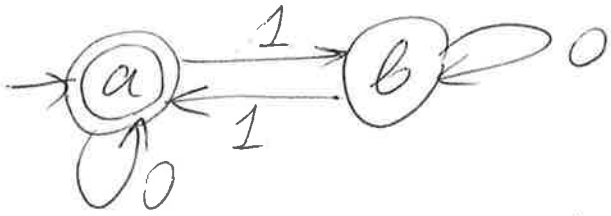


TEST 2, FALL 2025

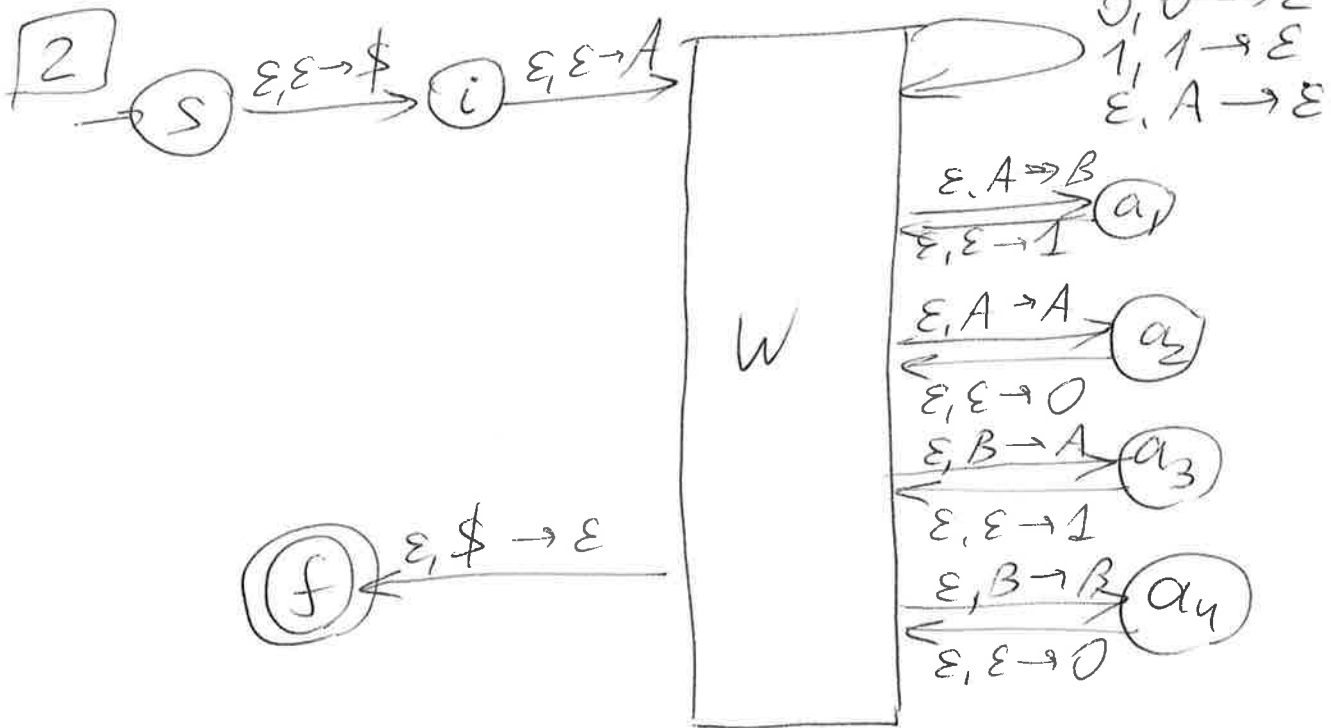
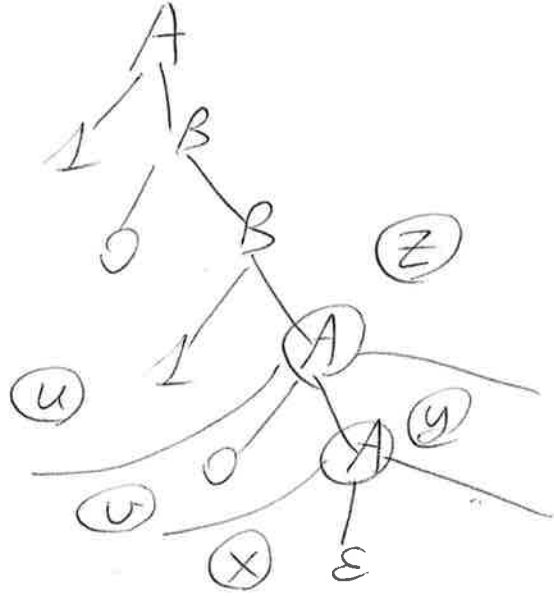


