed generated the word RSR in the corresponding grammar.