

# Algebraic Approach to Data Processing: Techniques and Applications

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## 1. Introduction

- How do we gain knowledge about the world?
- For example, how did we learn that if we drop an object, it will fall with the acceleration of  $9.81 \text{ m/sec}^2$ ?
- Well, the scientists dropped an object once, and observed this fall.
- Then they moved to a different location and repeated the same experiment – and got the exact same result.
- Then they turned by an angle – and also got the same result.
- After several such experiments, they concluded that:
  - the result of this experiment
  - does not change if we move to a different location or turn by some angle.
- In other words, they concluded that this process is *invariant* with respect to shifts and rotations.

## 2. Introduction (cont-d)

- In other cases, other transformations are appropriate.
- For example:
  - the whole idea of a wind tunnel – in which smaller-size airplane models used to be tested
  - is that the corresponding processes do not change if we re-scale the objects.
- In electrodynamics, all interactions remain the same if we replace all positive charges with negative ones, and vice versa.
- According to Special Relativity theory, processes do not change if everything starts moving with a constant speed, etc.
- In all these cases, we have some transformations with respect to which processes are invariant.
- In mathematics, studying the classes of such transformation is classified as part of *algebra*.

### 3. Introduction (cont-d)

- From this viewpoint, algebraic approach to designing (and optimizing) data processing algorithms is a very natural idea.
- In this thesis, we describe applications of this idea:
  - to various aspects of algorithmics,
  - to dynamic systems,
  - to physics,
  - to engineering,
  - to medicine,
  - to economics,
  - to social sciences,
  - to education, and
  - to mathematics.
- In this talk, I will give some examples of these applications.

Part I

# Why Squashing Functions in Multi-Layer Neural Networks

## 4. A Short Introduction

- In their successful applications, deep neural networks use a non-linear transformation  $s(z) = \max(0, z)$ .
- It is called a *rectified linear* activation function.
- Sometimes, more general transformations – called *squashing functions* – lead to even better results.
- In this talk, we provide a theoretical explanation for this empirical fact.
- To provide this explanation, let us first briefly recall:
  - why we need machine learning in the first place,
  - what are deep neural networks, and
  - what activation functions these neural networks use.

## 5. Machine Learning Is Needed

- For some simple systems, we know the equations that describe the system's dynamics.
- These equations may be approximate, but they are often good enough.
- With more complex systems (such as systems of systems), this is often no longer the case.
- Even when we have a good approximate model for each subsystem, the corresponding inaccuracies add up.
- So, the resulting model of the whole system is too inaccurate to be useful.
- We also need to use the records of the actual system's behavior when making predictions.
- Using the previous behavior to predict the future is called *machine learning*.

## 6. Deep Learning

- The most efficient machine learning technique is *deep learning*: the use of multi-layer neural networks.
- In general, on a layer of a neural network, we transform signals  $x_1, \dots, x_n$  into a new signal  $y = s \left( \sum_{i=1}^n w_i \cdot x_i + w_0 \right)$ .
- The coefficient  $w_i$  (called *weights*) are to be determined during training.
- $s(z)$  is a non-linear function called *activation function*.
- Most multi-layer neural networks use  $s(z) = \max(z, 0)$  known as *rectified linear* function.



## 7. Shall We Go Beyond Rectified Linear?

- Preliminary analysis shows that for some applications:
  - it is more advantageous to use different activation functions for different neurons;
  - specifically, this was shown for a special family of *squashing* activation functions

$$S_{a,\lambda}^{(\beta)}(z) = \frac{1}{\lambda \cdot \beta} \cdot \ln \frac{1 + \exp(\beta \cdot z - (a - \lambda/2))}{1 + \exp(\beta \cdot z - (a + \lambda/2))};$$

- this family contains rectified linear neurons as a particular case.
- We explain their empirical success of squashing functions by showing that:
  - their formulas
  - follow from reasonably natural symmetries.

## 8. How This Talk Is Structured

- First, we recall the main ideas of symmetries and invariance.
- Then, we recall how these ideas can be used to explain the efficiency of the sigmoid activation function

$$s_0(z) = \frac{1}{1 + \exp(-z)}.$$

- This function is used in the traditional 3-layer neural networks.
- Finally, we use this information to explain the efficiency of squashing activation functions.

## 9. Which Transformations Are Natural?

- From the mathematical viewpoint, we can apply any non-linear transformation.
- However, some of these transformations are purely mathematical, with no clear physical interpretation.
- Other transformation are *natural* in the sense that they have physical meaning.
- What are natural transformations?

## 10. Numerical Values Change When We Change a Measuring Unit And/Or Starting Point

- In data processing, we deal with numerical values of different physical quantities.
- Computers just treat these values as numbers.
- However, from the physical viewpoint, the numerical values are not absolute; they change:
  - if we change the measuring unit and/or
  - the starting point for measuring the corresponding quantity.
- The corresponding changes in numerical values are clearly physically meaningful, i.e., natural.
- For example, we can measure a person's height in meters or in centimeters.

## 11. Numerical Values Change (cont-d)

- The same height of 1.7 m, when described in centimeters, becomes 170 cm.
- In general, if we replace the original measuring unit with a new unit which is  $\lambda$  times smaller, then:
  - instead of the original numerical value  $x$ ,
  - we get a new numerical value  $\lambda \cdot x$  – while the actual quantity remains the same.
- Such a transformation  $x \rightarrow \lambda \cdot x$  is known as *scaling*.
- For some quantities, e.g., for time or temperature, the numerical value also depends on the starting point.
- For example, we can measure the time from the moment when the talk started.
- Alternatively, we can use the usual calendar time, in which Year 0 is the starting point.

## 12. Numerical Values Change (cont-d)

- In general, if we replace the original starting point with the new one which is  $x_0$  units earlier, than:
  - each original numerical value  $x$
  - is replaced by a new numerical value  $x + x_0$ .
- Such a transformation  $x \rightarrow x + x_0$  is known as *shift*.
- In general, if we change both the measuring unit and the starting point, we get a linear transformation:

$$x \rightarrow \lambda \cdot x + x_0.$$

- A usual example of such a transformation is a transition from Celsius to Fahrenheit temperature scales:

$$t_F = 1.8 \cdot t_C + 32.$$

## 13. Invariance

- Changing the measuring unit and/or starting point:
  - changes the numerical values but
  - does not change the actual quantity.
- It is therefore reasonable to require that physical equations do not change if we simply:
  - change the measuring unit and/or
  - change the starting point.
- Of course, to preserve the physical equations:
  - if we change the measuring unit and/or starting point for one quantity,
  - we may need to change the measuring units and/or starting points for other quantities as well.
- For example, there is a well-known relation  $d = v \cdot t$  between distance  $d$ , velocity  $v$ , and time  $t$ .

## 14. Invariance (cont-d)

- If we change the measuring units for measuring distance and time:
  - this formula remains valid –
  - but only if we accordingly change the units for velocity.
- For example:
  - if we replace kilometers with meters and hours with seconds,
  - then, to preserve this formula, we also need to change the unit for velocity from km/h to m/sec.



## 15. Natural Transformations Beyond Linear Ones

- In some cases, the relation between different scales is non-linear.
- For example, we can measure the earthquake energy:
  - in Joules (i.e., in the usual scale) or
  - in a logarithmic (Richter) scale.
- Which nonlinear transformation are natural?
- First, as we have argued, all linear transformations are natural.
- Second:
  - if we have a natural transformation  $f(x)$  from scale  $A$  to another  $B$ ,
  - then the inverse transformation  $f^{-1}(x)$  from scale  $B$  to scale  $A$  should also be natural.

## 16. Natural Transformations (cont-d)

- Third:
  - if  $f(x)$  and  $g(x)$  are natural scale transformation,
  - then we can apply first  $g(x)$  to get  $y = g(x)$  and then  $f$  to get  $f(y) = f(g(x))$ .
- Thus, the composition  $f(g(x))$  of two natural transformations should also be natural.
- The class of transformations that satisfies the 2nd and 3rd properties is called a *transformation group*.
- We also need to take into account that in a computer:
  - at any given moment of time,
  - we can only store the values of finitely many parameters.
- Thus, the transformations should be determined by a finite number of parameters.

## 17. Natural Transformations (cont-d)

- The smallest number of parameters needed to describe a family is known as the *dimension* of this family.
- E.g., that we need 3 coordinates to describe any point in space means that the physical space is 3-dimensional.
- In these terms, the transformation group  $T$  must be finite-dimensional.

## 18. Let Us Describe All Natural Transformations

- Interestingly, the above requirements uniquely determine the class of all possible natural transformation.
- This result can be traced back to Norbert Wiener, the father of cybernetics.
- In his seminal book *Cybernetics*, he noticed that:
  - when we approach an object from afar,
  - our perception of this object goes through several distinct phases.
- First, we see a blob; this means that:
  - at a large distance,
  - we cannot distinguish between images obtained each other by all possible continuous transformations.
- This phase corresponds to the group of all possible continuous transformations.

## 19. All Natural Transformations (cont-d)

- As we get closer, we start distinguishing angular parts from smooth parts, but still cannot compare sizes.
- This corresponds to the group of all projective transformations.
- After that, we become able to detect parallel lines.
- This corresponds to the group of all transformations that preserve parallel lines.
- These are linear (= affine) transformations.
- When we get even closer, we become able to detect the shapes, sizes, etc.

## 20. All Natural Transformations (cont-d)

- Wiener argued that there are no other transformation groups – since:
  - if there were other transformation groups,
  - after billions of years of evolution, we would use them.
- In precise terms, he conjectured that:
  - the only finite-dimensional transformation group that contain all linear transformations
  - is the groups of all projective transformations.
- This conjecture was later proven.
- For transformations of the real line, projective transformations are simply fractional-linear transformations

$$f(x) = \frac{a \cdot x + b}{c \cdot x + d}.$$

- So, natural transformations are fractional-linear ones.

## 21. Traditional Neural Networks (NN)

- Let us recall why traditional neural networks appeared in the first place.
- The main reason, in our opinion, was that computers were too slow.
- A natural way to speed up computations is to make several processors work in parallel.
- Then, each processor only handles a simple task, not requiring too much computation time.
- For processing data, the simplest possible functions to compute are linear functions.

## 22. Traditional Neural Networks (cont-d)

- However, we cannot only use linear functions – because then:
  - no matter how many linear transformations we apply one after another,
  - we will only get linear functions, and many real-life dependencies are nonlinear.
- So, we need to supplement linear computations with some nonlinear ones.
- In general, the fewer inputs, the faster the computations.
- Thus, the fastest to compute are functions with one input, i.e., functions of one variable.



## 23. Traditional Neural Networks (cont-d)

- So, we end up with a parallel computational device that has:
  - linear processing units (L) and
  - nonlinear processing units (NL) that compute functions of one variable.
- First, the input signals come to a layer of such devices; we will call such a layer a *d-layer*; d for *device*.
- Then, the results of this d-layer go to another d-layer, etc.
- The fewer d-layers we have, the faster the computations.

## 24. How Many d-Layers Do We Need?

- It can be proven that:
  - 1-d-layer schemes (L or NL) are not sufficient to approximate any possible dependence, and
  - 2-d-layer schemes (L-NL, linear layer followed by non-linear layer, or NL-L) are also not enough.
- Thus, we need at least 3-d-layer networks – and 3-d-layer networks can be proven to be sufficient.
- In a 3-d-layer network:
  - we cannot have two linear layers or two nonlinear d-layers following each other,
  - that would be equivalent to having one d-layer since, e.g., a composition of two L functions is also L.
- So, our only options are L-NL-L and NL-L-NL.

## 25. How Many d-Layers Do We Need (cont-d)

- Since linear transformations are faster to compute, the fastest scheme is L-NL-L.
- In this scheme:
  - first, each neuron  $k$  in the L d-layer combines the inputs into a linear combination

$$z_k = \sum_{i=1}^n w_{ki} \cdot x_i + w_{k0};$$

- then, in the next d-layer, each such signal is transformed into  $y_k = s_k(z_k)$  for some non-linear f-n;
  - finally, in the last linear d-layer, we form a linear combination of the values  $y_k$ :  $y = \sum_{k=1}^K W_k \cdot y_k + W_0$ .

## 26. How Many d-Layers Do We Need (cont-d)

- The resulting transformation takes the form

$$y = \sum_{k=1}^K W_k \cdot s_k \left( \sum_{i=1}^n w_{ki} \cdot x_i + w_{k0} \right) + W_0.$$

- Usually, we use the same function  $s(z)$  for all transformations.
- This is indeed the usual formula of the traditional neural network.

## 27. Traditional NN Mostly Used Sigmoid

- Originally, the sigmoid function was selected because:
  - it provides a reasonable approximation to
  - how biological neurons process their inputs.
- Several other nonlinear activation functions have been tried.
- However, in most cases, the sigmoid  $s_0(z)$  leads to the best approximation results.
- A partial explanation for this empirical success is that:
  - neural networks using sigmoid activation function  $s_0(z)$  have proven to be universal approximators;
  - i.e., the corresponding neural networks can approximate any continuous function.
- However, many other non-linear activation functions have the same universal approximation property.

## 28. So, Why Sigmoid?

- We have mentioned that the values of physical quantities change when we:
  - change the starting point,
  - i.e., shift all the data points by the same constant  $x_0$ .
- At first glance, it may seem that this does not apply to neural data processing, since usually:
  - before we apply a neural network,
  - we *normalize* the data, i.e., transform all the input values into the some fixed interval (e.g.,  $[0, 1]$ ).
- This normalization is based on all the values of the corresponding quantity that have been observed so far.
- The smallest of these values corresponds to 0 and the largest to 1.

## 29. Why Sigmoid (cont-d)

- However, as we will show, shift still makes sense even for the normalized data.
- Indeed, in real life, signals come with noise, in particular, with background noise.
- Often, a significant part of this noise is a constant which is added to all the measured signals.
- This constant noise component is, in general, different for different situations.
- We can try to get rid of this constant noise component by subtracting the corresponding constant.
- So, we replace:
  - each original numerical value  $x_i$
  - with a corrected value  $x_i - n_i$ .

## 30. Why Sigmoid (cont-d)

- After this correction, instead of the original value  $z_k$ , we get a corrected value

$$z'_k = \sum_{i=1}^n w_{ki} \cdot (x_i - n_i) + w_{k0} = z_k - h'_k.$$

- Here, we denoted  $h'_k \stackrel{\text{def}}{=} \sum_{i=1}^n w_{ki} \cdot n_i$ .
- The trouble is that we do not know the exact values of these constants  $n_i$ .
- So, depending on our estimates, we may subtract different values  $n_i$  and thus, different values  $h'_k$ :
  - if we change from one value  $h'_k$  to another one  $h''_k$ ,
  - then the resulting value of  $z_k$  is shifted by the difference  $h_k \stackrel{\text{def}}{=} h'_k - h''_k$ :  $z'_k = z_k + h_k$ .



## 31. Why Sigmoid (cont-d)

- This is exactly the same formula as for the shift corresponding to the change in the starting point.
- Since we do not know what shift is the best, all shifts within a certain range are equally possible.
- It is therefore reasonable to require that the formula  $y = s(z)$  for the nonlinear activation function:
  - should work for all possible shifts,
  - i.e., this formula should be, in this sense, *shift-invariant*.
- In other words:
  - if we start with the formula  $y = s(z)$  and we shift from  $z$  to  $z' = z + h$ ,
  - then we should have the same relation  $y' = s(z')$  for an appropriately transformed  $y' = f(y)$ .

## 32. Why Sigmoid (cont-d)

- For different shifts  $h$ , we will have, in general, different natural transformations  $f(y)$ .
- We have mentioned that all natural transformations  $f(y)$  are fractionally linear.
- Thus, for each  $h$ ,  $y' = s(z+h)$  should be fractional-linear in  $y = s(z)$ :

$$s(z+h) = \frac{a(h) \cdot s(z) + b(h)}{c(h) \cdot s(z) + d(h)}.$$

- It turns out that this implies the sigmoid  $s_0(z)$ .

### 33. Why Sigmoid: Derivation

- For  $h = 0$ , we should have  $s(z + h) = s(z)$ , thus, we should have  $d(0) \neq 0$ .
- It is reasonable to require that the function  $d(h)$  is continuous.
- In this case,  $d(h)$  is different from 0 for all small  $h$ .
- Then, we can divide both numerator and denominator of the above formula by  $d(h)$  and get a simpler formula:

$$s(z + h) = \frac{A(h) \cdot s(z) + B(h)}{C(h) \cdot s(z) + 1}, \text{ where } A(h) = a(h)/d(h), \dots$$

- For  $h = 0$ , we have  $s(z+h) = s(z)$ , so  $A(h) = 1$  and  $B(h) = C(h) = 0$ .
- It is also reasonable to require that the activation function  $s(z)$  be defined and smooth for all  $z$ .

