## Designing, Understanding, and Analyzing Unconventional Computation: The Important Role of Logic and Constructive **Mathematics**

Grigori Mints
Stanford University, gmints@stanford.edu

Vladik Kreinovich University of Texas at El Paso, vladik@utep.edu



## 1. Main Problem: Reminder and Possible Approaches

- *Problem:* computations are often too slow.
- Traditional approaches:
  - design faster super-computers (hardware);
  - design faster algorithms.
- Limitations of the traditional approaches:
  - re hardware: we use the same physical processes as before;
  - re algorithms: we solve the same exact problem as before while often, data are imprecise.
- Possible new approaches:
  - hardware: use unconventional physical (and biological) processes;
  - algorithms: perform computations only up to accuracy that matches the input accuracy.



# 2. Our Main Claim: Use Logic (and Constructive Mathematics)

- Alternative approaches (reminder):
  - use unconventional physical (and biological) processes;
  - perform computations only up to accuracy that matches the input accuracy.
- Our claim: for these approaches to succeed, it is crucial to further develop:
  - ullet the corresponding tools of mathematical logic, and
  - ullet the related methods of  $constructive\ mathematics.$



## 3. Why Logic?

- Claim: logic is useful on all the stages of solving a problem:
  - we *specify* a problem;
  - $\bullet$  we design and implement an algorithm;
  - ullet we *verify* the corresponding program.
- Specifying a problem:
  - sometimes, the problem is to solve an equation;
  - in general, a proper formulation requires quantifiers etc. (stable control, etc.) i.e., logic.
- Designing an algorithm: logic programming transforms a logical specification into an algorithm.
- Program verification: logic helps in reasoning about programs; e.g., pre-condition implies post-condition.
- $\bullet$  Proof assistant programs (based on logic) help to prove.



# 4. 1st Approach: Computations with Limited Accuracy

- Objective: compute f(x) with accuracy  $\varepsilon > 0$ .
- Idea: compute x only with accuracy  $\delta > 0$  for which  $d(x, x') \leq \delta$  implies  $d(f(x), f(x')) \leq \varepsilon$ .
- Constructive math: a number is  $r_n$  s.t.  $d(r_n, x) \leq 2^{-n}$ ; a constructive function is a pair of:
  - ullet an algorithm  $f:X \to Y$  and
  - an algorithm  $\varepsilon \to \delta$ .
- Status: this is developed only for a few problems.
- Research Direction I.1:
  - $\bullet \ \ develop \ general \ constructive \ mathematics \ techniques,$
  - with a special emphasis on problems requiring intensive computations (e.g., large-scale PDE).

Why Logic? 1st Approach: . . . 1st Approach: Need . . . Interval Computations Proof Mining 2nd Approach: . . . Potential Use of Potential Use of . . . Explicit Use of . . . Unconventional . . . Title Page **>>** Page 5 of 22 Go Back Full Screen Close

Our Main Claim: Use . . .

## 5. 1st Approach: Need for Logic

- Decomposition: solutions to complex problems usually come from combining solutions to subproblems.
- Logic is needed for combination: e.g., stable robust control means it is stable  $\forall$  possible parameter values.
- ullet Constructive logic: we want to preserve constructions:
  - $\exists x P(x)$  should mean that we can construct such x;
  - $\forall x \exists y \ P(x,y) \leftrightarrow \exists \text{ algorithm } \varphi : X \to Y \text{ s.t. } P(x,\varphi(x)).$
- It was originated by Kolmorogov: e.g.,  $A \vee \neg A$  is false.
- Fact: constructive logic is used in constructive math.
- Research Direction I.2:
  - develop general constructive logic techniques,
  - with a special emphasis on problems requiring intensive computations.



## 6. Interval Computations

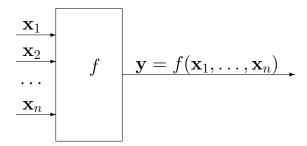
- Constructive math: algorithms work for all accuracies.
- In practice: we only need one given accuracy.
- Solution: interval computations (a.k.a. applied constructive mathematics).
- *Idea*: if we know an estimate  $\widetilde{x}$  w/accuracy  $\Delta$ , then

$$x \in \mathbf{x} = [\widetilde{x} - \Delta, \widetilde{x} + \Delta].$$

- Traditional approach: we also know probability distribution for  $\Delta x \stackrel{\text{def}}{=} \widetilde{x} x$  (usually Gaussian).
- Where it comes from: calibration using standard MI.
- *Problem:* calibration is not possible in:
  - fundamental science no better Measuring Instr. (MI);
  - manufacturing too expensive to calibrate.

