

# Why $1/(1 + d)$ Is an Effective Distance-Based Similarity Measure: Two Explanations

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## 1. Need for similarity measures

- Many of our decisions are based on the idea of similarity:
  - if some decision was effective in similar situations,
  - then it makes sense to apply a similar decision here.
- Suppose that we know, for each  $i$  from 1 to  $n$ , that a decision  $d_i$  was successful in situation  $s_i$ .
- It can be a control decision, a medical decision, a financial decision, etc.
- Then, to select a decision  $d$  in a new situation  $s$ , we should use the following natural rules:
  - if  $s$  is similar to  $s_1$ , then  $d$  should be similar to  $d_1$ ;
  - if  $s$  is similar to  $s_2$ , then  $d$  should be similar to  $d_2$ ;
  - ...
  - if  $s$  is similar to  $s_n$ , then  $d$  should be similar to  $d_n$ .

## 2. Need for similarity measures (cont-d)

- These rules use an imprecise (“fuzzy”) natural-language word “similar”.
- Such natural-language words are ubiquitous.
- To transform such rules into a precise decision making strategy, Lotfi Zadeh invented fuzzy methodology.
- In this methodology, each imprecise property can be described by assigning:
  - to each possible object,
  - the degree to which, according to the expert, this object satisfies this property.
- In our case, we can ask the expert:
  - for each pair of situations (or pair of decisions)  $a$  and  $b$
  - to estimate to what extent the following statement is true: “ $a$  and  $b$  are similar”.

### 3. Need for similarity measures (cont-d)

- In a computer, “true” is usually represented as 1, and “false” as 0.
- So it is natural to represent an intermediate degree of confidence by a number between 0 and 1.
- This way:
  - to estimate the degree of similarity  $s(a, b)$  between objects  $a$  and  $b$ ,
  - we ask an expert to mark his/her degree of similarity between the two objects on a scale of 0 to 1.
- The value  $s(a, b) = 1$  means that the objects are perfectly similar, practically indistinguishable.
- The value  $s(a, b) = 0$  means that the objects are completely dissimilar, i.e., that they have nothing in common.
- Values strictly between 0 and 1 describe the cases when there is some similarity, but there is some dissimilarity as well.

#### 4. Need for similarity measures (cont-d)

- Sometimes, experts are not comfortable providing numerical estimates of their degree of similarity.
- They can only give us binary answers: similar or not similar.
- Then we can ask several ( $n$ ) experts this question.
- If  $m$  of them answer that the objects are similar, use the ratio  $\frac{m}{n}$  as the desired degree of similarity.

## 5. Need for metric-based similarity measures

- In many practical cases, we have a large number of possible objects and situations.
- In such cases, it is not feasible to ask the experts about each possible pair.
- What can we do?
- Often, we have a naturally defined metric  $d(a, b)$  on the class  $S$  of some objects.
- In other words, we have a function  $d : S \times S \rightarrow [0, \infty)$  that satisfies the usual properties:
  - $d(a, b) = 0$  if and only if  $a = b$ ,
  - $d(a, b) = d(b, a)$  for all  $a$  and  $b$ , and
  - $d(a, c) \leq d(a, b) + d(b, c)$  for all  $a, b$ , and  $c$ .
- This metric describes to what extent the two objects are dissimilar.

## 6. Need for metric-based similarity measures (cont-d)

- Thus, a natural idea is to estimate the desired degree of similarity  $s(a, b)$  between the two objects based on this metric, as:

$$s(a, b) = f(d(a, b)) \text{ for some function } f(d).$$

- Which function  $f(d)$  should we choose?

## 7. Natural properties of the transformation $f(d)$

- The degree of similarity must satisfy the following natural properties.
- The degree of similarity  $s(a, b)$  should attain its largest value  $s(a, b) = 1$  if the objects are identical (under given representation).
- So, if  $d(a, b) = 0$ ; thus, we must have  $f(0) = 1$ .
- The larger the distance between the objects, the smaller the similarity between them.
- Thus, the function  $f(d)$  should be strictly decreasing: if  $d < d'$ , then we should have  $f(d) > f(d')$ .
- In the limit, when the objects are as far away from each other as possible, the resulting degree of similarity should be close to 0.
- In other words, as  $d \rightarrow \infty$ , we should have  $f(d) \rightarrow 0$ .
- There are many functions  $f(d)$  that satisfy these three properties.
- Which one should we choose?