### 34. Why Sigmoid: Derivation (cont-d)

- Indeed, on each interval, every continuous function:
  - can be approximated, with any desired accuracy,
  - by a smooth one – even by a polynomial.
- So, from the practical viewpoint, it is sufficient to only consider smooth activation functions.
- Multiplying both sides of the above formula by the denominator, we get:

$$s(z + h) = A(h) \cdot s(z) + B(h) - C(h) \cdot s(z + h) \cdot s(z).$$

- Let us take three different values  $z_i$ .
- Then, for each  $h$ , we get 3 linear equations for three unknown  $A(h)$ ,  $B(h)$ , and  $C(h)$ :

$$s(z_i + h) = A(h) \cdot s(z_i) + B(h) - C(h) \cdot s(z_i + h) \cdot s(z_i), \quad i = 1, 2, 3.$$

### 35. Why Sigmoid: Derivation (cont-d)

- Due to Cramer's rule, the solution to this system is:
  - a ratio of two determinants,
  - i.e., a ratio of two polynomials of the coefficients.
- Thus,  $A(h)$ ,  $B(h)$ , and  $C(h)$  are smooth functions of the values  $s(z_i + h)$ .
- Since the function  $s(z)$  is smooth, we conclude that all three functions  $A(h)$ ,  $B(h)$ , and  $C(h)$  are also smooth.
- Thus, we can differentiate both sides of the above equation by  $h$  and get

$$s'(z + h) = \frac{N(h)}{(C(h) \cdot s(z) + 1)^2}, \text{ where}$$

$$N(h) \stackrel{\text{def}}{=} (A'(h) \cdot s(z) + B'(h)) \cdot (C(h) \cdot s(z) + 1) - (A(h) \cdot s(z) + B(h)) \cdot (C'(h) \cdot s(z)).$$

### 36. Why Sigmoid: Derivation (cont-d)

- In particular, for  $h = 0$ , taking into account that  $A(h) = 1$  and  $B(h) = C(h) = 0$ , we conclude that

$$s'(z) = a_0 + a_1 \cdot s(z) + a_2 \cdot (s(z))^2, \text{ where } a_0 = B'(0), \dots$$

- So,  $\frac{ds}{dz} = a_0 + a_1 \cdot s + a_2 \cdot s^2$  and

$$\frac{ds}{a_0 + a_1 \cdot s + a_2 \cdot s^2} = dz.$$

- We can now integrate both sides of this formula and get an explicit expression of  $z(s)$ .
- Based on this expression, we can find the explicit formula for the dependence of  $s$  on  $z$ .

## 37. Why Sigmoid: Derivation (cont-d)

- The only non-linear dependencies  $s(z)$  are:
  - the sigmoid (plus some linear transformations before and after) and
  - the sigmoid's limit case  $\exp(z)$ .
- So, the sigmoid  $s_0(z)$  is the only shift-invariant activation function.
- This explains its efficiency in traditional neural networks.

## 38. We Need Multi-Layer Neural Networks

- The problem with traditional neural networks is that they waste a lot of bits:
  - for  $K$  neurons,
  - any of  $K!$  permutations results in exactly the same function.
- To decrease this duplication, we need to decrease the number of neurons  $K$  in each layer.
- So, instead of placing all nonlinear neurons in one layer, we place them in several consecutive layers.
- This is one of the main idea behind deep learning.



## 39. Which Activation Function Should We Use

- In the first nonlinear d-layer, we make sure that:
  - a shift in the input – corresponding to a different estimate of the constant noise component,
  - does not change the processing formula,
  - i.e., that results  $s(z+c)$  and  $s(z)$  can be obtained from each other by an appropriate transformation.
- We already know that this idea leads to the sigmoid function  $s_0(z)$ .
- This logic doesn't work if we try to find out what activation function we should use in the *next* NL d-layer.
- Indeed, the input to the 2nd NL d-layer is the output of the 1st NL d-layer.
- This input is *no longer* shift-invariant.

## 40. Which Activation Function (cont-d)

- This input is invariant with respect to some more *complex* (fractional linear) transformations.
- We know what to do when the input is shift-invariant.
- So a natural idea is to perform some *additional* transformation that will make the results shift-invariant.
- If we do that, then:
  - we will again be able to apply the sigmoid activation function  $s_0(z)$ ,
  - then again the additional transformation, etc.
- These additional transformations should transform generic fractional-linear operations into shift.

## 41. Which Activation Function (cont-d)

- Thus, the inverse of such a transformation should transform shifts into fractional-linear operations.
- But this is exactly what we analyzed earlier – transformations that transform shifts into fractional-linear.
- We already know the formulas  $s(z)$  for these transformations.
- In general, they are formed as follows:
  - first, we apply some linear transformation to the input  $z$ , resulting in a linear combination

$$Z = p \cdot z + q;$$

- then, we compute  $Y = \exp(Z)$ ; and
  - finally, we apply some fractional-linear transformation to the resulting value  $Y$ , getting  $y$ .

## 42. Which Activation Function (cont-d)

- So, to get the inverse transformation, we need to reverse all three steps, starting with the last one:
  - first, we apply a fractional-linear transformation to  $y$ , getting  $Y$ ;
  - then, we compute  $Z = \ln(Y)$ ; and
  - finally, we apply a linear transformation to  $Z$ , resulting in  $z$ .

### 43. This Leads Exactly to Squashing Functions

- What happens if we:
  - first apply a sigmoid-type transformation moving us from shifts to fractional-linear operations,
  - and then an inverse-type transformation?
- The last step of the sigmoid transformation and the first step of the inverse are fractional-linear.
- The composition of fractional-linear transformations is fractional-linear.
- So, we can combine these 2 steps into a single step.

## 44. This Leads to Squashing Functions (cont-d)

- Thus, the resulting combined activation function can thus be described as follows:
  - first, we apply some linear transformation  $L_1$  to the input  $z$ , resulting in a linear combination

$$Z = L_1(z) = p \cdot z + q;$$

- then, we compute  $E = \exp(Z) = \exp(L_1(z))$ ;
  - then, we apply a fractional-linear transformation  $F$  to  $E = \exp(Z)$ , getting  $T = F(E) = F(\exp(L_1(z)))$ ;
  - then, we compute  $Y = \ln(T) = \ln(F(\exp(L_1(z))))$ ;
  - and finally, we apply a linear transformation  $L_2$  to  $Y$ , resulting in the final value

$$y = s(z) = L_2(Y) = L_2(\ln(F(\exp(L_1(z)))).$$

## 45. This Leads to Squashing Functions (cont-d)

- One can check that these are exactly squashing function!
- Thus, squashing functions can indeed be naturally explained by the invariance requirements.

## 46. Example

- Let us provide a family of squashing functions that tend to the rectified linear activation function  $\max(z, 0)$ .
- For this purpose, let us take:
  - $L_1(z) = k \cdot z$ , with  $k > 0$ , so that
$$E = \exp(L_1(z)) = \exp(k \cdot z);$$
  - $F(E) = 1 + E$ , so that  $T = F(E) = \exp(k \cdot z) + 1$  and  $Y = \ln(T) = \ln(\exp(k \cdot z) + 1)$ ; and
  - $L_2(Y) = \frac{1}{k} \cdot Y$ , so that the resulting activation function takes the form  $s(z) = \frac{1}{k} \cdot \ln(\exp(k \cdot z) + 1)$ .
- Let us show that this expression tends to the rectified linear activation function when  $k \rightarrow \infty$ .
- When  $z < 0$ , then  $\exp(k \cdot z) \rightarrow 0$ , so  $\exp(k \cdot z) + 1 \rightarrow 1$ ,  $\ln(\exp(k \cdot z) + 1) \rightarrow 0$  and so  $s(z) \rightarrow 0$ .



## 47. Example (cont-d)

- On the other hand, when  $z > 0$ , then

$$\exp(k \cdot z) + 1 = \exp(k \cdot z) \cdot (1 + \exp(-k \cdot z)).$$

- Thus,  $\ln(\exp(k \cdot z) + 1) = k \cdot z + \ln(1 + \exp(-k \cdot z))$  and

$$s(z) = \frac{1}{k} \cdot \ln(\exp(k \cdot z) + 1) = z + \frac{1}{k} \cdot \ln(1 + \exp(-k \cdot z)).$$

- When  $k \rightarrow \infty$ , we have  $\exp(-k \cdot z) \rightarrow 0$ , hence

$$1 + \exp(-k \cdot z) \rightarrow 1, \quad \ln(1 + \exp(-k \cdot z)) \rightarrow 0.$$

- So  $\frac{1}{k} \cdot \ln(1 + \exp(-k \cdot z)) \rightarrow 0$  and indeed  $s(z) \rightarrow z$ .

Part II

# Natural Invariance Explains Empirical Success of Specific Membership Functions, Hedge Operations, and Negation Operations

## 48. Fuzzy Techniques: A Brief Reminder

- In many applications, we have knowledge formulated:
  - in terms of imprecise (“fuzzy”) terms from natural language,
  - like “small”, “somewhat small”, etc.
- To translate this knowledge into computer-understandable form, Lotfi Zadeh proposes *fuzzy techniques*.
- According to these techniques, each imprecise property like “small” can be described by assigning:
  - to each value  $x$  of the corresponding quantity,
  - a degree  $\mu(x)$  to which, according to the expert, this property is true.

## 49. Fuzzy Techniques (cont-d)

- These degrees are usually selected from the interval  $[0, 1]$ , so that:
  - 1 corresponds to full confidence,
  - 0 to complete lack of confidence, and
  - values between 0 and 1 describe intermediate degrees of confidence.
- The resulting function  $\mu(x)$  is known as a *membership function*.
- In practice, we can only ask finitely many questions to the expert.
- So we only elicit a few values  $\mu(x_1)$ ,  $\mu(x_2)$ , etc.
- Based on these values, we need to estimate the values  $\mu(x)$  for all other values  $x$ .

## 50. Fuzzy Techniques (cont-d)

- For this purpose, usually:
  - we select a family of membership functions – e.g., triangular, trapezoidal, etc. – and
  - we select a function from this family which best fits the known values.
- For terms like “somewhat small”, “very small”, the situation is more complicated.
- We can add different “hedges” like “somewhat”, “very”, etc., to each property.
- As a result, we get a large number of possible terms.

## 51. Fuzzy Techniques (cont-d)

- It is not realistically possible to ask the expert about each such term; instead:
  - practitioners estimate the degree to which, e.g., “somewhat small” is true
  - based on the degree to which “small” is true.
- In other words, with each linguistic hedge, we associate a function  $h$  from  $[0, 1]$  to  $[0, 1]$  that:
  - transforms the degree to which a property is true
  - into an estimate for the degree to which the hedged property is true.

## 52. Fuzzy Techniques (cont-d)

- Similarly to the membership functions:
  - we can elicit a few values  $h(x_i)$  of the hedge operation from the experts, and
  - then we extrapolate and/or interpolate to get all the other values of  $h(x)$ .
- Usually, a family of hedge operations is pre-selected.
- Then we select a specific operation from this family which best fits the elicited values  $h(x_i)$ .

## 53. Fuzzy Techniques (cont-d)

- Similarly:
  - instead of asking experts for their degrees of confidence in statements like “not small”,
  - we estimate these degrees based on their degrees of confidence in the positive statements.
- The corresponding operation  $n(x)$  is known as the *negation operation*.



## 54. Need to Select Proper Membership Functions, Hedge Operations, And Negation Operations

- Fuzzy techniques have been successfully applied to many application areas.
- However, this does not necessarily mean that every time we try to use fuzzy techniques, we get a success.
- The success (or not) often depends on which membership functions etc. we select:
  - for some selections, we get good results (e.g., good control),
  - for other selections, the results are not so good.
- There is a lot of empirical data about which selections work better.
- In this talk, we provide a general explanation for several of these empirically best selections.

## 55. Need to Select Proper Functions (cont-d)

- This explanation is based on the natural concepts of invariance.
- For symmetric membership functions that describe properties like “small”,
  - for which  $\mu(x) = \mu(-x)$  and the degree  $\mu(|x|)$  decreases with  $|x|$ ,
  - in many practical situations, the most empirically successful are so-called *distending* functions:

$$\mu(x) = \frac{1}{1 + a \cdot |x|^b}.$$

- Among hedge and negation operations, often, the most efficient are fractional linear functions:

$$h(x) = \frac{a + b \cdot x}{1 + c \cdot x}.$$

## 56. Re-Scaling

- The variable  $x$  describes the value of some physical quantity, such a distance, height, etc.
- When we process these values, we deal with numbers.
- Numbers depend on the selection of the measuring unit:
  - if we replace the original measuring unit with a new one which is  $\lambda$  times smaller,
  - then all the numerical values will be multiplied by  $\lambda$ :  $x \rightarrow X = \lambda \cdot x$ .
- For example, 2 meters become  $2 \cdot 100 = 200$  cm.
- This transformation from one measuring scale to another is known as *re-scaling*.

## 57. Scale-Invariance: Idea

- In many physical situations, the choice of a measuring unit is rather arbitrary.
- In such situations, all the formulas remain the same no matter what unit we use.
- For example, the formula  $y = x^2$  for the area of the square with side  $x$  remains valid:
  - if we replace the unit for measuring sides from meters with centimeters,
  - of course, we then need to appropriately change the unit for  $y$ , from  $\text{m}^2$  to  $\text{cm}^2$ .

## 58. Scale-Invariance (cont-d)

- In general, invariance of the formula  $y = f(x)$  means that:
  - for each re-scaling  $x \rightarrow X = \lambda \cdot x$ , there exists an appropriate re-scaling  $y \rightarrow Y$
  - for which the same formula  $Y = f(X)$  will be true for the re-scaled variables  $X$  and  $Y$ .

## 59. Let Us Apply This Idea to the Membership Function

- It is reasonable to require that:
  - the selection of the best membership functions
  - should also not depend on the choice of the unit for measuring the corresponding quantity  $x$ .
- So, it is reasonable to require that for each  $\lambda > 0$ :
  - there should exist some reasonable transformation  $y \rightarrow Y = T(y)$  of the degree of confidence
  - for which  $y = \mu(x)$  implies  $Y = \mu(X)$ .

## 60. So, What Are Reasonable Transformations of the Degree of Confidence?

- One way to measure the degree of confidence is to have a poll:
  - ask  $N$  experts how many of them believe that a given value  $x$  is, e.g., small,
  - count the number  $M$  of those who believe in this, and
  - take the ratio  $M/N$  as the desired degree  $y = \mu(x)$ .
- As usual with polls, the more people we ask, the more adequately we describe the general opinion.
- So, to get a more accurate estimate for  $\mu(x)$ , it is reasonable to ask more people.
- When we have a limited number of people to ask, it is reasonable to ask top experts in the field.

## 61. Reasonable Transformations (cont-d)

- When we start asking more people:
  - we are thus adding people who are less experienced,
  - and who may therefore be somewhat intimidated by the opinions of the top experts.
- This intimidation can be expressed in different ways.
- Some new people may be too shy to express their own opinion, so they will keep quiet; as a result:
  - if we add  $A$  people to the original  $N$ , we still have the same number  $M$  of people voting “yes”,
  - and the new ratio is  $Y = \frac{M}{N + A}$ .
- Here,  $Y = a \cdot y$ , where  $a \stackrel{\text{def}}{=} \frac{N}{N + A}$ .
- Some new people will be too shy to think on their own and will vote with the majority.



## 62. Reasonable Transformations (cont-d)

- So when  $M > N/2$ , we will have  $Y = \frac{M + A}{N + A}$ .
- Since  $M = y \cdot N$ , we will have  $Y = \frac{y \cdot N + A}{N + A} = a \cdot y + b$ , where  $a$  is the same as before and  $b = \frac{A}{N + A}$ .
- We may also have a situation in which:
  - a certain proportion  $c$  of the new people keep quiet while
  - the others vote with the majority.
- In this case, we have  $Y = \frac{M + (1 - c) \cdot A}{N + A} = a \cdot y + b$ , where  $a = (1 - c) \cdot \frac{A}{N + A}$ .

## 63. Reasonable Transformations (cont-d)

- In all these cases, we have a linear transformation

$$Y = a \cdot y + b.$$

- So, it seems reasonable to identify reasonable transformations with linear ones.
- We will call the corresponding scale-invariance L-scale-invariance (L for Linear).