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## 7. Interval Computations (cont-d)



- Given: algorithm  $y = f(x_1, ..., x_n)$  and  $\mathbf{x}_i = [\underline{x}_i, \overline{x}_i]$ .
- Compute: the corresponding range of y:

$$\mathbf{y} = [\underline{y}, \overline{y}] = \{ f(x_1, \dots, x_n) \mid x_1 \in [\underline{x}_1, \overline{x}_1], \dots, x_n \in [\underline{x}_n, \overline{x}_n] \}.$$

- Fact: this problem is NP-hard even for quadratic f.
- Challenge: find a good approximation  $Y \supseteq y$ .
- Applications: spaceflights, super-colliders, robotics, chemical engineering, nuclear safety, etc.
- Modal logic is efficiently used:  $\Box (y \in \mathbf{y}), \Diamond (x = \widetilde{x}).$

Our Main Claim: Use . . . Why Logic? 1st Approach: . . . 1st Approach: Need . . . Interval Computations Proof Mining 2nd Approach: . . . Potential Use of . . . Potential Use of . . . Explicit Use of . . . Unconventional . . . Title Page Page 8 of 22 Go Back Full Screen Close

#### 8. Proof Mining

- *Historically:* first existence proofs were *direct* (constr.).
- Currently: many proofs are indirect they prove  $\exists x \ P(x)$  without constructing such x.
- *Meta-results:* sometimes, we can extract ("mine") a constructive proof from a non-constructive one.
- Example (Kohlenbach): uniqueness  $\rightarrow$  computability.
- *Idea*: to find x s.t. f(x) = 0, compute  $\min_{B_{\varepsilon}(x_i)} |f(x)|$  w/increasing accuracy for  $x_i$  from  $\varepsilon$ -nets,  $\varepsilon = 2^{-1}, 2^{-2}, \dots$
- Research Direction I.4:
  - further develop proof mining,
  - with a special emphasis on its use to develop algorithms for realistic large-scale problems.



## 9. 2nd Approach: Unconventional Computations

- Main idea: use non-standard physical processes to speed up computations.
- Most well-known example: quantum computing:
  - search in an un-sorted array of size n in time  $\sqrt{n}$  (Grover);
  - factoring large integers in polynomial time (Shor).
- Limitation: the only provable speed-up is polynomial.
- Other schemes: can potentially lead to exponential speed-up.
- In other words: we can potentially solve NP-hard problems in polynomial time.
- Simplest example: acausal processes compute and send the result back in time.



## 10. Potential Use of Acausal Processes (cont-d)

- *Idea* (reminder): compute and send the result back in time.
- *Problem:* paradoxes of time travel e.g., killing your own grandfather before your father was conceived.
- Solution: some low probability event prevented the time traveller from this killing.
- Consequence: time travel (TT) can trigger events with probability  $p_0 \ll 1$ .
- Typical NP-hard problem: SAT given a propositional formula F(x), find  $x = (x_1, ..., x_n)$  s.t. F(x) holds.
- Usage (H. Moravec et al.): to solve SAT, generate n bits x and if  $\neg F(x)$ , launch TT.
- Why it works: for  $2^{-n} \gg p_0$ , TT is statistically improbable.



## 11. Potential Use of Curved Space-Time

- Parallelization a natural source of speed-up.
- Claim: in Euclidean space-time, parallelization only leads to a polynomial speed-up.
- Fact: the speed of all the physical processes is bounded by the speed of light c.
- Conclusion: in time T, we can only reach computational units at a distance  $\leq R = c \cdot T$ .
- The volume V(R) of this area (inside of the sphere of radius  $R = c \cdot T$ ) is proportional to  $R^3 \sim T^3$ .
- So, we can use  $\leq V/\Delta V \sim T^3$  computational elements.
- Interesting: in Lobachevsky space-time,  $V(R) \sim \exp(R)$ .
- Same is true for some more realistic space-time models.
- Hence, we can fit exponentially many processors and thus get an exponential speed-up.



## 12. Explicit Use of Kolmogorov Complexity

- ullet Fact: it's often difficult to describe biological processes.
- *Idea* (M. Gell-Mann): physical equations should include terms explicitly depending on complexity.
- ullet Natural formalization: Kolmogorov complexity

$$K(x) \stackrel{\text{def}}{=} \min\{\text{len}(p) : p \text{ generates } x\}.$$

- Conclusion: by observing physical and biological processes, we can measure the value K(x).
- Observation: K(x) is not algorithmically computable.
- Known results: ability to get non-computable values can speed up computations.
- Other schemes are based on:
  - quantum field theory (G. Kreisel),
  - $\bullet$  that every theory is approximate, etc.



## 13. Unconventional Computations (UC) and Constructive Mathematics

- Above UC schemes: use or propose a radically new physical process.
- Fact: some UC schemes were discovered by analyzing computability of known physical equations.
- Example (M. Pour-El et al.): even for wave equation, for some computable u(x, 0), u(x, T) is not computable.
- Desirable: extend the existing UC activity to the analysis of what computations can be sped up.
- Research Direction II.1. Use constructive mathematics to analyze:
  - how the use of physical processes (described by physically meaningful equations)
  - can speed up computations.



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