## 8. Empirical fact: an efficient transformation

- In many practical applications, the following function leads to reasonable similarity-based decisions

$$f(d) = \frac{1}{1 + d}.$$

- A natural question is: why this functions works well?
- In this talk, we provide two explanations of this empirical success.
- The fact that two different explanations lead to the same formula increases our confidence in both explanations.

## 9. Towards the first explanation

- When the degree of similarity comes from a poll of  $n$  experts, we only get  $n + 1$  possible degrees:  $0, \frac{1}{n}, \frac{2}{n}, \dots, \frac{n-1}{n}, 1$ .
- When  $n$  is small, these values provide a rather crude description of the actual degree of similarity.
- Thus, a natural way to increase the accuracy of the estimate is to ask more experts.
- This is similar to statistics, where we can estimate the probability of an event by taking the ratio  $m/n$  between:
  - the general number of situations  $n$  and
  - the number of cases  $m$  in which this event was observed.
- In statistics, the larger the sample size  $n$ , the more accurate this estimation of the probability.

## 10. Resulting problem

- To make our estimate more accurate, we ask the more knowledgeable experts.
- So, at first, we asked  $n$  top experts.
- Then, to increase the accuracy, we ask  $n'$  additional experts.
- These additional experts may be intimidated by the opinion of the top experts.
- This intimidation may be described in two ways.
- Additional experts may be unwilling to say anything: if top experts are disagreeing, who are we to voice our humble opinions?
- In this case, out of  $n + n'$  experts, we still have the same number  $m$  of experts who answer that the objects  $a$  and  $b$  are similar.
- Thus, instead of the original degree of similarity  $s = \frac{m}{n}$ , we have a new degree  $s' = \frac{m}{n + n'}$ .

## 11. Resulting problem (cont-d)

- One can easily see that the new degree  $s'$  can be obtained from the original degree by a transformation

$$s' = c_1 \cdot s, \text{ where } c_1 \stackrel{\text{def}}{=} \frac{n}{n + n'}.$$

- Alternatively, additional experts can simply side with the majority.
- We are looking for cases when there *is* a similarity – in this case, we can use this similarity to make a decision.
- So let us consider the case when originally, the majority of experts believed that the objects are similar.
- In this case, now we have  $m + n'$  experts who answer that the given objects  $a$  and  $b$  are similar.
- Thus, instead of the original degree of similarity  $s = \frac{m}{n}$ , we have a new degree  $s' = \frac{m + n'}{n + n'}.$

## 12. Resulting problem (cont-d)

- One can easily see that the new degree  $s'$  can be obtained from the original degree by a transformation  $s' = c_1 \cdot s + c_2$ , where  $c_2 \stackrel{\text{def}}{=} \frac{n'}{n + n'}$ .
- In both cases, we have linear transformations between different scales, i.e., linear functions  $s' = g(s)$ .

### 13. This is similar to measurements in general

- This possibility of a linear transformation between different scales is similar to the fact that in measurements:
  - we can select a different measuring unit, and
  - for some quantities like time or temperature, we can select a different starting point.
- When we use a measuring unit which is  $c_1$  times smaller, than all numerical values get multiplied by  $c_1$ :  $x \mapsto c_1 \cdot x$ .
- For example, when we replace meters with centimeters, then 1.7 m becomes 170 cm.
- When we use a starting point which is  $c_2$  units earlier than the original one, then this value  $c_2$  is added to all numerical values:  $x \mapsto x + c_2$ .
- If we change both the measuring unit and the starting point, then we get a general linear transformation  $c \mapsto c_1 \cdot x + c_2$ .