## 64. What Membership Functions We Consider

- We consider symmetric properties, for which

$$\mu(-x) = \mu(x).$$

- So it is sufficient to consider only positive values  $x$ .
- We consider properties like “small” for which  $\mu(x)$  decreases with  $x$  and  $\lim_{x \rightarrow \infty} \mu(x) = 0$ .
- We will call such membership functions s-membership functions (s for small).
- We say that an s-membership function  $\mu(x)$  is *L-scale-invariant* if:
  - for every  $\lambda > 0$ , there exist values  $a(\lambda)$  and  $b(\lambda)$
  - for which  $y = \mu(x)$  implies  $Y = \mu(X)$ , where

$$X = \lambda \cdot x \text{ and } Y = a(\lambda) \cdot y + b(\lambda).$$

## 65. What Membership Functions (cont-d)

- Unfortunately, this does not solve our problem:
- **Proposition 1.** *The only  $L$ -scale-invariant  $s$ -membership functions are constant functions  $\mu(x) = \text{const.}$*
- What does this result mean?
- We considered two possible types of reasonable transformations of the degrees of confidence.
- They both turned out to be linear.
- This was not enough.
- So probably there are other reasonable transformations of degrees of confidence.
- How can we describe such transformations?

## 66. What Membership Functions (cont-d)

- Clearly, if we have a reasonable transformation, then its inverse is also reasonable.
- Also, a composition of two reasonable transformations should be a reasonable transformation too.
- So, in mathematical terms, reasonable transformations should form a *group*.
- This group should be finite-dimensional, i.e.:
  - different transformations should be uniquely determined
  - by a finite number of parameters – since in the computer, we can store only finitely many parameters.

## 67. What Membership Functions (cont-d)

- We also know that linear transformations are reasonable; so, we are looking for:
  - a finite-dimensional group of transformations from real numbers to real numbers
  - that contains all linear transformations.
- It is known that all such transformations are piece-wise linear:  $\mu \rightarrow \frac{a \cdot \mu + b}{1 + c \cdot \mu}$ .
- Thus, we arrive at the following definitions.

## 68. Definitions and the Main Result

- We say that an s-membership function  $\mu(x)$  is *scale-invariant* if:
  - for every  $\lambda > 0$ , there exist  $a(\lambda)$ ,  $b(\lambda)$ , and  $c(\lambda)$
  - for which  $y = \mu(x)$  implies  $Y = \mu(X)$ , where

$$X = \lambda \cdot x \text{ and } Y = \frac{a(\lambda) \cdot y + b(\lambda)}{1 + c(\lambda) \cdot y}.$$

- **Proposition 2.** *The only scale-invariant s-membership functions are distending membership functions.*
- This result explains the empirical success of distending functions.

## 69. Which Hedge Operations and Negation Operations Should We Select

- We would like hedging and negation operations  $y = h(x)$  to be also invariant, i.e., that:
  - for each natural transformation  $X = T(x)$ , there should be a transformation  $Y = S(y)$
  - for which  $y = h(x)$  implies  $Y = h(X)$ .
- Now we know what are natural transformations of membership degrees – they are fractional-linear functions.
- Let us call this h-scale-invariance.
- **Proposition 3.** *The only h-scale-invariant functions are fractionally linear ones.*
- This result explains the empirical success of fractional-linear hedge operations and negation operations.



## 70. Proof of Proposition 1

- We will prove this result by contradiction.
- Let us assume that the function  $\mu(x)$  is not a constant, and let us derive a contradiction.
- Let us substitute the expressions for  $X$ ,  $Y$ , and  $y = \mu(x)$  into the formula  $Y = \mu(X)$ .
- Then, we conclude that for every  $x$  and for every  $\lambda$ , we have  $\mu(\lambda \cdot x) = a(\lambda) \cdot \mu(x) + b(\lambda)$ .
- It is known that monotonic functions are almost everywhere differentiable; due to the above formula:
  - if a function  $\mu(x)$  is differentiable at  $x = x_0$ ,
  - it is also differentiable at any point of the type  $\lambda \cdot x_0$  for every  $\lambda > 0$ ,
  - and thus, that it is differentiable for all  $x > 0$ .

## 71. Proof of Proposition 1 (cont-d)

- Since the function  $\mu(x)$  is not constant, there exist values  $x_1 \neq x_2$  for which  $\mu(x_1) \neq \mu(x_2)$ .
- For these values, the above formula has the form

$$\mu(\lambda \cdot x_1) = a(\lambda) \cdot \mu(x_1) + b(\lambda); \quad \mu(\lambda \cdot x_2) = a(\lambda) \cdot \mu(x_2) + b(\lambda).$$

- Subtracting the two equations, we get

$$\mu(\lambda \cdot x_1) - \mu(\lambda \cdot x_2) = a(\lambda) \cdot (\mu(x_1) - \mu(x_2)), \text{ thus}$$

$$a(\lambda) = \frac{\mu(\lambda \cdot x_1) - \mu(\lambda \cdot x_2)}{\mu(x_1) - \mu(x_2)}.$$

- Since the function  $\mu(x)$  is differentiable, we can conclude that the function  $a(\lambda)$  is also differentiable.
- Thus, the function  $b(\lambda) = \mu(\lambda \cdot x) - a(\lambda) \cdot \mu(x)$  is differentiable too.
- So, all three functions  $\mu(x)$ ,  $a(\lambda)$ , and  $b(\lambda)$  are differentiable.

## 72. Proof of Proposition 1 (cont-d)

- So, we can differentiate both sides of the equality

$$\mu(\lambda \cdot x) = a(\lambda) \cdot \mu(x) + b(\lambda) \text{ with respect to } \lambda.$$

- If we substitute  $\lambda = 1$ , we get  $x \cdot \mu'(x) = A \cdot \mu(x) + B$ , where we denoted  $A \stackrel{\text{def}}{=} a'(1)$ ,  $B \stackrel{\text{def}}{=} b'(1)$ .
- Here,  $\mu'(x)$ , as usual, indicates the derivative.
- Thus,  $x \cdot \frac{d\mu}{dx} = A \cdot \mu + B$ .
- We cannot have  $A = 0$  and  $B = 0$ , since then  $\mu'(x) = 0$  and  $\mu(x)$  would be a constant.
- Thus, in general, the expression  $A \cdot \mu + B$  is not 0, so

$$\frac{d\mu}{A \cdot \mu + B} = \frac{dx}{x}.$$

- If  $A = 0$ , then integration leads to  $\frac{1}{B} \cdot \mu(x) = \ln(x) + c$ , where  $c_0$  is the integration constant.

### 73. Proof of Proposition 1 (cont-d)

- Thus,  $\mu(x) = B \cdot \ln(x) + B \cdot c_0$ .
- This expression has negative values for some  $x$ , while all the values  $\mu(x)$  are in the interval  $[0, 1]$ .
- So, this case is impossible.
- If  $A \neq 0$ , then we have  $d(A \cdot \mu + B) = A \cdot d\mu$ , hence

$$\frac{d(A \cdot \mu + B)}{A \cdot \mu + B} = A \cdot \frac{dx}{x}.$$

- Integration leads to  $\ln(A \cdot \mu(x) + B) = A \cdot \ln(x) + c_0$ .
- By applying  $\exp(z)$  to both sides, we get  $A \cdot \mu(x) + B = \exp(c_0) \cdot x^A$ , i.e.,  $\mu(x) = A^{-1} \cdot \exp(c_0) \cdot x^A - B/A$ .
- This expression tends to infinity either for  $x \rightarrow \infty$  (if  $A > 0$ ) or for  $x \rightarrow 0$  (if  $A < 0$ ).
- In both cases, we get a contradiction with our assumption that  $\mu(x)$  is within the interval  $[0, 1]$ . Q.E.D.

## 74. Proof of Proposition 2

- Let us substitute the expressions for  $X$ ,  $Y$ , and  $y = \mu(x)$  into the formula  $Y = \mu(X)$ .
- Then, we conclude that for every  $x$  and for every  $\lambda$ :

$$\mu(\lambda \cdot x) = \frac{a(\lambda) \cdot \mu(x) + b(\lambda)}{1 + c(\lambda) \cdot \mu(x)}.$$

- Similarly to the previous proof, we can conclude that the function  $\mu(x)$  is differentiable for all  $x > 0$ .
- Multiplying both sides of the above equality by the denominator, we conclude that

$$\mu(\lambda \cdot x) + c(\lambda) \cdot \mu(x) \cdot \mu(\lambda \cdot x) = a(\lambda) \cdot \mu(x) + b(\lambda).$$

- So, for three different values  $x_i$ , we have the following three equations:

$$\mu(\lambda \cdot x_i) + c(\lambda) \cdot \mu(x_i) \cdot \mu(\lambda \cdot x_i) = a(\lambda) \cdot \mu(x_i) + b(\lambda), \quad i = 1, 2, 3.$$

## 75. Proof of Proposition 2 (cont-d)

- We thus have a system of three linear equations for three unknowns  $a(\lambda)$ ,  $b(\lambda)$ , and  $c(\lambda)$ .
- By Cramer's rule:
  - the solution to such a system
  - is a rational (hence differentiable) function of the coefficients and the right-hand sides.
- So, since  $\mu(x)$  is differentiable, we can conclude that  $a(\lambda)$ ,  $b(\lambda)$ , and  $c(\lambda)$  are differentiable.
- All the functions  $\mu(x)$ ,  $a(\lambda)$ ,  $b(\lambda)$ , and  $c(\lambda)$  are differentiable.
- So, we can differentiate both sides of the above formula with respect to  $\lambda$ .
- Let us substitute  $\lambda = 1$  and take into account that for  $\lambda = 1$ , we have  $a(1) = 1$  and  $b(1) = c(1) = 0$ .

## 76. Proof of Proposition 2 (cont-d)

- Then, we get  $x \cdot \frac{d\mu}{dx} = A \cdot \mu + B - C \cdot \mu^2$ , where  $A$  and  $B$  are the same as in the previous proof and  $C \stackrel{\text{def}}{=} c'(1)$ .
- For  $x \rightarrow \infty$ , we have  $\mu(x) \rightarrow 0$ , so  $\mu'(x) \rightarrow 0$ , and thus  $B = 0$  and  $x \cdot \frac{d\mu}{dx} = A \cdot \mu - C \cdot \mu^2$ .
- So,  $\frac{d\mu}{B \cdot \mu - C \cdot \mu^2} = \frac{dx}{x}$ .
- As we have shown in the previous proof, we cannot have  $C = 0$ , so  $C \neq 0$ .
- One can easily see that

$$\frac{1}{\mu - \frac{B}{C}} - \frac{1}{\mu} = \frac{\frac{B}{C}}{\mu \cdot \left(\mu - \frac{B}{C}\right)} = \frac{-B}{B \cdot \mu - C \cdot \mu^2}.$$

## 77. Proof of Proposition 2 (cont-d)

- Thus, by multiplying the equality  $\frac{d\mu}{B \cdot \mu - C \cdot \mu^2} = \frac{dx}{x}$  by  $-B$ , we get:  
$$\frac{d\mu}{\mu - \frac{B}{C}} - \frac{d\mu}{\mu} = -B \cdot \frac{dx}{x}.$$

- Integrating both sides, we get

$$\ln\left(\mu(x) - \frac{B}{C}\right) - \ln(\mu) = -B \cdot \ln(x) + c_0.$$

- By applying  $\exp(z)$  to both sides, we get

$$\frac{\mu(x) - \frac{B}{C}}{\mu(x)} = C_0 \cdot x^{-B}. \text{ so } 1 - \frac{B/C}{\mu} = C_0 \cdot x^{-B}.$$

- Hence  $\frac{B/C}{\mu} = 1 - C_0 \cdot x^{-B}$  and  $\mu(x) = \frac{B/C}{1 - C_0 \cdot x^{-B}}.$

- From the condition that  $\mu(0) = 1$ , we conclude that  $B < 0$  and  $B/C = 1.$



## 78. Proof of Proposition 2 (cont-d)

- From  $\mu(x) \leq 1$ , we conclude that  $C_0 < 0$ .
- So, we get the desired formula  $\mu(x) = \frac{1}{1 + |C_0| \cdot x^{|B|}}$ .
- The proposition is proven.

## 79. Proof of Proposition 3

- For constant functions the statement is trivial.
- Therefore, it is sufficient to prove for non-constant functions  $h(x)$ .
- Similarly to the previous proof, we can prove that the function  $h(x)$  is differentiable.
- Let  $x \in D$ , and let  $\lambda$  and  $x_0$  from an open neighborhood of 1 and 0 respectively be such that

$$\lambda \cdot x \in D \text{ and } x + x_0 \in D.$$

- Since the function  $h(x)$  is h-scale-invariant, there exist fractional-linear transformations for which

$$h(x + x_0) = \frac{a(x_0) \cdot h(x) + b(x_0)}{1 + c(x_0) \cdot h(x)} \text{ and}$$

$$h(\lambda \cdot x) = \frac{d(\lambda) \cdot h(x) + e(\lambda)}{1 + f(\lambda) \cdot h(x)}.$$

## 80. Proof of Proposition 3 (cont-d)

- Similarly to the previous proof, we can prove that the functions  $a(x_0)$ ,  $\dots$ , are differentiable.
- So, we can differentiate the  $\lambda$ -formula with respect to  $\lambda$  and take  $\lambda = 1$ , then we get:

$$x \cdot h' = D \cdot h + E - F \cdot h^2.$$

- Similarly, differentiating the  $h_0$ -formula with respect to  $x_0$  and taking  $x_0 = 0$ , we get:

$$h' = A \cdot h + B - C \cdot h^2.$$

- Let us consider two cases:  $C \neq 0$  and  $C = 0$ .
- Let us first consider the case when  $C \neq 0$ .
- By completing the square, we get  $h' = A \cdot h + B - C \cdot h^2 = \hat{A} - C \cdot (h - h_0)^2$  for some  $\hat{A}$  and  $h_0$ , i.e.,

$$h' = \hat{A} - C \cdot H^2, \text{ where } H \stackrel{\text{def}}{=} h - h_0.$$

## 81. Proof of Proposition 3 (cont-d)

- Substituting  $h = H + h_0$  into the right-hand side, we conclude that

$$x \cdot h' = \widehat{D} \cdot H + \widehat{E} - F \cdot H^2 \text{ for some } \widehat{D} \text{ and } \widehat{E}.$$

- Dividing the two equations, we get

$$x = \frac{\widehat{D} \cdot H + \widehat{E} - F \cdot H^2}{\widehat{A} - C \cdot H^2}, \text{ so } \frac{dx}{dH} =$$

$$\frac{(\widehat{D} - 2F \cdot H)(\widehat{A} - C \cdot H^2) - (\widehat{D} \cdot H + \widehat{E} - F \cdot H^2)(-2C \cdot H)}{(\widehat{A} - C \cdot H^2)^2} =$$

$$\frac{\widehat{A} \cdot \widehat{D} - 2(\widehat{A} \cdot F - C \cdot \widehat{E}) \cdot H + C \cdot \widehat{D} \cdot H^2}{(\widehat{A} - C \cdot H^2)^2}.$$

- On the other hand,

$$\frac{dx}{dH} = \frac{1}{\frac{dH}{dx}} = \frac{1}{\widehat{A} - C \cdot H^2}.$$

## 82. Proof of Proposition 3 (cont-d)

- The right-hand sides of these two formulas must be equal, so for all  $H$ , we have

$$\widehat{A} \cdot \widehat{D} - 2(\widehat{A} \cdot F - C \cdot \widehat{E}) \cdot H + C \cdot \widehat{D} \cdot H^2 = \widehat{A} - C \cdot H^2.$$

- Since the two polynomials of  $H$  are equal, the coefficients at 1,  $H$ , and  $H^2$  must coincide.
- Comparing the coefficients at  $H^2$ , we get  $C \cdot \widehat{D} = -C$ .
- Since  $C \neq 0$ , we conclude that  $\widehat{D} = -1$ .
- Comparing the coefficients at 1, we get  $\widehat{A} \cdot \widehat{D} = \widehat{A}$ , i.e.,  $-\widehat{A} = \widehat{A}$  and thus  $\widehat{A} = 0$ .
- Comparing the coefficients at  $H$  and taking into account that  $\widehat{A} = 0$ , we get  $0 = \widehat{A} \cdot F - C \cdot \widehat{E} = -C \cdot \widehat{E}$ .
- Since  $C \neq 0$ , this implies  $\widehat{E} = 0$ .

