

# Random Interval Arithmetic is Closer to Common Sense: An Observation

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# 1. Commonsense Arithmetic

- We have a bridge whose weight we know with an accuracy of 1 ton.
- On this bridge, we place a car whose weight we know with an accuracy of 5 kg.
- The accuracy of the overall weight is still 1 ton.
- This is what an engineer or a physicist would say.
- *Related joke:*
  - in 2000, a dinosaur was 14,000,000 years old;
  - so, in 2005, it must be 14,000,005 years old.
- *What is desired:* if  $\Delta_a \gg \Delta_b$ , and
  - we add “uncertainty approximately  $\Delta_b$ ” to “uncertainty approximately  $\Delta_a$ ”,
  - we should get “uncertainty approximately  $\Delta_a$ ”.

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## 2. Traditional Interval Arithmetic Does not Have the Desired Property

- A natural way of dealing with approximately known values is *interval arithmetic*.
- The value  $\tilde{a}$  with an accuracy  $\Delta_a$  is interpreted as an interval  $[\tilde{a} - \Delta_a, \tilde{a} + \Delta_a]$ .
- *Specifics*:
  - we know  $\tilde{a}$  with uncertainty  $\Delta_a$ ;
  - we know  $\tilde{b}$  with uncertainty  $\Delta_b$ ;
  - then,  $\mathbf{a} = [\tilde{a} - \Delta_a, \tilde{a} + \Delta_a]$ ,  $\mathbf{b} = [\tilde{b} - \Delta_b, \tilde{b} + \Delta_b]$ , and
  - so, the set of possible values of  $c = a + b$  is an interval

$$\mathbf{c} = \mathbf{a} + \mathbf{b} = [(\tilde{a} + \tilde{b}) - (\Delta_a + \Delta_b), (\tilde{a} + \tilde{b}) + (\Delta_a + \Delta_b)].$$

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### 3. Traditional Interval Arithmetic Does not Have the Desired Property

- *Situation:*

- we know  $\tilde{a}$  with uncertainty  $\Delta_a$ ;
- we know  $\tilde{b}$  with uncertainty  $\Delta_b$ ;
- we conclude that

$$c \in [(\tilde{a} + \tilde{b}) - (\Delta_a + \Delta_b), (\tilde{a} + \tilde{b}) + (\Delta_a + \Delta_b)].$$

- *Interpretation:* we thus interpret this interval as

$$“\tilde{a} + \tilde{b} \text{ with uncertainty } \Delta_a + \Delta_b”.$$

- *Conclusion:* if we know  $a$  with uncertainty 1 ton, and we know  $b$  with uncertainty 5 kg, then the resulting uncertainty in  $a + b$  is 1.005 ton.
- *Problem:* how can we modify interval arithmetic?

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## 4. Interval Arithmetic: Origins

- *Objective:* analyze how:
  - the uncertainty in input data, and
  - the round-off imprecision of computer operationsaffect the results of the computations.
- *Traditional approach:* statistical techniques.
- *Problem:*
  - we must know the exact probability distributions of the input and round-off errors;
  - in practice, we *don't know* these distributions.
- *What we do know:* upper bounds on the errors – i.e., *intervals* that contain them.
- *e.g.:* space navigation under uncertainty (NASA, 1950s).
- *Interval arithmetic* was developed.

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## 5. Interval Arithmetic: Limitations

- *Problem:* producing the exact bounds on the inaccuracy of the output is often difficult (NP-hard).
- *Discussion:* the origin of interval techniques is in NASA-related problems that required high reliability.
- *Conclusion:* the emphasis in interval computations has always been on getting the *validated* results.
- Interval techniques produce estimates that are guaranteed to contain (enclose) the actual error.
- *Limitation:* it is often desirable,
  - in addition to guaranteed “overestimates”,
  - to produce a reasonable estimate of the size of the actual error,
  - an estimate that may be only valid with a certain probability.

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## 6. Interval Arithmetic: Main Idea

- *Main idea:* we follow computations step by step.
- *Specifics:* for each intermediate computation step  $z := x \odot y$ ,
  - once we have already computed the intervals  $\mathbf{x} = [\underline{x}, \bar{x}]$  and  $\mathbf{y} = [\underline{y}, \bar{y}]$  of possible values of  $x$  and  $y$ ,
  - we compute the interval for  $z$ .
- *Traditional interval arithmetic:* apply interval arithmetic operation to  $\mathbf{x}$  and  $\mathbf{y}$  corresponding to the worst case.
- *Example:* for addition,

$$\mathbf{z} = [\underline{x} + \underline{y}, \bar{x} + \bar{y}].$$

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## 7. Random Interval Arithmetic

- *New idea (Vignes et al.) – motivation:*
  - depending on the relative monotonicity of the  $x$  and  $y$  relative to inputs,
  - the intervals  $\mathbf{z}$  can change from the worst-case situation to the best-case situation.

- *Best case arithmetic:* (a.k.a. dual or inner): e.g., for addition,

$$\mathbf{z} = [\min(\underline{x} + \bar{y}, \bar{x} + \underline{y}), \max(\underline{x} + \bar{y}, \bar{x} + \underline{y})].$$

- *Reasonable assumption:*
  - monotonicity in the same direction and
  - monotonicity in different directions

are equally frequent.