## 14. This is similar to measurements in general (cont-d)

- In measurements, we often also have nonlinear transformations.
- The energy of an earthquake can be measured either by its energy, or by the logarithm of its energy – which is the usual Richter scale.
- Similarly, the energy of a signal can be measured in the usual energy units, or in decibels, which is the logarithmic scale.
- In some applications, more complex transformations are used as well.
- Similarly to this, we can potentially envision non-linear transformation between different scales of degree of similarity.
- What form can these transformations have?

## 15. What are possible nonlinear transformations?

- Let us analyze what are reasonable transformations in general.
- First of all, all linear transformations are reasonable.
- If a transformation from one scale to another is reasonable, then an inverse transformation is also reasonable.
- If we have two reasonable transformations, then:
  - applying them one after another – i.e., performing a superposition of these transformations
  - should also lead to a reasonable transformation.
- Thus, the class of all reasonable transformations should be closed under taking the inverse and under taking the superposition.
- In mathematics, such classes are called *transformation groups*.
- Finally, our goal is to use this information in computer-aided decision making.

## 16. What are possible nonlinear transformations (cont-d)

- In each computer, we can only store finitely many values.
- So it makes sense to limit ourselves to classes of transformations which are determined by finitely many parameters.
- Such transformation groups are called *finite-dimensional*; so:
  - the question of which transformations are reasonable can be reformulated as
  - a question of what are the possible finite-dimensional transformation groups that contain all linear transformations.
- A general description of such groups was conjectured by Norbert Wiener, the father of Cybernetics.

## 17. What are possible nonlinear transformations (cont-d)

- This conjecture was proved in the 1960s.
- In particular, for functions of one variables, all the transformations from each such group must be fractionally linear:

$$g(x) = \frac{A \cdot x + B}{1 + C \cdot x}.$$

## 18. Let us apply this conclusion to our case

- Both the similarity measure  $s(a, b) = f(d(a, b))$  and the original metric  $d(a, b)$  describe the similarity between the two objects  $a$  and  $b$ .
- Thus, we can consider similarity and metric as representing the same quantity in two different scales.
- So, based on what we have concluded, the transformation  $f(d)$  between these two scales must be fractionally-linear:

$$f(d) = \frac{A \cdot d + B}{1 + C \cdot d} \text{ for some } A, B, \text{ and } C.$$

- To find the values of these three parameters, let us recall the above-mentioned properties of the function  $f(d)$ :
  - that  $f(0) = 1$ ,
  - that  $f(d) \rightarrow 0$  as  $d \rightarrow \infty$ , and
  - that  $f(d)$  is a decreasing function of  $d$ .

## 19. Let us apply this conclusion to our case (cont-d)

- Substituting  $d = 0$  into the general formula and equating the result to 1, we conclude that  $B = 1$ , so  $f(d) = \frac{A \cdot d + 1}{1 + C \cdot d}$ .
- For  $d \rightarrow \infty$ , this expression tends to  $\frac{A}{C}$ .
- Thus, the fact that this limit should be equal to 0 means that  $\frac{A}{C} = 0$ , i.e., that  $A = 0$ .
- Thus, the desired nonlinear transformation has the form

$$f(d) = \frac{1}{1 + C \cdot d}.$$

- The requirement that the function  $f(d)$  is decreasing leads to  $C > 0$ .

## 20. From “almost exactly” to “exactly”.

- This is almost exactly the desired formula.
- Let us take into account that the distance  $d(a, b)$  can also be described by using different measuring units:
  - if for distance, we select a measuring unit which is  $C$  times smaller than the original one,
  - then the new numerical values of the distance take the form

$$d' = C \cdot d.$$

- If we describe the distance in these new units, then the above formula takes exactly the desired form  $f(d') = \frac{1}{1 + d'}$ .
- Thus, we have indeed explained the emergence of the empirical formula – it is the only formula corresponding to natural requirements.