### 83. Proof of Proposition 3 (cont-d)

- So, the above formula for  $x$  takes the form

$$x = \frac{\widehat{D} \cdot H - F \cdot H^2}{-C \cdot H^2} = \frac{\widehat{D} - F \cdot H}{-C \cdot H}.$$

- Thus  $x$  is a fractional linear function of  $H$ .
- Hence  $H$  (and therefore  $h = H + h_0$ ) is also a fractional linear function of  $x$ .
- Let us now consider the case when  $C = 0$ .
- Then,  $h' = A \cdot h + B$  and  $x \cdot h' = D \cdot h + E - F \cdot h^2$ , so:

$$x = \frac{x \cdot h'}{h'} = \frac{D \cdot h + E - F \cdot h^2}{A \cdot h + B}.$$

- If  $F = 0$ , then  $x$  is a fractional linear function of  $h(x)$  and hence,  $h$  is also a fractional-linear function of  $x$ .
- So, it is sufficient to consider the case when  $F \neq 0$ .

## 84. Proof of Proposition 3 (cont-d)

- In this case, by completing the square, we can find constants  $\widehat{D}$ ,  $h_0$ , and  $\widehat{B}$  for which, for  $H = h - h_0$ :

$$\begin{aligned}x \cdot h' &= D \cdot h + E - F \cdot h^2 = \widetilde{D} - F \cdot H^2 \text{ and} \\h' &= A \cdot h + B = A \cdot H + \widehat{B}.\end{aligned}$$

- Dividing the first equation by the second one, we have

$$\begin{aligned}x &= \frac{\widetilde{D} - F \cdot H^2}{A \cdot H + \widehat{B}}, \text{ thus} \\ \frac{dx}{dH} &= \frac{(-2F \cdot H) \cdot (A \cdot H + \widehat{B}) - (\widehat{D} - F \cdot H^2) \cdot A}{(A \cdot H + \widehat{B})^2} \\ &= \frac{-A \cdot \widehat{D} - 2\widehat{B} \cdot F \cdot H - A \cdot F \cdot H^2}{(A \cdot H + \widehat{B})^2}.\end{aligned}$$

- On the other hand,  $\frac{dx}{dH} = -\frac{1}{\frac{dH}{dx}} = \frac{1}{A \cdot H + \widehat{B}}.$

## 85. Proof of Proposition 3 (cont-d)

- By equating the two expressions for the derivative and multiplying both sides by  $(A \cdot H + \widehat{B})^2$ , we get:

$$-A \cdot \widehat{D} - 2\widehat{B} \cdot F \cdot H - A \cdot F \cdot H^2 = A \cdot H + \widehat{B}.$$

- Thus  $A \cdot F = 0$ ,  $A = -2\widehat{B} \cdot F$ , and  $-A \cdot \widehat{D} = \widehat{B}$ .
- If  $A = 0$ , then we have  $\widehat{B} = 0$ , so  $h' = 0$  and  $h$  is a constant.
- However, we consider the case when the function  $h(x)$  is not a constant.
- Thus,  $A \neq 0$ , hence  $F = 0$ , and the above formula describes  $x$  as a fractional-linear function of  $H$ .
- Both for  $C \neq 0$  and  $C = 0$ ,  $x$  is fractionally linear in  $H$  (hence in  $h$ ).
- Since the inverse of a fractional linear is fractional linear, the function  $h(x)$  is also fractional linear. Q.E.D.



Part III

# Why Ellipsoids in Mechanical Analysis of Wood Structures

## 86. Formulation of the Problem

- Many constructions are made of wood.
- Wood is one of the oldest materials used in construction.
- During the past millennia, people have developed a lot of skills for working with wood.
- However, in spite of this experience, wood remains one of the most difficult materials to handle.
- The main reason for this difficulty is that:
  - in contrast to many other construction materials which are mostly homogeneous and isotropic,
  - wood is highly inhomogeneous and anisotropic.

## 87. Formulation of the Problem (cont-d)

- At each point in the wooden beam:
  - both the average values and fluctuations of the local mechanical properties
  - depend on whether the direction is longitudinal, radial or tangential with respect to the grain.
- In designing wooden constructions, it is important:
  - to properly describe and to properly take into account
  - this inhomogeneity and anisotropy.
- How can we describe local fluctuations of mechanical characteristics?
- These fluctuations are caused by many different relatively small factors.

## 88. Formulation of the Problem (cont-d)

- It is known that the distribution of the joint effect of a large number of small factors is close to Gaussian.
- This follows from the Central Limit Theorem, according to which:
  - this distribution tends to Gaussian
  - when the number of factors increases.
- To describe a Gaussian distribution, it is sufficient to describe its first and second moments.
- For a general random field  $f(x)$ , this means that we need to describe:
  - its mean values  $E[f(x)]$  (where  $E[\cdot]$  denotes the expected value) and
  - its covariances  $E[f(x) \cdot f(y)]$ .

## 89. Formulation of the Problem (cont-d)

- For fluctuations, the mean is 0, so we only need to describe covariances.
- In statistics, it is often convenient:
  - instead of explicitly describing covariances,
  - to describe the standard deviations and correlations:

$$\sigma[f(x)] \stackrel{\text{def}}{=} \sqrt{E[(f(x))^2]}; \quad \rho(x, y) \stackrel{\text{def}}{=} \frac{E[f(x) \cdot f(y)]}{\sigma[f(x)] \cdot \sigma[f(y)]}.$$

- Then, covariances can be reconstructed as

$$E[f(x) \cdot f(y)] = \sigma[f(x)] \cdot \sigma[f(y)] \cdot \rho(x, y).$$

- An interesting property of the corresponding correlation functions was recently empirically found.

## 90. Formulation of the Problem (cont-d)

- This property is about:
  - *iso-correlation* surfaces corresponding to each spatial location  $x$ ,
  - i.e., surfaces formed by all the points  $y$  for which the correlation  $\rho(x, y)$  is equal to a constant  $\rho_0$ .
- Empirical analysis shows that:
  - for each point  $x$ ,
  - the corresponding surfaces are well approximated by concentric homothetic ellipsoids.
- This property helps narrow down possible functions  $\rho(x, y)$  when we analyze mechanical properties of wood.
- Thus, it has a potential to make mechanical analysis of wooden structures more efficient.

## 91. Formulation of the Problem (cont-d)

- The problem is that so far, this property was purely empirical, it had no theoretical justification.
- Thus, engineers were reluctant to use it.
- It is known that sometimes:
  - empirical properties found under some conditions
  - do not work well when conditions change.
- We want to make this property more reliable and thus, more practically useful.
- It is therefore desirable to come up with a theoretical explanation.
- In this talk, we provide a desired theoretical explanation for this empirical fact.

## 92. Our Explanation: Main Idea

- We show that there exists the smallest dimension  $d$  for which:
  - it is possible to have an affine-invariant optimality criterion
  - on the space of all such  $d$ -dimensional classes.
- We also show that for any such criterion, the optimal family consists of concentric homothetic ellipsoids.
- Thus, such families of ellipsoids provide the optimal approximation to the actual surfaces:
  - at least in the *first* approximation, i.e.,
  - approximation corresponding to the smallest possible number of parameters.



## 93. Family of Sets: Towards a Precise Definition

- For each spatial point  $x$ , we would like to describe:
  - for each possible value  $\rho_0$  of the correlation  $\rho(x, y)$ ,
  - the set  $S_{\rho_0}(x) = \{y : \rho(x, y) \geq \rho_0\}$ .
- What are the natural properties of these families of sets?
- The first property is coverage.
- For each  $y$ , there is some value of  $\rho(x, y)$ .
- So for this  $x$ , the union of all these sets  $S_{\rho_0}(x)$  coincides with the whole space.
- The second property is monotonicity.
- If  $\rho(x, y) \geq \rho_0$  and  $\rho_0 \geq \rho'_0$ , then  $\rho(x, y) \geq \rho'_0$ .
- So, the sets  $S_{\rho_0}(x)$  should be inclusion-monotonic:

$$\text{if } \rho_0 \leq \rho'_0, \text{ then } S_{\rho'_0}(x) \subseteq S_{\rho_0}(x).$$

## 94. Family of Sets (cont-d)

- The third property is boundedness.
- From the physical viewpoint:
  - the further away is the point  $y$  from the point  $x$ ,
  - the less the physical quantities corresponding to these points are correlated.
- As the distance increases, this correlation should tend to 0.
- Thus, each set  $S_{\rho_0}(x)$  is *bounded*.
- The fourth property is continuity.
- In physics:
  - most processes are continuous,
  - with the exception of processes like fracturing, which we do not consider here.

## 95. Family of Sets (cont-d)

- We can therefore conclude that the correlation  $\rho(x, y)$  continuously depends on  $y$ , so:
  - if we have  $\rho(x, y_n) \geq \rho_0$  for some sequence of points  $y_n$  that converges to a point  $y$  ( $y_n \rightarrow y$ ),
  - then we should have  $\rho(x, y) = \lim_{n \rightarrow \infty} \rho(x, y_n) \geq \rho_0$ .
- Thus, if  $y_n \in S_{\rho_0}(x)$  and  $y_n \rightarrow y$ , then  $y \in S_{\rho_0}(x)$ .
- So, each set  $S_{\rho_0}(x)$  is *closed*.
- Similarly, it is reasonable to conclude that the set  $S_{\rho_0}(x)$  should continually depend on  $\rho_0$ :
  - if the two values  $\rho_0$  and  $\rho'_0$  are close,
  - then the corresponding sets  $S_{\rho_0}(x)$  and  $S_{\rho'_0}(x)$  should also be close.
- A natural way to describe closeness between (bounded closed) sets is to use the so-called Hausdorff distance.

## 96. Family of Sets (cont-d)

- We say that the sets  $A$  and  $B$  are  $\varepsilon$ -close if:
  - every point  $a \in A$  is  $\varepsilon$ -close to some point  $b \in B$ , i.e.,  $d(a, b) \leq \varepsilon$ ,  
and
  - every point  $b \in B$  is  $\varepsilon$ -close to some point  $a \in A$ .
- The Hausdorff distance  $d_H(A, B)$  is defined as the smallest  $\varepsilon$  for which the sets  $A$  and  $B$  are  $\varepsilon$ -closed.
- It can be shown that this distance can be equivalently defined as follows:

$$d_H(A, B) = \max \left( \sup_{a \in A} d(a, B), \sup_{b \in B} d(b, A) \right), \text{ where}$$
$$d(a, B) \stackrel{\text{def}}{=} \inf_{b \in B} d(a, b).$$

- What is the set of possible values of the parameter?
- In this family of sets, correlation value is a parameter.

## 97. Family of Sets (cont-d)

- Correlations can take any value from  $-1$  to  $1$ .
- When  $y = x$ , the correlation is clearly equal to  $1$ .
- When  $y \rightarrow \infty$ , we get values close to  $0$ .
- Since the function  $\rho(x, y)$  is continuous, this function takes all intermediate values.
- So, the possible values of the correlation form some interval.
- In some cases, we may have all possible negative values.
- In other cases, only some negative values, in yet other cases, we only have non-negative values.
- So, in general, we will consider all possible intervals of possible value of  $\rho_0$ .
- This interval may be closed – if there are points with limit correlation, or it can be open.

## 98. Definition

- So, we arrive at the following definition.
- Let  $N \geq 2$  be an integer.
- Let  $I$  be an interval.
- By a *family of sets*, we mean a set  $\{S_c : c \in I\}$  of bounded closed sets  $S_c \subseteq \mathbb{R}^N$  for which:
  - the dependence of  $S_c$  on  $c$  is continuous: if  $c_n \rightarrow c$ , then  $d_H(S_{c_n}, S_c) \rightarrow 0$ ;
  - the family  $S_c$  is monotonic: if  $c < c'$ , then  $S_{c'} \subseteq S_c$ ; and
  - the union of all the sets  $S_c$  coincides with the whole space.

## 99. Comments

- According to this definition, the family remains the same if we simply re-parameterize the family.
- For example:
  - instead of the original parameter  $c$ ,
  - we can use a new parameter  $c' = c + c_0$  or  $c' = \lambda \cdot c$  for some constants  $c_0$  and  $\lambda$ .
- In our specific problem, we are interested in the 3-D case  $N = 3$ .
- However, we can envision similar problem in the plane  $N = 2$  or in higher-dimensional spaces.
- So, in this talk, we consider the general case  $N \geq 2$ .

## 100. Comments (cont-d)

- We are specifically interested:
  - in *concentric homothetic families of ellipsoids*, i.e.,
  - in families of the type  $S_c = c \cdot E + a$ , where  $a$  is a given vector, and  $E$  is an ellipsoid with center 0.



## 101. Class of Families of Sets

- For different situation, in general:
  - we get different correlations and thus,
  - we get different families of sets.
- We would like to find a general class of such families that would, ideally, cover all such situations.
- We can use different parameters to differentiate different families from this class.
- In other words, a class can be described as a method for assigning:
  - to each possible combination of values of these parameters,
  - a specific family.
- As before, it makes sense to require that the resulting mapping is continuous.

## 102. Class of Families of Sets (cont-d)

- Here is a precise definition.
- Let  $N \geq 2$  and  $r > 0$  be integers.
- By an *r-parametric class of families of sets*, we mean a mapping that assigns,
  - to each element  $p = (p_1, \dots, p_r)$  from an open  $r$ -dimensional set  $D \subseteq \mathbb{R}^r$ ,
  - a family  $\{S_c(p)\}$  so that the dependence of  $S_c(p)$  on  $c$  and  $p$  is continuous.

### 103. Optimality Criteria: General Idea

- Out of all possible classes, we want to select a class which is, in some reasonable sense, optimal.
- For this, we need to be able to describe when some classes are better than others.
- In other words, we need to have an *order* on the set of all the classes.
- It would be nice to have a *total* (*linear*) order, in the sense that:
  - for every two classes,
  - we should be able to tell which one is better.
- However, it may be sufficient to have a *partial* order – as long as this order enables us to select the best class.
- It is OK if for some not-best classes, we do not have an opinion of which of them is better.

## 104. Optimality Criteria: General Idea (cont-d)

- In practice, usually, optimality criteria are described in numerical form:
  - we have an objective function  $f(a)$  that assigns a numerical value to each possible alternative  $a$ , and
  - we want to select an alternative for which this value is the largest possible,
  - or, depending on the context, the smallest possible.
- For example:
  - a company wants to maximize its profit,
  - a city wants to upgrade its road system so as to minimize the average travel time, etc.
- However, often, we need to go somewhat beyond this approach.

## 105. Optimality Criteria: General Idea (cont-d)

- For example, a company may have two (or more) different projects that lead to the same expected profit.
- In this case, we can use this non-uniqueness to optimize something else.
- For example:
  - out of all most-profitable projects,
  - we can select the one that leads to the smallest possible long-term environmental impact.
- In this case, we have a more complex criterion for comparing alternatives: we say that  $a$  is better if:
  - either  $f(a) > f(a')$
  - or  $f(a) = f(a')$  and  $g(a) > g(a')$ , for some other numerical criterion  $g(a)$ .

## 106. Optimality Criteria: General Idea (cont-d)

- If this still does not select us a unique alternative, we can optimize yet something else, etc.
- In view of this possibility, in this talk, we do not restrict ourselves to numerical optimization criteria.
- Instead, we use the most general definition of the optimality criterion, when:
  - for some pairs of alternatives  $a$  and  $a'$ , we know that  $a$  is better (we will denote it by  $a' < a$ ),
  - for some pairs of alternatives  $a$  and  $a'$ , we know that  $a'$  is better ( $a < a'$ ), and
  - for some pairs of alternatives  $a$  and  $a'$ ,  $a$  and  $a'$  are of the same value (we will denote it by  $a \sim a'$ ).
- Clearly, if  $a'$  is better than  $a$ , and  $a''$  is better than  $a'$ , then  $a''$  should be better than  $a$ , etc.

## 107. Optimality Criteria: General Idea (cont-d)

- Thus, we arrive at the following definition
- Let  $A$  be a set; elements of this set will be called *alternatives*.
- By an *optimality criterion*, we mean a pair of binary relations  $(<, \sim)$  on the set  $A$  for which:
  - if  $a < a'$  and  $a' < a''$ , then  $a < a''$ ;
  - if  $a < a'$  and  $a' \sim a''$ , then  $a < a''$ ;
  - if  $a \sim a'$  and  $a' < a''$ , then  $a < a''$ ;
  - if  $a \sim a'$  and  $a' \sim a''$ , then  $a \sim a''$ ;
  - if  $a \sim a'$ , then  $a' \sim a$ ;
  - if  $a < a'$ , then we cannot have  $a' < a$  or  $a \sim a'$ .
- Such a pair of relations is sometimes called a *partial pre-order*.

