

# Computing at Least One of Two Roots of a Polynomial is, in General, not Algorithmic

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## 1. Overview

- In our previous work, we provided a theoretical explanation for an empirical fact that:
  - it is easier to find a unique root
  - than to find multiple roots.
- In this paper, we strengthen that explanation.
- Namely, we show that finding one of many roots is also difficult.

## 2. Background

- For each class of problems:
  - before we start designing an algorithm for solving all problems from this class,
  - it is desirable to know whether such a general algorithm is indeed possible.
- Among these problems are:
  - solving systems of equations (in particular, equations),
  - finding optima of a given function on a given box, etc.
- It is known that there exists an algorithm which is applicable:
  - to every system  $f_1(x_1, \dots, x_n) = 0, \dots, f_m(x_1, \dots, x_n) = 0$  with computable functions  $f_i$
  - which has exactly one solution on a given box  $\mathbf{x}_1 \times \dots \times \mathbf{x}_n$ .

### 3. Background (cont-d)

- This known algorithm computes the corresponding solution. solution.
- To be more precise, it produces, for every  $\varepsilon > 0$ , an  $\varepsilon$ -approximation to the solution.
- It is also known that no algorithm is possible:
  - which is applicable to every computable system which has exactly two solutions and
  - which would return *both* solutions.
- The proof shows that such an algorithm is not possible even for computable polynomial equations.

#### 4. Mathematical comment

- This algorithmic impossibility is due to the fact that we allow computable polynomials.
- For polynomials with rational (or even algebraic) coefficients, solution problems are algorithmically decidable.

## 5. Practical comment

- This result is in good accordance with the empirical fact that in general:
  - it is easier to find a point  $(x_1, \dots, x_n)$ , in which a given system of equations has a unique solution
  - than when this system has several solutions.

## 6. Formulation of the problem

- A natural question is:
  - since in the general two-roots case, we cannot return *both* roots,
  - maybe we can return at least one of them?

## 7. What was known and what we plan to do

- It is known that no such algorithm is possible for *general* computable functions.
- This construction requires a computable function which is more complex than a polynomial.
- In this paper, we show that already for computable *polynomial* equations, it is impossible to compute even one of the roots.
- In this proof, we will use the polynomials:
  - with the smallest possible number of variables and
  - of the smallest possible degree.
- We also prove a similar result for optimization problems.

## 8. Definitions

- We want to formulate our result in precise terms.
- For this purpose, we need to recall the definition of a computable number.
- Crudely speaking, a real number is computable if it can be computed with an arbitrary accuracy.
- Formally, a real number  $x$  is called *computable* if there exists an algorithm that:
  - given an integer  $k$ ,
  - returns a rational number  $r$  for which  $|x - r_k| \leq 2^{-k}$ .
- By a *computable polynomial*, we mean a polynomial with computable coefficients.
- In line with our promise, we will prove this result for the case of the smallest possible number of variables: one.

## 9. First result

- *No algorithm is possible:*
  - *that is applicable to any computable polynomial function  $f(x)$  with exactly two roots,*
  - *and that returns one of these roots.*
- Our proof will use the known fact that no algorithm is possible for detecting:
  - whether a given constructive real number  $\alpha$  is non-negative
  - or whether it is non-positive.

- For every computable real number  $\alpha$ , we can form a polynomial

$$f_\alpha(x) = [(x - 1)^2 + \alpha] \cdot [(x + 1)^2 - \alpha].$$

- This polynomial is equal to 0 if one of the two factors is equal to 0.
- When  $\alpha = 0$ , this polynomial  $f_\alpha(x)$  has exactly two roots: 1 and  $-1$ .

## 10. First result (cont-d)

- When  $\alpha > 0$ , the first factor is positive, so  $f_\alpha(x) = 0$  if and only if  $(x + 1)^2 = \alpha$ , hence  $x + 1 = \pm\sqrt{\alpha}$ .
- So, for such  $\alpha$ , the polynomial  $f_\alpha(x)$  has exactly two roots:  $x = -1 \pm \sqrt{\alpha}$ .
- Similarly, when  $\alpha < 0$ , the polynomial  $f_\alpha(x)$  has exactly two roots:  $x = 1 \pm \sqrt{|\alpha|}$ .
- If we could compute one of roots, then:
  - by computing this root with enough accuracy and comparing it with 1,
  - we could tell whether this root is close to 1 or to  $-1$ .
- According to our description of the roots:
  - if this root is close to 1, then  $\alpha \leq 0$ ;
  - if this root is close to  $-1$ , then  $\alpha \leq 0$ .

## 11. First result (cont-d)

- We have mentioned that we cannot check whether  $\alpha \geq 0$  or  $\alpha \leq 0$ .
- We thus cannot return one of the roots.
- The proposition is proven.

## 12. Comment

- In the proof, we used a 4th degree polynomial.
- Let us give reasons why we cannot use a polynomial of a lower degree.
- Since we need polynomials with two roots, we must use polynomials of degree at least 2.
- If a quadratic polynomial has exactly 2 roots, then we can find these roots by using a standard formula.
- So these roots are easy to compute.
- For a cubic polynomial  $f(x)$ , the only way to have exactly two real roots is to have one double root.
- At this root, the derivative  $f'(x)$  is equal to 0.
- And the solution to the quadratic equation  $f'(x) = 0$  can be easily found.

### 13. Second result

- *No algorithm is possible:*
  - *which is applicable to any computable polynomial  $f(x)$  that attains its minimum at exactly two points,*
  - *and which returns one of these points.*
- *Proof:* it is sufficient to consider  $f_\alpha^2(x)$ , where  $f_\alpha(x)$  is the polynomial from the previous proof.
- This new polynomial is always non-negative, and it attains its minimum 0 if and only if  $f_\alpha(x) = 0$ .

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