## 21. Main idea behind the second explanation

- In the first explanation, we focused on analyzing what is the actual dependence between the distance and the similarity.
- In this explanation, we ignored the fact that similarity usually comes from people marking a value on the interval  $[0, 1]$ .
- In reality, such markings are very uncertain.
- There is a well-known “seven plus minus two law” according to which, in particular:
  - when we do such types of markings,
  - we, in effect, only distinguish between 5 to 9 different values.
- Thus, the accuracy with which we mark the similarity value ranges:
  - from 11% (corresponding to 9 classes on the interval  $[0, 1]$ )
  - to 20% (corresponding to 5 classes on this interval).
- This inaccuracy can be easily observed.

## 22. Main idea behind the second explanation (cont-d)

- If we ask people to mark the same thing again, they may use somewhat different values (within this accuracy).
- With such imprecise values, it makes sense:
  - not to seek exact matching of the dependence  $s = f(d)$ ,
  - but rather to look for functions which are the fastest to compute.
- Indeed, as we have mentioned, the ultimate goal of assigning similarity values is to make decisions.
- Often, we need to make decision as soon as possible.
- So, the question becomes: of all the functions  $f(d)$  that satisfy the above three conditions, which ones are the fastest to compute?

## 23. Which functions are the fastest to compute?

- In a computer, the only directly hardware supported operations are arithmetic ones: addition, subtraction, multiplication, and division.
- Everything else is computed as a sequence of such arithmetic operations, for which the operands are:
  - either constants,
  - or the input values,
  - or the results of previous arithmetic operations.
- For example:
  - when we ask a computer to compute the values  $\exp(x)$ ,
  - what the computer will actually compute is the sum of the first few terms of the Taylor series for this functions:

$$\exp(x) \approx 1 + \frac{x}{1!} + \frac{x^2}{2!} + \dots + \frac{x^k}{k!}.$$

## 24. Which functions are the fastest to compute (cont-d)

- So, the computation time of each computation is, crudely speaking, proportional to the number of arithmetic operations.
- So, the fastest computations are the ones that use the smallest number of such arithmetic operations.

## 25. Computing $f(d)$ must include division

- Let us first explain that computing the function  $f(d)$  must include division.
- Indeed, if this computation only included addition, subtraction, and multiplication, then we would compute a polynomial.
- Polynomials do not tend to 0 as  $d \rightarrow \infty$ .
- Thus, at least one arithmetic operation must be division.

## 26. Can we have just one division?

- Can we just have one division? Not really.
- In this case, when we start with  $d$  and constants, the only things we can get by division are

$$\frac{C}{d}, \quad \frac{d}{C}, \quad \text{and} \quad \frac{d}{d} = 1.$$

- The first two expressions do not satisfy the property  $f(0) = 1$ .
- The third expression is not decreasing to 0 as  $d$  increases.
- Thus, we cannot have only one arithmetic operation, we must have at least one more arithmetic operation.

27. Which functions  $f(d)$  can be computed in two computational steps?

- The empirical expression requires two arithmetic operations:
  - first, we add 1 and  $d$ , and
  - then, we divide 1 by  $1 + d$ .
- So, this is clearly one of the fastest-to-compute functions  $f(d)$ .
- What other functions  $f(d)$  satisfying all three requirements we can compute in two arithmetic operations – one of which is division.

## 28. What if we perform division first

- If we perform division first, we get

$$\frac{C}{d} \text{ or } \frac{d}{C}.$$

- If we start with the first of these options, then
  - on the next step, as a second input to the second arithmetic operation,
  - we can have a constant or the original value  $d$ .
- Thus, we have the following options.

- If the second operation is addition or subtraction, we get

$$\frac{C}{d} + C' \text{ or } \frac{C}{d} \pm d.$$

- None of these expressions satisfies the condition  $f(0) = 1$ .
- If the second operation is multiplication, we get

$$\frac{C}{d} \cdot C' = \frac{C \cdot C'}{d} \text{ or } \frac{C}{d} \cdot d = C.$$

## 29. What if we perform division first (cont-d)

- Here, we do not get any new functions.
- If we second operation is division, then we get:

$$\frac{\frac{C}{d}}{C'} = \frac{C/C'}{d}, \quad \frac{C'}{\frac{C}{d}} = \frac{C'}{C} \cdot d,$$

$$\frac{\frac{C}{d}}{d} = \frac{C}{d^2}, \quad \frac{d}{\frac{C}{d}} = \frac{1}{C} \cdot d^2.$$

- The first and third expressions do not satisfy the requirement that  $f(0) = 1$ .
- The second and fourth are polynomials – and we have already mentioned that the transformation  $f(d)$  cannot be a polynomial.