## 108. Optimality Criteria: General Idea (cont-d)

- Let  $(<, \sim)$  be an optimality criterion on a set  $A$ .
- An alternative  $a_{\text{opt}}$  is called *optimal* with respect to this criterion if for every alternative  $a \in A$ , we have

$$a < a_{\text{opt}} \text{ or } a \sim a_{\text{opt}}.$$



## 109. We Need A Final Optimality Criterion

- For the optimality criterion to be useful, it must select *at least one* optimal alternative.
- If the criterion selects *several* alternatives as optimal, this means that this criterion is not final.
- We can use the resulting non-uniqueness:
  - to optimize something else,
  - i.e., in effect, to come up with a better optimality criterion.
- If for this better criterion, we still have several optimal alternatives, we should modify this criterion again.
- Finally, we get a criterion for which there is exactly one optimal alternative.
- We will call such criteria *final*.

## 110. For Our Problem, an Optimality Criterion Must Be Affine-Invariant

- In our case, we want to compare different classes (of families of sets).
- In selecting optimality criteria, it is reasonable to take into account that:
  - while we want to deal with sets of points in physical space,
  - from the mathematical viewpoint, we deal with sets of tuples of real numbers.
- Real numbers (coordinates) describing each point depend on what coordinate system we use.
- If we select a different starting point, then all the coordinates are shifted  $x_i \rightarrow x_i + a_i$ .

## 111. Affine-Invariant (cont-d)

- If we select different axes for the coordinates, we get a rotation  $x_i \rightarrow \sum_{j=1}^N r_{ij} \cdot x_j$  for an appropriate matrix  $r_{ij}$ .
- These transformations make sense for the *isotropic* case, when:
  - all the properties of a material
  - are the same in all directions.
- Wood is an example of an *anisotropic* material.
- For example, it is easier to cut it along the orientation of the original tree than across that orientation.
- It is known that in many cases:
  - the description of an anisotropic material can be reduced to the isotropic case
  - if we apply an appropriate affine transformation.

## 112. Affine-Invariant (cont-d)

- This usually comes from the fact that, e.g.:
  - mechanical properties of a body can be described by a symmetric matrix, and
  - a symmetric matrix becomes symmetric if we use its eigenvectors as new axes.
- In view of this, it is reasonable to require that our optimality criterion is invariant:
  - not only with respect to shifts and rotations,
  - but also with respect to all possible affine (linear) transformations.
- Thus, we arrive at the following definitions.
- Let  $N > 2$  be an integer.

### 113. Affine-Invariant (cont-d)

- By an *affine transformation*, we mean
$$(Tx)_i = a_i + \sum_{j=1}^N b_{ij} \cdot x_j$$
for some reversible matrix  $b_{ij}$ .
- Let  $T$  be an affine transformation.
- Let  $S \subseteq \mathbb{R}^N$  be a set.
- By the *result*  $T(S)$  of applying  $T$  to  $S$ , we mean the set  $\{T(s) : s \in S\}$ .
- Let  $F = \{S_c : c \in I\}$  be a family of sets.
- By the *result*  $T(F)$  of applying  $T$  to  $F$ , we mean the family  $\{T(S_c) : c \in I\}$ .
- Let  $C = \{S_c(p)\}$  be class of families.
- By the *result*  $T(C)$  of applying  $T$  to  $C$ , we mean the class  $\{T(S_c(p))\}$ .

## 114. Affine-Invariant (cont-d)

- Let  $A$  be a set of alternatives, let  $(<, \sim)$  be an optimality criterion of the set  $A$ .
- Let  $\mathcal{T}$  be a class of transformations  $A \rightarrow A$ .
- We say that  $(<, \sim)$  is  $\mathcal{T}$ -invariant if for all  $T \in \mathcal{T}$  and  $a, a' \in A$ , we have:
  - if  $a < a'$  then  $T(a) < T(a')$ , and
  - If  $a \sim a'$ , then  $T(a) \sim T(a)$ .

## 115. Main Result

- Let  $N > 0$  and  $r > 0$  be integers.
- We consider sets in  $\mathbb{R}^N$ .
- Let  $(<, \sim)$  be a final affine-invariant optimality criterion on the set of all  $r$ -parametric classes of families.
- Then  $r \geq r_{\min} \stackrel{\text{def}}{=} \frac{N \cdot (N + 3)}{2} - 1$ , and:
  - for  $r = r_{\min}$ ,
  - the optimal class consists of concentric homothetic families of ellipsoids.
- This result indeed shows that:
  - the class of concentric homothetic families of ellipsoids
  - is the simplest (= fewer parameters) of all possible optimal classes.

## 116. Proof

- Since the optimality criterion is final, there exists exactly one optimal class  $C_{\text{opt}}$  for which:

$$C < C_{\text{opt}} \text{ or } C \sim C_{\text{opt}} \text{ for all classes } C.$$

- Let us prove that the optimal class  $C_{\text{opt}}$  is itself affine-invariant, i.e., that  $T(C_{\text{opt}}) = C_{\text{opt}}$  for each affine  $T$ .
- Indeed, due to optimality, for each class  $C$  and for each affine transformation class  $T$ , for  $T^{-1}(C)$ , we have:

$$T^{-1}(C) < C_{\text{opt}} \text{ or } T^{-1}(C) \sim C_{\text{opt}}.$$

- Since the criterion is affine-invariant, we have:

$$T(T^{-1}(C)) < T(C_{\text{opt}}) \text{ or } T(T^{-1}(C)) \sim T(C_{\text{opt}}).$$

- Here, by the definition of the inverse transformation:

$$T(T^{-1}(C)) = C.$$



## 117. Proof (cont-d)

- So we conclude that for every class  $C$ , we have:

$$C < T(C_{\text{opt}}) \text{ or } C \sim T(C_{\text{opt}}).$$

- By definition of optimality, this means that the class  $T(C_{\text{opt}})$  is optimal.
- However, our optimality criterion is final, which means that there is only one optimal class.
- Thus, indeed,  $T(C_{\text{opt}}) = C_{\text{opt}}$ .
- Since the optimal class is affine-invariant, with each family  $F$  this class also contains the family  $T(F)$ .
- This means that for each set  $S_c$  from each family, some family from the optimal class contains the set  $T(S_c)$ .
- Let us show that  $r \geq \frac{N \cdot (N + 3)}{2} - 1$ .

## 118. Proof (cont-d)

- Indeed, it is known that:
  - for every non-degenerate bounded set  $S$  (i.e., not contained in a proper subspace),
  - among all ellipsoids that contain  $S$ , there exists a unique ellipsoid of the smallest volume.
- It is also known that this correspondence between a set and the corresponding ellipsoid is affine-invariant:
  - if an ellipsoid  $E$  corresponds to the set  $S_c$ , then,
  - for each affine transformation  $T$ , to the set  $T(S_c)$  there corresponds the ellipsoid  $T(E)$ .
- It is known that every two ellipsoids can be obtained from each other by an affine transformation.

## 119. Proof (cont-d)

- Thus:
  - the family of all ellipsoids corresponding to all the sets from all the families
  - consists of all the ellipsoids.
- How many ellipsoids are there?
- A general ellipsoid can be determined by a quadratic formula  $\sum_{ij} a_{ij} \cdot$ 
$$x_i \cdot x_j + \sum_{i=1}^N a_i \cdot x_i \leq 1.$$
- Here,  $a_{ij}$  is a symmetric matrix  $a_{ij}$  and  $a_i$  is a vector.
- It is easy to see that different combinations of the matrix and the vector lead to different ellipsoids.
- We need  $N$  values  $a_1, \dots, a_N$  to describe a vector.

## 120. Proof (cont-d)

- Out of  $N^2$  elements of the matrix:
  - we need  $N$  values to describe its diagonal values  $a_{ii}$ , and
  - we need  $\frac{N^2 - N}{2}$  to describe non-diagonal elements.
- We divide by two since the matrix is symmetric:

$$a_{ij} = a_{ji}.$$

- Thus, overall, we need

$$N + N + \frac{N^2 - N}{2} = \frac{N \cdot (N + 3)}{2} \text{ values.}$$

- So, the set of all ellipsoids is:

$$\frac{N \cdot (N + 3)}{2}\text{-dimensional.}$$

## 121. Proof (cont-d)

- To each set  $S_c$  from families from the optimal class, we assign an ellipsoid.
- Thus, the dimension of the set of such sets should also be at least  $\frac{N \cdot (N + 3)}{2}$ -dimensional.
- These sets are divided into 1-parametric families.
- So the dimension  $r$  of the class of such families cannot be smaller than the above dimension minus 1.
- Thus, indeed,  $r \geq \frac{N \cdot (N + 3)}{2} - 1$ .

## 122. Proof (cont-d)

- Let us now prove that:
  - for the smallest possible dimension

$$r = r_{\min} \stackrel{\text{def}}{=} \frac{N \cdot (N + 3)}{2} - 1,$$

- all the sets  $S_c$  from the each family of the optimal class are ellipsoids.
- Indeed, we showed that each ellipsoid is associated with some set  $S_c$  from one of these families.
- The unit ball with a center at 0 is clearly an ellipsoid.
- Let us consider the set  $S_c$  which is associated with this unit ball.
- A unit ball is invariant with respect to all the rotations around its center.

### 123. Proof (cont-d)

- If the associated set  $S_c$  is not equal to the unit ball, this means that:
  - this set is not invariant
  - with respect to at least some rotations.
- In other words:
  - the group of all rotations that leave this set invariant
  - is a proper subgroup of the group of all rotations.
- This implies that the dimension of this group is smaller than the dimension of the group of all rotations.
- Thus, that there exists at least 1-parametric family  $\mathcal{R}$  of rotations  $R$  w.r.t. which the set  $S_c$  is not invariant.
- The optimal class is affine-invariant.

## 124. Proof (cont-d)

- Thus, all the sets  $R(S_c)$  are also sets from some family from the optimal class.
- For all these sets, the same unit ball is the smallest-volume ellipsoid.
- Thus, for this particular ellipsoid – the unit ball:
  - we have at least a 1-dimensional family of sets  $S_c$
  - associated with this ellipsoid.
- By applying a generic affine transformation:
  - we can find a similar at-least-1-dimensional family of sets
  - corresponding to each ellipsoid.



## 125. Proof (cont-d)

- Thus:
  - the dimension of the set of all sets  $S_c$
  - is at least one larger than the dimension of the family of all ellipsoids,
  - i.e. at least  $\frac{N \cdot (N + 3)}{2} + 1 = r_{\min} + 2$ .
- However, we have a  $r_{\min}$ -dimensional class of 1-dimensional families of sets.
- So the overall dimension of the set of all the sets  $S_c$  cannot be larger than  $r_{\min} + 1$ .
- This contradiction shows that the set  $S_c$  cannot be different from the enclosing minimal-volume ellipsoid.
- Thus, indeed, each set from each family from the optimal class is an ellipsoid.

## 126. Completing the Proof

- To complete the proof, we need to prove that ellipsoids in each family are concentric and homothetic.
- We have proven that each ellipsoid appears as an appropriate smallest-volume set.
- We know that each set  $S_c$  coincides with its smallest-volume enclosure.
- So, each ellipsoid appears as one of the sets  $S_c$  from one of the families from the optimal class.
- Let us again consider the unit ball centered at 0:
  - if the 1-dimensional family  $F_0$  containing this ball is not invariant with respect to all possible rotations,
  - then we have at least a 1-dimensional group of different families containing the same ellipsoid.

## 127. Completing the Proof (cont-d)

- We have:
  - an  $r_{\min}$ -dimensional class of 1-dimensional families
  - covering the whole  $(r_{\max} + 1)$ -dimensional family of ellipsoids.
- Thus, all elements of all families are different.
- So we cannot have several families containing the same ellipsoid.
- This argument shows that the family  $F_0$  containing the unit ball *should be* rotation-invariant.
- All the sets from this family are included in each other and thus, cannot be rotated into each other.
- This means that each ellipsoid from this family  $F_0$  must be rotation-invariant.

## 128. Completing the Proof (cont-d)

- This means that each ellipsoid from this family must be a ball concentric with our selected unit ball.
- Thus, it be homothetic to the original ball.
- For any other family  $F$ :
  - by selecting any ellipsoid  $E$  from this family and
  - by applying the affine transformation that transforms the above unit ball into  $E$ ,
  - we get a new family  $T(F_0)$  of concentric homothetic ellipsoids.
- An ellipsoid can only belong to one family.
- We thus conclude that the family  $F$  also consists of concentric homothetic ellipsoids.
- The result is proven.

## 129. Conclusions

- Wood is one the oldest construction materials; however:
  - in spite of several thousand years of experience with wooden constructions,
  - predicting and estimating mechanical properties of wooden constructions remains a difficult problem.
- One of the main reasons for this difficulty is that:
  - in contrast to many other constructions materials which are largely homogeneous and isotropic,
  - wood is highly inhomogeneous and anisotropic.
- Recently, a new property of wooden materials was discovered.
- It has a potential to make mechanical analysis of wooden structures more efficient.

## 130. Conclusions (cont-d)

- Namely, for wood:
  - iso-correlation surfaces (i.e., surfaces of equal correlation)
  - are well-approximated by concentric homothetic ellipsoids.
- The problem is that this property is purely empirical.
- It has no theoretical explanation and thus, engineers are understandably reluctant to rely on it.
- In this talk, we provide a theoretical explanation for this empirical fact.
- Thus, we make this property more reliable and therefore more useable.

Part IV

# Why Spiking Neural Networks Are Efficient: A Theorem

## 131. Why Spiking Neural Networks (NN)

- At this moment, artificial neural networks are the most successful – and the most promising – direction in AI.
- Artificial neural networks are largely patterned after the way the actual biological neural networks work.
- This patterning makes perfect sense:
  - after all, our brains are the result of billions of years of improving evolution,
  - so it is reasonable to conclude that many features of biological neural networks are close to optimal,
  - not very efficient features would have been filtered out in this long evolutionary process.
- However, there is an important difference between the current artificial NN and biological NN.



## 132. Why Spiking NN (cont-d)

- In hardware-implemented artificial NN each value is represented by the intensity of the signal.
- In contrast, in the biological neural networks, each value is represented by a frequency instantaneous spikes.
- Since simulating many other features of biological neural networks has led to many successes.
- So, a natural idea is to also try to emulate the spiking character of the biological neural networks.

### 133. Spiking Neural Networks Are Indeed Efficient

- Interestingly, adding spiking to artificial neural networks has indeed led to many successful applications.
- They were especially successful in processing temporal (and even spatio-temporal) signals.
- A biological explanation of the success of spiking neural networks makes perfect sense.
- However, it would be nice to supplement it with a clear mathematical explanation.
- It is especially important since:
  - in spite of all the billions years of evolution,
  - we humans are not perfect as biological beings,
  - we need medicines, surgeries, and other artificial techniques to survive, and
  - our brains often make mistakes.

## 134. Looking for Basic Functions

- In general, to represent a signal  $x(t)$  means to approximate it as a linear combination of some basic functions.
- For example, it is reasonable to represent a periodic signal as a linear combination of sines and cosines.
- Often, it makes sense to represent the observed values as a linear combination of:
  - functions  $t$ ,  $t^2$ , etc., representing the trend and
  - sines and cosines that describe the periodic part of the signal.
- We can also take into account that the amplitudes of the periodic components can also change with time.
- So, we end up with terms of the type  $t \cdot \sin(\omega \cdot t)$ .

## 135. Looking for Basic Functions (cont-d)

- For radioactivity, the observed signal is:
  - a linear combination of functions  $\exp(-k \cdot t)$
  - that represent the decay of different isotopes.
- So, in precise terms, selecting a representation means selecting an appropriate family of basic functions.
- In general, elements  $b(t)$  of a family can be described as  $b(t) = B(c_1, \dots, c_n, t)$  corr. to diff.  $c = (c_1, \dots, c_n)$ .
- Sometimes, there is only one parameter, as in sines and cosines.
- In control, typical are functions  $\exp(-k \cdot t) \cdot \sin(\omega \cdot t)$ , with two parameters  $k$  and  $\omega$ , etc.

## 136. Dependence on Parameters Is Continuous

- We want the dependence  $B(c_1, \dots, c_n, t)$  to be computable.
- It is known that all computable functions are, in some reasonable sense, continuous.
- Indeed, in real life, we can only determine the values of all physical quantities  $c_i$  with some accuracy.
- Measurements are always not 100% accurate, and computations always involve some rounding.
- For any given accuracy, we can provide the value with this accuracy.
- Thus, the approximate values of  $c_i$  are the only thing that  $B(c_1, \dots, c_n, t)$ -computing algorithm can use.
- This algorithm can ask for more and more accurate values of  $c_i$ .