- $$\boxed{1} \quad \begin{array}{l} A \rightarrow 1B \\ A \rightarrow 0A \\ A \rightarrow \epsilon \end{array} \quad \begin{array}{l} B \rightarrow 1A \\ B \rightarrow 0B \end{array}$$

$$u = 101$$

$$v = 0$$

$$x = y = z = 1$$



read				1			0			1			0				
state	S	i	w	a ₁	w	w	a ₄	w	w	a ₃	w	w	a ₂	w	w	w	S
stack	\$	A	B	1	B	B	0	B	A	1	A	A	0	A	\$		
		\$	\$	B	\$	\$	B	\$	\$	A	\$	\$	A	\$			
				\$			\$			\$			\$				

3. $A \rightarrow IB$

2. $A \rightarrow OA$

3. $B \rightarrow IA$

4. $B \rightarrow OB$

5. $A \rightarrow \varepsilon$

Prelim-step

6. $A_0 \rightarrow A$

Step 0

7. $A \rightarrow 0$

8. $B \rightarrow 1$

9. $A_0 \rightarrow \varepsilon$

Step 1

10. $A_0 \rightarrow IB$

11. $A_0 \rightarrow OA$

12. $A_0 \rightarrow 0$

Step 2

13. $V_0 \rightarrow 0$

14. $V_1 \rightarrow 1$

15. $A \rightarrow V_1 B$

16. $A \rightarrow V_0 A$

17. $B \rightarrow V_1 A$

18. $B \rightarrow V_0 B$

19. $A_0 \rightarrow V_1 B$

20. $A_0 \rightarrow V_0 A$

Step 3

Not Needed

④

$$E \rightarrow 0$$

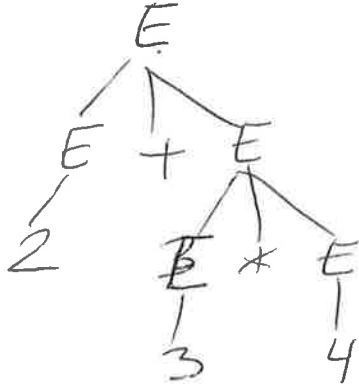
⋮

$$E \rightarrow 9$$

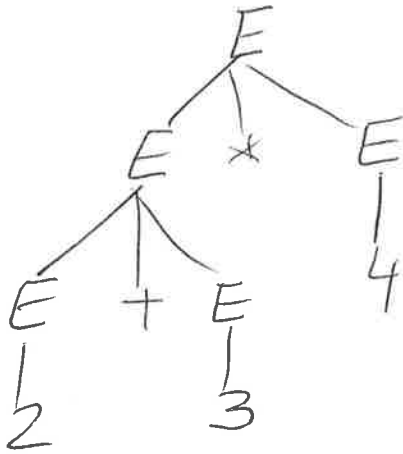
$$E \rightarrow E + E$$

$$E \rightarrow E * E$$

Then $2+3*4$ has two derivations:



← This is what we want



← this does not help in compiling

5-6

Proof by contradiction.

Let us assume that L is CFG.

Then by pumping lemma

$\exists p \forall w \in L (\text{len}(w) \geq p \rightarrow \exists u, v, x, y, z / w = uvxyz$

$\& \text{len}(vy) > 0 \& \text{len}(vxy) \leq p \& \forall i (uv^i xy^i z \in L)$

Let us take $w = c^p b^p a^p = \underbrace{c \dots c}_{p \text{ times}} \underbrace{b \dots b}_{p \text{ times}} \underbrace{a \dots a}_{p \text{ times}}$

Since $\text{len}(vxy) \leq p$, vxy cannot contain all 3 letters - otherwise it should contain at least one a , at least one c , and all p b 's - to the total of $p+2 > p$. Thus, we have the following possibilities:

- (1) vxy is in a 's,
- (2) vxy is in a 's and b 's,
- (3) vxy is in b 's
- (4) vxy is in b 's and c 's,
- (5) vxy is in c 's

In case 1, v and y only have a 's so when we go from $uvxyz$ to $uvvxyyz$ we add a 's but we don't add b 's or c 's. In $uvxyz$, we had equal number of a 's, b 's, and c 's, now balance is disrupted so $uvvxyyz \notin L$. But by pumping lemma for $i=2$, $uvvxyyz \in L$. Contradiction shows that this case is not possible.

In case 2, v and y only have a 's and b 's. So when we go from $uvxyz$ to $uvvxyyz$, we add a 's and/or b 's, but not c 's. So balance is disrupted, and $uvvxyyz \notin L$ - but by PL $uvvxyyz \in L$. Other cases are similar. In all cases we get contradiction, so L is not CFG.