### 30. What if we first compute $d/C$

- What if first compute  $d/C$ ?
- On the next step, as a second input to the second arithmetic operation, we can have a constant or the original value  $d$ .
- If the second operation is addition, subtraction, or multiplication, we get a polynomial.
- We have already mentioned that the function  $f(d)$  cannot be a polynomial.
- This, the only possible option is when the second arithmetic operation is division.

### 31. What if we first compute $d/C$ (cont-d)

- In this case, we get the following options:

$$\frac{\frac{1}{C} \cdot d}{C'} = \frac{C}{C'} \cdot d, \quad \frac{C'}{\frac{1}{C} \cdot d} = \frac{C \cdot C'}{d},$$

$$\frac{\frac{1}{C} \cdot d}{d} = \frac{1}{C}, \quad \frac{d}{\frac{1}{C} \cdot d} = C.$$

- In the first case we get a polynomial.
- In the second case, we do not satisfy the requirement that  $f(0) = 1$ .
- In the third and fourth cases, we get constants.
- So, none of these options lead to functions  $f(d)$  that satisfy all three requirements.

## 32. What if division is the second arithmetic operation

- The cases when division is the first arithmetic operation do not lead to a function  $f(d)$  that satisfies all three conditions.
- So, we need to perform division only as a second arithmetic operation.
- In this case, the first arithmetic operation is addition, subtraction, or multiplication.
- Thus, as a result of the first arithmetic operation, we get  $d+C$ ,  $C-d$ , or  $C \cdot d$ .
- When the first arithmetic operation results in  $d+C$ , we have  $d$ , constants, and  $d+C$ .
- Thus, we have the following division options.
- The first option is  $\frac{C'}{d+C}$ .

### 33. What if division is the 2nd arithmetic operation (cont-d)

- The requirement that  $f(0) = 1$  leads to  $C' = C$ , so this expression is equal to  $\frac{C}{d+C} = \frac{1}{1+C^{-1} \cdot d}$ .
- This is exactly the expression that, as we have shown, is equivalent to the desired one after an appropriate re-scaling of distance.
- The second option is  $\frac{d}{d+C}$  which does not satisfy the condition  $f(0) = 1$ .
- The third option is  $\frac{d+C}{C'} = \frac{1}{C'} \cdot d + \frac{C}{C'}$ .
- This is a polynomial, so it cannot satisfy all three conditions.
- The fourth option is  $\frac{d+C}{d} = 1 + \frac{C}{d}$ .
- This option does not satisfy the condition  $f(0) = 1$ .

### 34. What if division is the 2nd arithmetic operation (cont-d)

- When the first arithmetic operation is subtraction, the conclusions are similar.
- When first operation results in  $C \cdot d$ , we have  $d$ , constants, and  $C \cdot d$ .
- Thus, we have the following division options.
- The first option is  $\frac{C'}{C \cdot d} = \frac{C''}{d}$ , where  $C'' \stackrel{\text{def}}{=} \frac{C'}{C}$ .
- So, in this case, we do not get a new function.
- The second option is  $\frac{d}{C \cdot d} = \frac{1}{C}$ , a constant function which is not decreasing.
- The third option is  $\frac{C \cdot d}{C'} = \frac{C}{C'} \cdot d$ , a polynomial.
- The fourth option is  $\frac{C \cdot d}{d} = C$ , a constant.

## 35. Summarizing

- We have considered all possible options; so
  - out of all functions  $f(d)$  that satisfy all three requirements,
  - the only functions that can be computed the fastest – in two arithmetic steps – are the functions of type  $1/(1 + C \cdot d)$ .
- We showed that these functions are, in effect, equivalent to the desired formula  $1/(1 + d)$ .
- Thus, we get the second explanation of the effectiveness of the empirical formula – that this function is the fastest to compute.

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