## 137. Dependence Is Continuous (cont-d)

- However, at some point it must produce the result.
- At this point, we only know approximate values of  $c_i$ .
- So, we only know the interval of possible values of  $c_i$ .
- And for all the values of  $c_i$  from this interval:
  - the result of the algorithm provides, with the given accuracy,
  - the approximation to the desired value  $B(c_1, \dots, c_n, t)$ .
- This is exactly what continuity is about!
- One has to be careful here, since the real-life processes may actually be discontinuous.
- Sudden collapses, explosions, fractures do happen.

## 138. Dependence Is Continuous (cont-d)

- For example, we want to make sure that:
  - a step-function which is equal to 0 for  $t < 0$  and to 1 for  $t \geq 0$  is close to
  - an “almost” step function which is equal to 0 for  $t < 0$ , to 1 for  $t \geq \varepsilon$  and to  $t/\varepsilon$  for  $t \in (0, \varepsilon)$ .
- In such situations:
  - we cannot exactly describe the value at moment  $t$ ,
  - since the moment  $t$  is also measured approximately.
- What we can describe is its values at a moment close to  $t$ .

### 139. Dependence Is Continuous (cont-d)

- In other words, we can say that the two functions  $a_1(t)$  and  $a_2(t)$  are  $\varepsilon$ -close if:
  - for each  $t_1$ , there are  $\varepsilon$ -close  $t_{21}, t_{22}$  such that  $a_1(t_1)$  is  $\varepsilon$ -close to a convex combination of  $a_2(t_{2i})$ ;
  - for each  $t_2$ , there are  $\varepsilon$ -close  $t_{11}, t_{12}$  such that  $a_2(t_2)$  is  $\varepsilon$ -close to a convex combination of  $a_1(t_{1i})$ .



## 140. Additional Requirement

- We consider linear combinations of basic functions.
- So, it does not make sense to have two basic functions that differ only by a constant.
- If  $b_2(t) = C \cdot b_1(t)$ , then there is no need to consider the function  $b_2(t)$  at all.
- In each linear combination we can replace  $b_2(t)$  with

$$C \cdot b_1(t).$$

## 141. We Would Like to Have the Simplest Possible Family

- How many parameters  $c_i$  do we need? The fewer parameters:
  - the easier it is to adjust the values of these parameters, and
  - the smaller the probability of *overfitting* – a known problem of machine learning and data analysis in general.
- We cannot have a family with no parameters at all; this would mean, in effect, that:
  - we have only one basic function  $b(t)$  and
  - we approximate every signal by an expression  $C \cdot b(t)$  obtained by its scaling.

## 142. Simplest Possible Family (cont-d)

- This will be a very lousy approximation to real-life processes:
  - these processes are all different,
  - they do not resemble each other at all.
- So, we need at least one parameter.
- We are looking for the simplest possible family.
- So, we should therefore consider families depending on a single parameter  $c_1$ .
- In precise terms, we need functions  $b(t) = B(c_1, t)$  corresponding to different values of the parameter  $c_1$ .

## 143. Most Observed Processes Are Limited in Time

- From our viewpoint, we may view astronomical processes as going on forever.
- In reality, even they are limited by billions of years.
- In general, the vast majority of processes that we observe and that we want to predict are limited in time.
- A thunderstorm stops, a hurricane ends, after-shocks of an earthquake stop, etc.
- From this viewpoint:
  - to get a reasonable description of such processes,
  - it is desirable to have basic functions which are also limited in time,
  - i.e., which are equal to 0 outside some finite time interval.

## 144. Limited in Time (cont-d)

- This need for finite duration is one of the main reasons in many practical problems:
  - a decomposition into wavelets performs much better than
  - a more traditional Fourier expansion into linear combinations of sines and cosines.

## 145. Shift- and Scale-Invariance

- Processes can start at any moment of time.
- Suppose that we have a process starting at moment 0 which is described by a function  $x(t)$ .
- What if we start the same process  $t_0$  moments earlier?
- At each moment  $t$ , the new process  $x'(t)$  has been happening for the time period  $t + t_0$ , so  $x'(t) = x(t + t_0)$ .
- There is no special starting point.
- So it is reasonable to require that the class of basic function not change if we change the starting point:

$$\{B(c_1, t + t_0)\}_{c_1} = \{B(c_1, t)\}_{c_1}.$$

- Similarly, processes can have different speed.

## 146. Shift- and Scale-Invariance (cont-d)

- Some processes are slow, some are faster:
  - if a process starting at 0 is  $x(t)$ ,
  - then a  $\lambda$  times faster process is characterized by the function  $x'(t) = x(\lambda \cdot t)$ .
- There is no special speed.
- So it is reasonable to require that the class of basic function not change if we change the process's speed:

$$\{B(c_1, \lambda \cdot t)\}_{c_1} = \{B(c_1, t)\}_{c_1}.$$

- Now, we are ready for the formal definitions.

## 147. Definitions and the First Result

- We say that a function  $b(t)$  is *limited in time* if it equal to 0 outside some interval.
- We say that a function  $b(t)$  is a *spike* if it is different from 0 only for a single value  $t$ .
- This non-zero value is called the *height* of the spike.
- Let  $\varepsilon > 0$  be a real number.
- We say that the numbers  $a_1$  and  $a_2$  are  $\varepsilon$ -close if

$$|a_1 - a_2| \leq \varepsilon.$$

- We already had a definition of the functions  $a_1(t)$  and  $a_2(t)$  being  $\varepsilon$ -close.



## 148. Definitions and the First Result (cont-d)

- We say that a mapping  $B(c_1, t)$  is *continuous* if, for every  $c_1$  and  $\varepsilon > 0$ , there exists  $\delta > 0$  such that:
  - if  $c'_1$  is  $\delta$ -close to  $c_1$ ,
  - then the function  $b(t) = B(c_1, t)$  is  $\varepsilon$ -close to the function  $b'(t) = B(c'_1, t)$ .
- By a *family of basic functions*, we mean a continuous mapping for which:
  - for each  $c_1$ , the function  $b(t) = B(c_1, t)$  is limited in time, and
  - if  $c_1 \neq c'_1$ , then  $B(c'_1, t) \not\equiv C \cdot B(c_1, t)$ .
- We say that a family  $B(c_1, t)$  is *shift-invariant* if for each  $t_0$ :  
 $\{B(c_1, t)\}_{c_1} = \{B(c_1, t + t_0)\}_{c_1}$ .
- We say that a family  $B(c_1, t)$  is *scale-invariant* if for each  $\lambda > 0$ :  
 $\{B(c_1, t)\}_{c_1} = \{B(c_1, \lambda \cdot t)\}_{c_1}$ .

## 149. The First Result (cont-d)

- **Proposition.** *If a family of basic functions  $B(c_1, t)$  is shift- and scale-invariant, then:*
  - *for every  $c_1$ , the corresponding function  $b(t) = B(c_1, t)$  is a spike, and*
  - *all these spikes have the same height.*
- This result provides a possible explanation for the efficiency of spikes.

## 150. Proof

- Let us assume that the family of basic functions  $B(c_1, t)$  is shift- and scale-invariant.
- Let us prove that all the functions  $b(t) = B(c_1, t)$  are spikes.
- First, we prove that none of the functions  $B(c_1, t)$  is identically 0.
- Indeed, the zero function can be contained from any other function by multiplying by 0.
- This would violate the definition of a family of basic functions).
- Let us prove that each function from the given family is a spike.
- Indeed, each of the functions  $b(t) = B(c_1, t)$  is not identically zero, i.e., it attains non-zero values for some  $t$ .

## 151. Proof (cont-d)

- By definition, each of these functions is limited in time.
- So, the values  $t$  for which the function  $b(t)$  is non-zero are bounded by some interval.
- Thus, the values  $t_- \stackrel{\text{def}}{=} \inf\{t : b(t) \neq 0\}$  and  $t_+ \stackrel{\text{def}}{=} \sup\{t : b(t) \neq 0\}$  are finite, with  $t_- \leq t_+$ .
- Let us prove that we cannot have  $t_- < t_+$ .
- Indeed, in this case, the interval  $[t_-, t_+]$  is non-degenerate; thus:
  - by an appropriate combination of shift and scaling,
  - we will be able to get this interval from any other non-degenerate interval  $[a, b]$ .
- The family is shift- and scale-invariant.
- Thus, the correspondingly re-scaled function  $b'(t) = b(\lambda \cdot t + t_0)$  also belongs to the family  $B(c_1, t)$ .

## 152. Proof (cont-d)

- For this function, the corresponding values  $t'_-$  and  $t'_+$  will coincide with  $a$  and  $b$ .
- All these functions are different – so, we will have a 2-dimensional family of functions.
- This contradicts to our assumption that the family  $B(c_1, t)$  is one-dimensional.
- We cannot have  $t_- < t_+$ , so  $t_- = t_+$ , i.e., every function from our family is a spike.
- Let us prove that all the spikes have the same height.
- Indeed, let  $b_1(t)$  and  $b_2(t)$  be any two functions from the family.

### 153. Proof (cont-d)

- Both functions are spikes, so:
  - the value  $b_1(t)$  is only different from 0 for some value  $t_1$ , its height is  $h_1 \stackrel{\text{def}}{=} b_1(t_1)$ ;
  - similarly, the value  $b_2(t)$  is only different from 0 for some value  $t_2$ , its height is  $h_2 \stackrel{\text{def}}{=} b_2(t_2)$ .
- Since the family  $\mathcal{B}$  is shift-invariant, for  $t_0 \stackrel{\text{def}}{=} t_1 - t_2$ , the shifted function  $b'_1(t) \stackrel{\text{def}}{=} b_1(t + t_0)$  is also in  $\mathcal{B}$ .
- The shifted function is non-zero when  $t + t_0 = t_1$ , i.e., when  $t = t_1 - t_0 = t_2$ , and it has the same height  $h_1$ .
- If  $h_1 \neq h_2$ , we would have  $b'_1(t) = C \cdot b_2(t)$  for  $C \neq 1$ .
- Thus, the heights must be the same.
- The proposition is proven.

## 154. But Are Spiked Neurons Optimal?

- We showed that spikes naturally appear if we require reasonable properties like shift- and scale-invariance.
- This provides some justification for the spiked neural networks.
- However, the ultimate goal of neural networks is to solve practical problems.
- A practitioner is not interested in invariance or other mathematical properties.
- A practitioner wants to optimize some objective function.
- So, from the practitioner's viewpoint, the main question is: are spiked neurons optimal?

## 155. Different Practitioners Have Different Optimality Criteria

- In principle:
  - we can pick one such criterion (or two or three) and
  - analyze which families of basic functions are optimal with respect to these particular criterion.
- However, this will not be very convincing to a practitioner who has a different optimality criterion.
- An ideal explanation should work for *all* reasonable optimality criteria.
- To achieve this goal, let us analyze which optimality criteria can be considered reasonable.



## 156. What Is an Optimality Criterion: Analysis

- At first glance, the answer to this question may sound straightforward,
- We have an objective function  $J(a)$  that assigns, to each alternative  $a$ , a numerical value  $J(a)$
- We want to select an alternative for which the value of this function is the largest possible.
- If we are interested in minimizing losses, the value is the smallest possible.
- This formulation indeed describes many optimality criteria, but not all of them.
- Indeed, assume, for example, we are looking for the best method  $a$  for approximating functions.
- A natural criterion may be to minimize the mean squared approximation error  $J(a)$  of the method  $a$ .

## 157. What Is an Optimality Criterion (cont-d)

- If there is only one method with the smallest possible mean squared error, then this method is selected.
- But what if there are several different methods with the same mean squared error.
- This, by the way, is often the case.
- In this case, we can use this non-uniqueness to optimize something else; e.g., we can select:
  - out of several methods with the same mean squared error,
  - the method for which the average computation time  $T(a)$  is the smallest.
- The actual optimality criterion cannot be described by single objective function, it is more complex.

## 158. What Is an Optimality Criterion (cont-d)

- Namely, we say that a method  $a'$  is better than a method  $a$  if:
  - either  $J(a) < J(a')$ ,
  - or  $J(a) = J(a')$  and  $T(a) < T(a')$ .
- This additional criterion may still leave us with several equally good methods.
- We can use this non-uniqueness to optimize yet another criterion: e.g., worst-case computation time, etc.
- This criterion must enable us to decide which alternatives are better (or of the same quality).
- Let us denote this by  $a \leq a'$ .
- Clearly, if  $a \leq a'$  and  $a' \leq a''$ , then  $a \leq a''$ , so the relation  $\leq$  must be transitive (a.k.a. *pre-orders*).

## 159. An Optimality Criterion Must Be Final

- In terms of the relation  $\leq$ , optimal means better than (or of the same quality as) all other alternatives:

$$a \leq a_{\text{opt}} \text{ for all } a.$$

- If we have several optimal alternatives, then we can use this non-uniqueness to optimize something else.
- So, the corresponding criterion is not final.
- For a *final* criterion, we should have only one optimal alternative.

## 160. An Optimality Criterion Must Be Invariant

- In real life, we deal with real-life processes  $x(t)$ , in which values of different quantities change with time  $t$ .
- The corresponding numerical values of time  $t$  depend:
  - on the starting point that we use for measuring time and
  - on the measuring unit.
- For example, 1 hour is equivalent to 60 minutes.
- Numerical values are different, but from the physical viewpoint, this is the same time interval.
- We are interested in a universal technique for processing data.

## 161. Criterion Must Be Invariant (cont-d)

- It is therefore reasonable to require that:
  - the relative quality of different techniques should not change
  - if we change the starting point for measuring time or a measuring unit.
- Let us describe all this in precise terms.

## 162. Definitions and the Main Result

- Let a set  $A$  be given; its elements will be called *alternatives*.
- By an *optimality criterion*  $\leq$  on the set  $A$ , we mean a transitive relation (i.e., a *pre-order*) on this set.
- An element  $a_{\text{opt}}$  is called *optimal* with respect to the criterion  $\leq$  if for all  $a \in A$ , we have  $a \leq a_{\text{opt}}$ .
- An optimality criterion is called *final* if there exists exactly one optimal alternative.
- For each family  $B(c_1, t)$  and for each  $t_0$ , by its *shift*  $T_{t_0}(B)$ , we mean a family  $B(c_1, t + t_0)$ .
- We say that an optimality criterion on the class of all families is *shift-invariant* if
  - for every two families  $B$  and  $B'$  and for each  $t_0$ ,
  - $B \leq B'$  implies that  $T_{t_0}(B) \leq T_{t_0}(B')$ .

## 163. Definitions and the Main Result (cont-d)

- For each family  $B(c_1, t)$  and for each  $\lambda > 0$ , by its *scaling*  $S_\lambda(B)$ , we mean a family  $B(c_1, \lambda \cdot t)$ .
- We say that an optimality criterion on the class of families is *scale-invariant* if:
  - for every two families  $B$  and  $B'$  and for each  $\lambda > 0$ ,
  - $B \leq B'$  implies that  $S_\lambda(B) \leq S_\lambda(B')$ .
- **Proposition.**
  - Let  $\leq$  be a final shift- and scale-invariant optimality criterion on the class of all families of basic f-s.
  - Then, all elements of the optimal family are spikes of the same height.



## 164. Discussion

- Techniques based on representing signals as a linear combination of spikes are known to be very efficient.
- In different applications, efficiency mean different things: faster computations, more accurate results, etc.
- In different situations, we may have different optimality criteria.
- Our result shows that no matter what optimality criterion we use, spikes are optimal.
- This explains why spiking NN have been efficient in several different situations, with different criteria.

## 165. Proof

- Let us prove that the optimal family  $B_{\text{opt}}$  is itself shift- and scale-invariant.
- Then this result will follow from the previous Proposition.
- Indeed, let us consider any transformation  $T$  – be it shift or scaling.
- By definition of optimality, for any other family  $B$ , we have  $B \leq B_{\text{opt}}$ .
- In particular, for every  $B$ , this is true for  $T^{-1}(B)$ , i.e.,  $T^{-1}(B) \leq B_{\text{opt}}$ .
- Here,  $T^{-1}$  denotes the inverse transformation.
- Due to invariance,  $T^{-1}(B) \leq B_{\text{opt}}$  implies that  $T(T^{-1}(B)) \leq T(B_{\text{opt}})$ , i.e., that  $B \leq T(B_{\text{opt}})$ .

