Solution to Homework 5

Task: No matter how hard you study (S), you need some rest (R). Ideally, you should study no more 8 hours a day, with at least 16 hours remaining, i.e., you should have twice as much time to rest as to study. Prove that the following language is not regular: the language L of all sequences of Ss and Rs in which there are at least twice as many Rs as Ss. For example, the word $SRSRRRR$ is in the language L, while the word $SRSRR$ is not in L.

Solution. We will prove this result by contradiction. Let us assume that the language L is regular, and let us show that this assumption leads to a contradiction.

Since the language L is regular, according to the Pumping Lemma, there exists an integer p such that every word from L whose length $\text{len}(w)$ is at least p can be represented as a concatenation $w = xyz$, where:

- y is non-empty;
- the length $\text{len}(xy)$ does not exceed p, and
- for every natural number i, the word $xy^iz \stackrel{\text{def}}{=} xy\ldots yz$, in which y is repeated i times, also belongs to the language L.

Let us take the word

$$w = S^pR^{2p} = S\ldots SR\ldots R,$$

in which first the symbol S is repeated p times and then the symbol R is repeated $2p$ times. The length of this word is $p + 2p = 3p > p$. So, by pumping lemma, this word can be represented as $w = xyz$ with $\text{len}(xy) \le p$. The word $w = xyz$ starts with xy, and the length of xy is smaller than or equal to p. Thus, xy is among the first p symbols of the word w – and these symbols are all Ss. So, the word y only has Ss.

In the original word $w = xyz$, we had S repeated p times and R repeated $2p$ times. When we go from the word $w = xyz$ to the word $xyyz$, we add Ss, and we do not add any Rs. Thus, we still have R repeated $2p$ times, but the number of times S is repeated is now larger than p. In any word from the language L, we should have at least twice as many Rs as Ss. In the word $xyyz$, this condition is not satisfied. Thus, the word $xyyz$ cannot be in the language L.

On the other hand, by Pumping Lemma, the word $xyyz$ must be in the language L. So, we proved two opposite statements:

- that this word *is not* in L and
- that this word *is* in L.

This is a contradiction.

The only assumption that led to this contradiction is that L is a regular language. Thus, this assumption is false, so the language L is not regular.