

lec-03  
Date: 01/24/2017

P.r.:

+ , \* , - , / , % , if-then , or , and ,  
= , ≠ , < , ≤ , ≥ , > ,  
 $a \div b = \max(a - b, 0) = \begin{cases} a - b & \text{if } a \geq b \\ 0 & \text{otherwise} \end{cases}$

Instead of -- ,

$$\text{prev}(n) = \begin{cases} n-1 & \text{if } n \geq 1 \\ 0 & \text{otherwise} \end{cases} = n \div 1$$

So,  $\begin{cases} \text{prev}(0) = 0 \\ \text{prev}(m+1) = m \end{cases}$

$$\begin{cases} f(0) = 0 \\ f(m+1) = m \end{cases}$$

$$\begin{cases} f(0) = g() \\ f(m+1) = h(m, f(m)) \end{cases}$$

$g = 0, h = \pi_1^2$

So,  $\text{prev} = \text{PR}(0, \pi_1^2)$   
Th<sup>m</sup> prev is p.r.

Th<sup>m</sup>-2:

$\div$  is p.r

Proof:

$$a \div b = a \underbrace{\div 1 \div 1 \dots \div 1}_{b \text{ times}}$$

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int sub = a ;
for (int i=1 ; i <= b ; i++)
  sub = prev(sub) ; sub - 1 ;
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14

~~sub~~(~~sub~~)

$$\text{sub}(a, 0) = a$$

$$\text{sub}(a, m+1) = \text{prev}(\text{sub}(a, m))$$

$$f(n_1, 0) = n_1$$

$$f(n_1, m+1) = \text{prev}(f(n_1, m))$$

$$f(n_1, 0) = g(n_1)$$

$$f(n_1, m+1) = h(n_1, m, f(n_1, m))$$

$$g(n_1) = n_1 = \pi_1^1$$

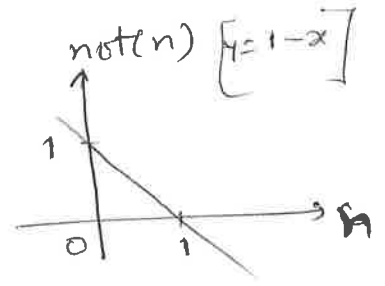
$$h = \text{prev}(f(n_1, m))$$

$$= \text{prev}(\pi_3^3)$$

$$\therefore = \text{PR}(\pi_1^1, \text{prev}(\pi_3^3))$$

Java boolians are: not, and, or  
 !, &&, ||

not(0) = 1  
 not(1) = 0



So,  $not(n) = 1 - n$   
 $= \sim(0) - n$   
 not is p.r.

and (a, b)

b/a	0	1
0	0	0
1	0	1

In engineering,  
 it's multiplication

and is p.r.

or (a, b)

it's addition

b/a	0	1
0	0	1
1	1	1

~~$a \vee b = \neg(\neg a \ \& \ \neg b)$~~

$a \vee b = \neg(\neg a \ \& \ \neg b)$  → de Morgan law  
 or (a, b) = not (and (not(a), not(b)))  
 $= 1 - (1 - a) * (1 - b) = 1 - (1 - a)(1 - b)$   
 not is p.r.  
 and is p.r.  
 $= 1 - (1 - a - b + a * b)$   
 $= a + b - a * b$   
 so, or is p.r.

16

$$\text{eq to } O(n) = \begin{cases} \text{true} & \text{if } n=0 \\ \text{false} & \text{otherwise} \end{cases}$$

$$\text{eq to } O(0) = 1$$

$$f(0) = 1$$

$$\text{eq to } O(m+1) = 0$$

$$f(m+1) = 0$$

$$f(0) = g()$$

$$f(m+1) = h(m, f(m))$$

$$g = 1, h = 0$$

$$= \sigma \circ 0$$

$$\text{eq to } O = PR(\sigma \circ 0, 0)$$

If  $a = 0$ , how <sup>will we define</sup>  $a = b$ ?

$$a = b \Rightarrow a - b = 0 \Leftrightarrow a \leq b,$$

$$\text{bec } a - b = \begin{cases} a - b, & \text{if } a > b \\ 0, & \text{otherwise} \end{cases}$$

$$a \leq b \Leftrightarrow \text{eq to } O(a - b)$$

So,  $\leq$  is p.r.

$$a = b \Leftrightarrow \underbrace{a \leq b}_{\text{p.r.}} \ \&\& \ \underbrace{b \leq a}_{\text{p.r.}}$$

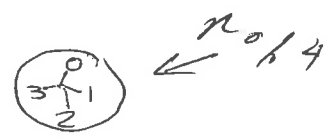
So,  $=$  is p.r.

$$a \neq b \Leftrightarrow \neg(a = b)$$

So,  $\neq$  is p.r.

$$a < b \Leftrightarrow a \leq b \ \&\& \ \neg a = b$$

So,  $<$  is p.r.



$0 \% a$

$$\text{rem}(a, 0) = 0$$

$$\text{rem}(a, m+1) = \begin{cases} \text{rem}(a, m) + 1 & \text{if } \text{rem}(a, m) + 1 < a \\ 0 & \text{otherwise} \end{cases}$$

m	0	1	2	3	4	5	6	7	8	...
m/4	0	1	2	3	0	1	2	3	0	...

Th<sup>m</sup>: IF  $P, g, h$  are p.r. then  
 if  $P(\bar{n})$  then  $g(\bar{n})$  else  $h(\bar{n})$  is also p.r.

Brainstorm:

either  $P$  is true and we have  $g(\bar{n})$   
 or  $P$  is false & " "  $h(\bar{n})$

$\text{abs}(n) = (n \geq 0) : n \quad -n$

$P(\bar{n}) * g(\bar{n}) \quad + \quad (1 - P(\bar{n})) * h(\bar{n})$

IF  $P(\bar{n})$  is true, then  $g(\bar{n})$   
 "  $P(\bar{n})$  " false, "  $h(\bar{n})$

So, if-then-else is p.r.

if  $(\text{rem}(a, m) + 1 < a)$   $\{ \text{rem}(a, m) + 1 \}$   
 else  $(0)$

So, rem is p.r.

$m/a$

$$\text{div}(a, 0) = 0$$

$$\text{div}(a, m+1) = \begin{cases} \text{div}(a, m) + 1 & \text{if } \text{rem}(a, m+1) = 0 \\ \text{div}(a, m) & \text{otherwise} \end{cases}$$

$m$	0	1	2	3	4	5	6	7	8	9	...
$m/4$	0	1	2	3	0	1	2	3	0	1	...
$m/4$	0	0	0	0	1	1	1	1	2	2	...

### HW-2

1. if  $p$  then  $f$ , else if  $q$  then  $g$   
else  $h$

2. Prove from scratch that  
remainder is p.r.