## 166. Proof (cont-d)

- This is true for each family  $B$ , thus the family  $T(B_{\text{opt}})$  is optimal.
- However, our optimality criterion is final, i.e., there is only one optimal family.
- Thus, we have  $T(B_{\text{opt}}) = B_{\text{opt}}$ .
- So, the optimal family  $B_{\text{opt}}$  is indeed invariant with respect to any of the shifts and scalings.
- Now, by applying the previous Proposition, we conclude the proof of this proposition.

## 167. Conclusions

- A usual way to process signals is to approximate each signal by a linear combinations of basic functions.
- Examples: sinusoids, wavelets, etc.
- In the last decades, a new approximation turned out to be very efficient in many practical applications.
- Namely, approximation of a signal by a linear combination of spikes.
- In this talk, we provide a possible theoretical explanation for this empirical success.

## 168. Conclusions (cont-d)

- Our main explanation is that:
  - for every reasonable optimality criterion on the class of all possible families of basic functions,
  - the optimal family is the family of spikes,
  - provided that the optimality criterion is scale- and shift-invariant.

Part V

# Why Most Empirical Distributions Are Few-Modal

## 169. Empirical Distributions: We Expect Them to Be Multi-Modal

- Continuous distributions are characterized by their probability density functions  $\rho(x)$ .
- In principle, a probability density function can be any non-negative function.
- The only condition is that the overall probability should be equal to 1, i.e., that  $\int \rho(x) dx = 1$ .
- In such situations, it is natural to expect that:
  - in general, we will observe generic functions with this property,
  - e.g., functions which are random with respect to some reasonable measure on the set of all functions.

## 170. Empirical Distributions (cont-d)

- The first such measure was Wiener measure, corresponding to random walk.
- Later, many other random measures have been proposed.
- In most of these random measures, almost all functions are truly random, similar to random walk.
- They are very “wiggly”, they have infinitely many local maxima and minima.
- In probabilistic terms, we expect the empirical probability density functions to be multi-modal.

## 171. Empirical Distributions Are Mostly Few-Modal

- In reality, empirical distributions are mostly either unimodal, or bimodal, or – in rare cases – trimodal.
- In other words, they are usually few-modal.
- Why?
- In science and engineering, the few-modality is often easy to explain.
- E.g., the distributions are normal or Gamma, or, in general, follow some theoretically justified law.
- But few-modal distributions are ubiquitous also:
  - in situations where we do not have exact equations,
  - such as econometrics.
- Why?
- In this talk, we provide a theoretical explanation for the few-modality of empirical distributions.



## 172. Main Idea

- Of course, the space of all possible probability density functions is infinite-dimensional.
- So to exactly describe each such function, we need to describe the values of infinitely many parameters.
- In practice, at each moment of time, we can only use finitely many parameters.
- So, we need to look into appropriate finite-dimensional families of probability density functions.
- And we need explain why functions from this appropriate family are few-modal.
- To answer this question, let us describe natural properties of such families  $F$  of distributions  $\rho(c_1, \dots, c_n.x)$ .
- How do we gain the information about the distributions?

## 173. We Want Smoothness

- It is reasonable to require that:
  - small changes in the values of the parameters  $c_i$  and/or small changes in  $x$
  - should lead to small changes in the probability density.
- In other words, we want the function  $\rho(c_1, \dots, c_n, x)$  to be smooth.

## 174. Combining Pieces of Knowledge

- Suppose that:
  - one piece of evidence is described by a probability density function (pdf)  $\rho_1(x)$ , and
  - another – independent – piece of evidence – leads to pdf  $\rho_2(x)$ .
- If these were evidences about two different quantities  $x_1$  and  $x_2$ , then:
  - due to independence, we would conclude that
  - the distribution of the pair  $(x_1, x_2)$  follows a product distribution  $\rho_1(x_1) \cdot \rho_2(x_2)$ .
- In our case, however, we know that this is the same quantity, i.e., that  $x_1 = x_2$ .
- Thus, to get the resulting distribution, we need to restrict the product distribution to the case  $x_1 = x_2$ .

## 175. Combining Pieces of Knowledge (cont-d)

- In precise terms, we need to consider conditional distribution under the condition that  $x_1 = x_2$ .
- This means that we need to consider the distribution

$$\rho(x) = c \cdot \rho_1(x) \cdot \rho_2(x).$$

- Here  $c$  is a normalizing constant – which can be determined by the condition that  $\int \rho(x) dx = 1$ .
- Thus, it is reasonable to require that:
  - for every two distribution  $\rho_1(x)$  and  $\rho_2(x)$  from the desired family  $F$ ,
  - their normalized product  $c \cdot \rho_1(x) \cdot \rho_2(x)$  should also belongs to this family.

## 176. Knowledge Can Come In Parts

- Sometimes, we gain the knowledge right away.
- In many other cases, knowledge comes in small steps.
- Suppose that:
  - the resulting knowledge is described by a probability density function  $\rho(x)$ , and
  - it comes via several ( $n$ ) independent similar pieces of knowledge,
  - each step characterized by some probability density function  $\rho_1(x)$ .
- Then, based on the previous subsection, we can conclude that  $\rho(x) = c \cdot (\rho_1(x))^n$  for some constant  $c$ .
- So,  $\rho_1(x) = c_1 \cdot (\rho(x))^{1/n}$  for an appropriate normalizing coefficient  $c_1$ .

## 177. Knowledge Can Come In Parts (cont-d)

- Thus, it is reasonable to require that:
  - for every distribution  $\rho_1(x)$  from the desired family  $F$  and
  - for every natural number  $n > 1$ ,
  - the normalized distribution  $c_1 \cdot (\rho(x))^{1/n}$  should also belong to the family.

## 178. Scale- and Shift-Invariance

- The numerical value of a quantity depends:
  - on the starting point for measuring this quantity
  - and on the measuring unit.
- When we change numerical values, the expression for the probability distribution also changes.
- It is reasonable to require that:
  - if we simply change the starting point and/or the measuring unit in a distribution from the family  $F$ ,
  - then we should still get a distribution from the same family.
- What if we change the starting point, i.e.,
  - we replace the original starting point
  - with a new one which is  $a$  units larger.

## 179. Scale- and Shift-Invariance (cont-d)

- Then in the new units  $y = x - a$ , the distribution:
  - described by pdf  $\rho(x)$
  - will now be described by  $\rho_1(y) = \rho(y + a)$ .
- Similarly, we can change the measuring unit, i.e.:
  - replace the original measuring unit
  - with a new one which is  $\lambda$  times larger.
- Then in the new units  $y = x/\lambda$ , the distribution
  - described by the pdf  $\rho(x)$
  - will now be described by  $\rho_1(y) = \lambda \cdot \rho(\lambda \cdot y)$ .



## 180. Definitions

- Let  $n$  be a natural number.
- By an  $n$ -parametric family of distributions, we mean a family  $F = \{f(c_1, \dots, c_n, x)\}_{c_1, \dots, c_n}$  of pdfs, where:
  - the values  $(c_1, \dots, c_n)$  go over some set  $U$ , and
  - the function  $f(c_1, \dots, c_n, x)$  is continuously differentiable over the closure of this set.
- We say that a family  $F$  allows combining knowledge if:
  - for every two functions  $\rho_1(x), \rho_2(x) \in F$ ,
  - there exists a real number  $c > 0$  for which the product  $c \cdot \rho_1(x) \cdot \rho_2(x)$  also belongs to  $F$ .

## 181. Definitions (cont-d)

- We say that a family  $F$  *allows partial knowledge* if:
  - for every function  $\rho(x)$  from this family and for every natural number  $n$ ,
  - there exists a real number  $c > 0$  for which the function  $c \cdot (\rho(x))^{1/n}$  also belongs to  $F$ .
- We say that a family  $F$  is *shift-invariant* if:
  - for every function  $\rho(x)$  from this family and for every real number  $a$ ,
  - the function  $\rho(x + a)$  also belongs to  $F$ .
- We say that a family  $F$  is *scale-invariant* if:
  - for every function  $\rho(x)$  from this family and for every real number  $\lambda > 0$ ,
  - the function  $\lambda \cdot \rho(\lambda \cdot x)$  also belongs to  $F$ .

## 182. Main Result

- **Proposition.**

- *Let  $F$  be a shift- and scale-invariant  $n$ -parametric family that allows combining and partial knowledge.*
- *Then, every function  $\rho \in F$  has the form  $\rho(x) = \exp(P(x))$  for some polynomial of degree  $\leq n$ .*

- **Corollary.**

- *Let  $F$  be a shift- and scale-invariant  $n$ -parametric family that allows combining and partial knowledge.*
- *Then, every function  $\rho \in F$  has no more than  $n - 1$  local maxima and local minima.*

- This explain why empirical distributions are few-modal.

### 183. Proof of the Corollary

- Indeed, at local maxima and minima, the derivative  $\rho'(x) = \exp(P(x)) \cdot P'(x)$  is equal to 0.
- This is equivalent to  $P'(x) = 0$ .
- The derivative  $P'(x)$  is a polynomial of degree  $\leq n - 1$ .
- Such polynomials can have no more than  $n - 1$  zeros.

## 184. Proof of the Main Result

- Let  $F$  be a family that satisfies all the given properties.
- To simplify the problem, let's consider a family  $G$  of all the functions  $c \cdot \rho(x)$ , where  $c > 0$  and  $\rho(x) \in F$ .
- By definition, every function from the family  $F$  is also an element of  $G$ .
- To show this, it is sufficient to take  $c = 1$ .
- We will prove the desired form for all the function from the class  $G$ .
- This will automatically imply that all the functions from the family  $F$  also have this property.
- What is the dimension of the family  $G$ ?
- I.e., how many parameters do we need to specify each function from this family?

## 185. Proof (cont-d)

- To describe a function from  $G$ , we need to specify:
  - the value  $c$  (1 parameter), and
  - the function  $\rho(x) \in F$  – which requires  $n$  parameters.
- Thus,  $n+1$  parameters are sufficient, and the dimension of the family  $G$  is  $\leq n+1$ .
- For the family  $G$ , allowing combining knowledge leads to a simpler property: that
  - for every two functions  $f_1(x), f_2(x) \in G$
  - their product  $f_1(x) \cdot f_2(x)$  also belong to  $G$ .
- Indeed,  $f_i(x) \in G$  means that  $f_i(x) = c_i \cdot \rho_i(x)$  for some  $c_i > 0$  and  $\rho_i(x) \in F$ .
- Thus, the product  $f(x) = f_1(x) \cdot f_2(x)$  of these functions has the form  $f(x) = c_1 \cdot c_2 \cdot \rho_1(x) \cdot \rho_2(x)$ .

## 186. Proof (cont-d)

- By the property of allowing combining knowledge, for some  $c > 0$ , we have  $\rho_0(x) = c \cdot \rho_1(x) \cdot \rho_2(x) \in F$ .
- Thus,  $f(x) = \frac{c_1 \cdot c_2}{c} \cdot (c \cdot \rho_1(x) \cdot \rho_2(x)) = c_0 \cdot \rho_0(x)$ , where  $c_0 \stackrel{\text{def}}{=} \frac{c_1 \cdot c_2}{c}$ .
- So indeed,  $f(x) \in G$ .
- Similarly, from the other properties of the family  $F$ , we can make the following conclusions:
  - that for every function  $f(x) \in G$  and for every natural number  $n$ , we have  $(f(x))^{1/n} \in G$ ;
  - that for every function  $f(x) \in G$  and for every real number  $a$ , we have  $f(x + a) \in G$ ;
  - that for every function  $f(x) \in G$  and for every real number  $\lambda > 0$ , we have  $f(\lambda \cdot x) \in G$ .

## 187. Proof (cont-d)

- We can simplify the problem even more if:
  - instead of the family  $G$ ,
  - we consider the family  $g$  of all the functions of the type  $F(x) = \ln(f(x))$ , where  $f(x) \in G$ .
- To such functions, we also add the limit functions.
- Adding limit cases does not increase the dimension, so the dimension of the family  $g$  is still  $\leq n + 1$ .
- In terms of this new family, we need to prove that all the functions from  $g$  are polynomials of order  $\leq n$ .
- The fact that the family  $G$  is closed under multiplication means that the family  $g$  is closed under addition.



## 188. Proof (cont-d)

- The fact that the family  $G$  is closed under taking the  $n$ -th root means that:
  - the family  $g$  is closed
  - under multiplication by  $1/n$  for each natural number  $n$ .
- Together with closing under addition, this means that:
  - for every two natural numbers  $m$  and  $n$ ,
  - the function  $\frac{m}{n} \cdot F(x) = \frac{1}{n} \cdot F(x) + \dots + \frac{1}{n} \cdot F(x)$  ( $m$  times) also belongs to the family  $g$ .
- In other words, for every  $F(x) \in g$  and for every rational number  $r$ , we have  $r \cdot F(x) \in g$ .
- Every real number is a limit of rational numbers.
- E.g., it is a limit of numbers obtained if we only keep the first  $N$  digits in the decimal or binary expansion.

## 189. Proof (cont-d)

- Since we added all limit cases, we can conclude that  $r \cdot F(x) \in g$  for all non-negative real numbers  $r$  as well.
- One can easily show that shift- and scale-invariance properties are also satisfied for the new family:
  - that for every function  $F(x) \in g$  and for every real number  $a$ , we have  $F(x + a) \in g$ ;
  - that for every function  $F(x) \in g$  and for every real number  $\lambda > 0$ , we have  $F(\lambda \cdot x) \in g$ .
- As a final simplification, we consider the family  $h$  of all the differences  $d(x) = F_1(x) - F_2(x)$  between  $F_i(x) \in g$ .
- To describe each of the functions  $F_1(x)$  and  $F_2(x)$ , we need  $n + 1$  parameters.
- So the dimension of the new family does not exceed  $2 \cdot (n + 1)$ .

## 190. Proof (cont-d)

- For every function  $F(x) \in g$ , the function  $2F(x)$  also belongs to the family  $g$ .
- So, we can conclude that the difference  $F(x) = (2F(x)) - F(x)$  also belongs to the family  $h$ . Thus,  $g \subseteq h$ .
- The family  $h$  is also closed under addition.
- Indeed, if  $d_1(x) = F_{11}(x) - F_{12}(x)$  and  $d_2(x) = F_{21}(x) - F_{22}(x)$  for some  $F_{ij}(x) \in g$ , then

$$\begin{aligned}d_1(x) + d_2(x) &= (F_{11}(x) - F_{12}(x)) + (F_{21}(x) - F_{22}(x)) = \\&= (F_{11}(x) + F_{21}(x)) - (F_{12}(x) + F_{22}(x)).\end{aligned}$$

- Since  $g$  is closed under addition, the sums  $F_{11}(x) + F_{21}(x)$  and  $F_{12}(x) + F_{22}(x)$  also belong to  $g$ .
- Thus, indeed, the sum  $d_1(x) + d_2(x)$  is a difference between two functions from  $g$  and is, thus, in  $h$ .

## 191. Proof (cont-d)

- We can also prove that the family  $h$  is closed under multiplication by any real number  $c$ .
- Indeed, let  $d(x) = F_1(x) - F_2(x)$ .
- If  $c > 0$ , then  $c \cdot d(x) = (c \cdot F_1(x)) - (c \cdot F_2(x))$ , where both  $c \cdot F_1(x)$  and  $c \cdot F_2(x)$  belong to the family  $g$ .
- If  $c < 0$ , then  $c \cdot F(x) = |c| \cdot F_2(x) - |c| \cdot F_1(x)$ , where also  $|c| \cdot F_2(x)$  and  $|c| \cdot F_1(x)$  belong to the family  $g$ .
- So, the family  $h$  is closed under addition and under multiplication by any real number.
- Thus,  $h$  is a linear space.
- Let  $d \leq 2n + 2$  denote the dimension of this linear space.
- Let us select a basis  $e_1(x), \dots, e_d(x)$ .