- *Result:* on each step, we pick traditional or inner arithmetic with equal probability.

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## 8. Random Interval Arithmetic Has the Desired Property

- *Example:* addition  $\mathbf{c} = \mathbf{a} + \mathbf{b}$ .
- *Traditional arithmetic:* the half-width is:

$$\Delta_c^t = \Delta_a + \Delta_b.$$

- *Dual arithmetic:*  $\Delta_c^d = \max(\Delta_a, \Delta_b) - \min(\Delta_a, \Delta_b)$ .
- *Random interval arithmetic:* uses each operation with probability 50%.
- So, the average half-width of  $\mathbf{c}$  is  $\Delta_c^r = \frac{\Delta_c^t + \Delta_c^d}{2}$ .
- *Fact:*  $\Delta_c^t = \Delta_a + \Delta_b = \max(\Delta_a, \Delta_b) + \min(\Delta_a, \Delta_b)$ .
- *Conclusion:*  $\Delta_c^r = \max(\Delta_a, \Delta_b)$ .
- *Good news:* this is exactly the intuitive property that we have been trying to formalize.

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## 9. What If We Add $n$ Values?

- *Problem:*
  - we know each quantity  $a_i$  with an accuracy  $\Delta_i$ ;
  - what is the (expected value of) the accuracy in  $a = a_1 + \dots + a_n$ ?
- First, we add  $a_1 + a_2$ ; the resulting accuracy is  $\max(\Delta_1, \Delta_2)$ .
- To estimate the uncertainty of the next intermediate result  $(a_1 + a_2) + a_3$ , we take, as an estimate of the uncertainty in  $a_1 + a_2$ , the value  $\max(\Delta_1, \Delta_2)$ .
- Then, the average uncertainty in  $(a_1 + a_2) + a_3$  will be equal to

$$\max(\max(\Delta_1, \Delta_2), \Delta_3) = \max(\Delta_1, \Delta_2, \Delta_3).$$

- Similarly, we conclude that the average uncertainty in  $a_1 + \dots + a_n$  is equal to

$$\max(\Delta_1, \dots, \Delta_n).$$

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## 10. Computing $f(a_1, a_2)$

- When  $\Delta a_i \stackrel{\text{def}}{=} a_i - \tilde{a}_i \ll a_i$ , we can safely linearize the expression for  $f(a_1, a_2)$ :

$$f(a_1, a_2) = f(\tilde{a}_1 + \Delta a_1, \tilde{a}_2 + \Delta a_2) =$$

$$f(\tilde{a}_1, \tilde{a}_2) + \frac{\partial f}{\partial a_1} \cdot \Delta a_1 + \frac{\partial f}{\partial a_2} \cdot \Delta a_2.$$

- So, when  $\Delta a_i \in [\Delta_i, \Delta_i]$ , the worst-case half-width in  $a = f(a_1, a_2)$  is equal to

$$\Delta^t = \left| \frac{\partial f}{\partial a_1} \right| \cdot \Delta_1 + \left| \frac{\partial f}{\partial a_2} \right| \cdot \Delta_2.$$

- The result of applying dual interval arithmetic is

$$\Delta^d = \left| \left| \frac{\partial f}{\partial a_1} \right| \cdot \Delta_1 - \left| \frac{\partial f}{\partial a_2} \right| \cdot \Delta_2 \right|.$$

- Thus, the average half-width – corresponding to random interval arithmetic – is equal to

$$\Delta^r = \max \left( \left| \frac{\partial f}{\partial a_1} \right| \cdot \Delta_1, \left| \frac{\partial f}{\partial a_2} \right| \cdot \Delta_2 \right).$$

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## 11. Computing $f(a_1, \dots, a_n)$

- Similarly, for  $n > 2$  variables, we conclude that

$$\Delta^r = \max \left( \left| \frac{\partial f}{\partial a_1} \right| \cdot \Delta_1, \dots, \left| \frac{\partial f}{\partial a_n} \right| \cdot \Delta_n \right).$$

- In interval computations, we estimate the range of a function over a box  $[\underline{a}_1, \bar{a}_1] \times \dots \times [\underline{a}_n, \bar{a}_n]$ .
- If a box is not too narrow, the estimates are too wide.
- To improve the estimates, we:
  - bisect the box along one of the directions and
  - repeat the estimation for each of the two half-boxes.
- The optimal direction in a direction  $a_i$  in which the product  $\left| \frac{\partial f}{\partial a_i} \right| \cdot \Delta_i$  is the largest possible.
- The above value  $\Delta^r$  is exactly the value of this maximum.

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## 12. Relation with Fuzzy Logic

- *Our formula:*

$$\Delta_c = \max(\Delta_1, \Delta_b).$$

- *Fuzzy logic – objective:*

- we know:

- \* the degree of belief  $a = d(A)$  in a statement  $A$  and

- \* the degree of belief  $b = d(B)$  in a statement  $B$ ,

- we want to estimate the degree of belief  $c = d(C)$  in  $C \stackrel{\text{def}}{=} A \vee B$ .

- In the most widely used (and most practically successful version) of fuzzy logic,

$$c = \max(a, b).$$

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