## 192. Proof (cont-d)

- This means that all functions from the space  $g$  have the form  $c_1 \cdot e_1(x) + \dots + c_d \cdot e_d(x)$ .
- We know that the family  $g$  is shift- and scale-invariant.
- Thus, we can conclude that the family  $h$  is also shift- and scale-invariant.
- Shift-invariance means that for each  $d(x) \in h$  and for each real number  $a$ , we have  $d(x + a) \in h$ .
- In particular, this is true for the basis functions

$$e_1(x), \dots, e_d(x).$$

- Thus, for each  $i$  and  $a$ , there exist coefficients  $c_{ij}(a)$  depending on  $a$  for which

$$e_i(x + a) = c_{i1}(a) \cdot e_1(x) + \dots + c_{id}(a) \cdot e_d(x).$$

### 193. Proof (cont-d)

- In particular, for each  $i$ , we can select  $d$  different values

$$x_1, \dots, x_d.$$

- Then we get the following system of  $d$  linear equations for determining the coefficients  $c_{ij}(a)$ :

$$e_i(x_1 + a) = c_{i1}(a) \cdot e_1(x_1) + \dots + c_{id}(a) \cdot e_d(x_1),$$

$$\dots$$

$$e_i(x_d + a) = c_{i1}(a) \cdot e_1(x_d) + \dots + c_{id}(a) \cdot e_d(x_d).$$

- Here, the coefficients  $e_j(x_k)$  are constants.
- So the values  $c_{ij}(a)$  are linear combinations of the right-hand sides  $e_i(x_k + a)$ .
- Since the functions  $e_i(x)$  are differentiable, the values  $c_{ij}(a)$  are also differentiable functions of  $a$ .

## 194. Proof (cont-d)

- So, both sides of the following equality are differentiable:  $e_i(x+a) = c_{i1}(a) \cdot e_1(x) + \dots + c_{id}(a) \cdot e_d(x)$ .
- Thus, we can differentiate them with respect to  $a$  and then plug in  $a = 0$ .
- As a result, we get the following system of differential equations, where  $C_{ij} \stackrel{\text{def}}{=} c'_{ij}(0)$ :

$$e'_1(x) = C_{11} \cdot e_1(x) + \dots + C_{1d} \cdot e_d(x),$$

...

$$e'_d(x) = C_{d1} \cdot e_1(x) + \dots + C_{dd} \cdot e_d(x),$$

- In other words, for  $e_i(x)$ , we get a system of linear differential equations with constant coefficients.

## 195. Proof (cont-d)

- It is known that each solution of such system is a linear coefficient of the functions  $x^p \cdot \exp(\alpha \cdot x)$ , where:
  - the value  $p$  is a natural number and
  - $\alpha$  is a – possible complex – eigenvalue of the matrix  $C_{ij}$ .
- Similarly, scale-invariance means that for each function  $d(x) \in h$  and for each positive real number  $\lambda > 0$ , we have  $d(\lambda \cdot x) \in h$ .
- In particular, this is true for the basis functions  $e_i(x)$ .
- For an auxiliary variable  $X \stackrel{\text{def}}{=} \ln(x)$ :
  - replacing  $x$  with  $\lambda \cdot x$  corresponds to
  - replacing  $X$  with  $X + a$ , where  $a \stackrel{\text{def}}{=} \ln(\lambda)$ .



## 196. Proof (cont-d)

- So, for the correspondingly re-scaled functions  $E_i(X) \stackrel{\text{def}}{=} e_i(\exp(X))$ , we conclude that:
  - for each such function and for each real number  $a$ ,
  - the function  $E_i(X+a)$  is a linear combination of functions  $E_1(X), \dots, E_d(X)$ .
- We already know, from the previous parts of this proof, that this implies that:
  - each function  $E_i(X)$
  - is a linear combination of the functions  $X^p \cdot \exp(\alpha \cdot X)$ .
- Thus, each function  $e_i(x) = E_i(\ln(x))$  is a linear combination of expressions

$$(\ln(x))^p \cdot \exp(\alpha \cdot \ln(x)) = (\ln(x))^p \cdot x^\alpha.$$

## 197. Proof (cont-d)

- One can see that:
  - the only possibility for a function to be represented in both forms
  - is to avoid logarithms and exponential functions altogether.
- So,  $e_i(x)$  is a linear combination of the terms  $x^p$  for natural  $p$ , i.e., a polynomial.
- Thus, each function from the class  $g$  is a polynomial, as a linear combination of  $d$  polynomials  $e_i(x)$ .
- Since  $g \subseteq h$ , all functions from the class  $g$  are also polynomials.
- What is the order of these polynomials?
- Let  $D$  be the order of a polynomial  $F(x)$  from the  $g$ .

## 198. Proof (cont-d)

- For a polynomial of order  $D$ , in general,  $F(x)$ ,  $F(x + h)$ ,  $F(x + 2h)$ ,  $\dots$ ,  $F(x + D \cdot h)$  are linearly independent.
- Indeed, for  $h \rightarrow 0$ , this is equivalent to linear independence of  $x^D$ ,  $x^{D-1}$ ,  $\dots$ ,  $1$ .
- Thus, in the generic case, the corresponding determinant is different from 0.
- Since we have  $D + 1$  independent functions, thus, the family  $g$  has dimension  $D + 1$ .
- But we know that the dimension of this family is  $\leq n + 1$ .
- From  $D + 1 \leq n + 1$ , we conclude that  $D \leq n$ .
- Thus, all functions  $F(x) = \ln(f(x))$  from the class  $g$  are polynomials of order  $\leq n$ .

## 199. Proof (cont-d)

- Thus, all functions  $F(x) = \ln(f(x))$  from the class  $g$  are polynomials of order  $\leq n$ .
- Hence, each function  $f(x) = \exp(F(x))$  from the class  $F$  has the desired form.
- The proposition is proven.

## Part VI

# Why a Classification Based on Linear Approximation to Dynamical Systems Often Works Well in Nonlinear Cases

## 200. Dynamical Systems Are Ubiquitous

- We want to describe the state of a real-life system at any given moment of time.
- So, we need to know the values  $x = (x_1, \dots, x_n)$  of all the quantities that characterize this system.
- For example:
  - to describe the state of a mechanical system consisting of several pointwise objects,
  - we need to know the position and velocities of all these objects.
- To describe the state of an electric circuit, we need to know the currents and voltages, etc.

## 201. Dynamical Systems Are Ubiquitous (cont-d)

- In many real-life situation, the corresponding systems are deterministic – in the sense that:
  - the future states of the system
  - are uniquely determined by its current state.
- Sometimes, to make the system deterministic:
  - we need to enlarge its description
  - so that it incorporates all the objects that affect its dynamics.
- For example:
  - the system consisting of Earth and Moon is not deterministic in its original form,
  - since the Sun affects its dynamics.
- However, once we add the Sun, we get a system with a deterministic behavior.

## 202. Dynamical Systems Are Ubiquitous (cont-d)

- That the future dynamics of the system is uniquely determined by its current state means, in particular:
  - that the rate  $\dot{x}$  with which the system changes is also uniquely determined by its current state,
  - i.e., that we have  $\dot{x} = f(x)$ , for some function  $f(x)$ .
- This equation can be described coordinate-wise, as

$$\dot{x}_i = f_i(x_1, \dots, x_n).$$

- Systems that satisfy such equations are known as *dynamical systems*.

## 203. Simplest Case: Linear Systems

- The simplest case is when the rate of change  $f_i(x_1, \dots, x_n)$  of each variables is a linear function, i.e., when

$$\dot{x}_i = a_{i0} + \sum_{j=1}^n a_{ij} \cdot x_j.$$

- In almost all such cases, the matrix  $a_{ij}$  is non-degenerate.
- Then, we can select constants  $s_i$  so that:
  - for the shifted variables  $y_i = x_i + s_i$ ,
  - the system gets a simpler form  $\dot{y}_i = \sum_{j=1}^n a_{ij} \cdot y_j$ .
- Indeed, substituting  $x_i = y_i - s_i$  into the above formula, and taking into account that  $\dot{y}_i = \dot{x}_i$ , we conclude that

$$\dot{x}_i = a_{i0} + \sum_{j=1}^n a_{ij} \cdot (y_j - s_j) = a_{i0} + \sum_{j=1}^n a_{ij} \cdot y_j - \sum_{j=1}^n a_{ij} \cdot s_j.$$



## 204. Simplest Case: Linear Systems (cont-d)

- Thus, if we select the value  $s_j$  for which  $a_{i0} = \sum_{j=1}^n a_{ij} \cdot s_j$  for each  $i$ , we will indeed get the desired formula.
- For the linear equation, the general solution is known: it is a linear combination of expressions  $t^k \cdot \exp(\lambda \cdot t)$ :
  - where  $\lambda$  is an eigenvalue of the matrix  $\|a_{ij}\|$  – which is, in general, a complex number  $\lambda = a + i \cdot b$ ,
  - and  $k$  is a natural number which does not exceed the multiplicity of this eigenvalue.
- In real-number terms, we get a linear combination of the expressions  $t^k \cdot \exp(a \cdot t) \cdot \sin(b \cdot t + \varphi)$ .
- Depending on the values of  $\lambda$ , we have the following types of behavior.

## 205. Simplest Case: Linear Systems (cont-d)

- When  $a < 0$  for all the eigenvalues, then the system is *stable*:
  - no matter what state we start with,
  - it asymptotically tends to the state

$$y_1 = \dots = y_n = 0.$$

- When  $a > 0$  for at least one eigenvalue, then the system is *unstable*.
- In this case, the deviation from the 0 state exponentially grows with time.
- When  $a = 0$  and  $b \neq 0$ , we get an *oscillatory behavior*.

## 206. Simplest Case: Linear Systems (cont-d)

- When  $a = b = 0$ , we get a *transitional behavior*, when a system:
  - linearly (or quadratically etc.) moves
  - from one state to another.
- Interestingly, a similar classification works well for nonlinear dynamical systems as well, but why?
- In this talk, we will try to explain this fact.

## 207. We Need Finite-Dimensional Approximations

- We want to describe how the state  $x(t) = (x_1(t), \dots, x_n(t))$  of a dynamical system changes with time  $t$ .
- In general, the set of all possible smooth functions  $x_i(t)$  is infinite-dimensional.
- In other words, we need infinitely many parameters to describe it.
- However, in practice, at any given moment, we can only have finitely many parameters.
- Thus, it is reasonable to look for finite-parametric approximations.

## 208. Finite-Dimensional Approximations (cont-d)

- A natural idea is:
  - to fix some smooth functions  $e_k(t) = (e_{k1}(t), \dots, e_{kn}(t))$ ,  $1 \leq k \leq K$ , and
  - consider linear combinations

$$x(t) = \sum_{k=1}^K c_k \cdot e_k(t).$$

## 209. Shift-Invariance

- For dynamical systems, there is no fixed moment of time.
- The equations remain the same:
  - if we change the starting point for measuring time,
  - i.e., if we replace the original temporal variable  $t$  with the new variable  $t' = t + t_0$ .
- It is therefore reasonable to require that:
  - the approximating family be invariant
  - with respect to the same transformation.
- In other words, we require that all shifted functions  $e_k(t + t_0)$  can also be represented in the same form.
- Let us show that this reasonable requirement explains the above phenomenon.

## 210. Towards the Explanation

- Reminder: for each  $i$ , we have  $x_i(t) = \sum_{k=1}^K c_k \cdot e_{ki}(t)$ .
- The fact that shifted functions can be represented in this form means that for each  $k$ ,  $i$ , and  $t_0$ , we have

$$e_{ki}(t + t_0) = \sum_{\ell=1}^K c_{k\ell i}(t_0) \cdot e_{\ell i}(t), \text{ for some coefficients } c_{k\ell i}(t_0).$$

- Let us fix  $i$  and  $k$  and select  $K$  different moments of time  $t_m$ ,  $m = 1, \dots, K$ .
- For these moments of time, we get:

$$e_{ki}(t_m + t_0) = \sum_{\ell=1}^K c_{k\ell i}(t_0) \cdot e_{\ell i}(t_m).$$

- Thus, we get  $K$  linear equations for determining  $K$  unknowns  $c_{k1i}(t_0), \dots, c_{kKi}(t_0)$ .

## 211. Towards the Explanation (cont-d)

- Cramer's formula:
  - describes the solution to a system of linear equations
  - as a rational (and thus, smooth) function of its coefficients and right-hand sides.
- Thus, each coefficient  $c_{k\ell i}(t_0)$  is a smooth function of the values  $e_{ki}(t_m + t_0)$  and  $e_{\ell i}(t_m)$ .
- Since the functions  $e_{ki}(t)$  are smooth, the dependence of the coefficients  $c_{k\ell i}(t_0)$  on  $t_0$  is also differentiable.
- All the functions involved in the formula
$$x_i(t) = \sum_{k=1}^K c_k \cdot e_{ki}(t)$$
are differentiable.
- So we can differentiate this formula with respect to  $t_0$  and get  $\dot{e}_{ki}(t + t_0) = \sum_{\ell=1}^K \dot{c}_{k\ell i}(t_0) \cdot e_{\ell i}(t)$ .



## 212. Towards the Explanation (cont-d)

- In particular, for  $t_0 = 0$ , we get

$$\dot{e}_{ki}(t) = \sum_{\ell=1}^K a_{k\ell i} \cdot e_{\ell i}(t), \text{ where } a_{k\ell i} \stackrel{\text{def}}{=} \dot{c}_{k\ell i}(t_0).$$

- So, the functions  $e_{ki}(t)$  satisfy the system of linear differential equations with constant coefficients.
- We have already mentioned that:
  - the solutions to such systems
  - are exactly the functions leading to a known classification of linear dynamical system behaviors.
- This explains why for nonlinear systems, we also naturally observe similar types of behavior.

Part VII

# When Revolutions Happen: Algebraic Explanation

## 213. When Revolutions Happen

- People usually believe that revolutions happen when life under the old regime becomes intolerable.
- However, a historical analysis shows that the usual understanding is wrong.
- Most revolutions happen *not* when the situation is at its worst.
- They usually happen when the situation has been improving for some time and then suddenly gets worse.
- Although, by the way, it never gets as bad as it was before the improvement started.

## 214. How Can We Explain This?

- Experiments show that in most situations, people act rationally:
  - the more their needs are satisfied, in general,
  - the happier they are.
- So why right before the revolution:
  - when the level of living is higher (often much higher) than in the recent past,
  - people are so much less happy that they start a revolution?
- How can we explain this unexpected (and somewhat counterintuitive) behavior?

## 215. Traditional Decision Theory: A Brief Reminder

- In traditional decision theory, people's preferences are described by numerical values called *utilities*.
- The actions of a person are determined:
  - not just by this person's current level of satisfaction – as described by the current utility value  $u_0$ ,
  - but also by the expected future utility values  $u_1$ ,  $u_2$ , etc.
- If we have  $m$  dollars, we can place it in a bank and get  $(1 + \alpha)^t \cdot m$  at time  $t$ , where  $\alpha$  is the interest rate.
- Thus, \$1 at time  $t$  is equivalent to  $q^t$  dollars now, where  $q \stackrel{\text{def}}{=} \frac{1}{1 + \alpha}$ .
- So, if we get  $m_0$  now,  $m_1$  in the next year, etc., this is equivalent to getting the following amount now:

$$m_0 + q \cdot m_1 + q^2 \cdot m_2 + \dots$$

## 216. This General Approach Requires Extrapolation

- The future amounts are based on extrapolation:
  - we select a family of functions characterized by a few parameters
$$u_t = f(p_1, \dots, p_n, t),$$
  - then we find the values  $\hat{p}_1, \dots, \hat{p}_n$  of the parameters that best fit the observed data  $u_0, u_{-1}, u_{-2}$ , etc.,
  - and then we use these values to predict future values as  $f(\hat{p}_1, \dots, \hat{p}_n, t)$ .
- Let's use models that linearly depend on  $p_i$ :
  - then, matching parameters to data means easy-to-solve solving systems of linear equations,
  - while solving systems of nonlinear equations is, in general, NP-hard.
- Thus, we consider models  $u_t = \sum_{i=1}^n p_i \cdot f_i(t)$ , where  $f_i(t)$  are given functions, and  $p_i$  are parameters.

## 217. Which Basis Functions $f_i(t)$ Should We Choose?

- Most transitions are smooth; so, it's reasonable to require that all the functions  $f_i(t)$  are smooth.
- Another reasonable requirement is related to the fact that the numerical value of time depends:
  - on the choice of a measuring unit – years or months,
  - and on the choice of a starting time.
- If we change a measuring unit by a new one which is  $a$  times smaller, then  $t \rightarrow a \cdot t$ .
- If we replace the original starting point with the new one,  $b$  units in the past  $t \rightarrow t + b$ .
- The general formulas for extrapolation should not depend on such an arbitrary things as:
  - selecting a unit of time or
  - selecting a starting point.

## 218. Choosing $f_i(t)$ (cont-d)

- It is therefore reasonable to assume that the approximating family  $\left\{ \sum_{i=1}^n p_i \cdot f_i(t) \right\}$  will not change:

$$\left\{ \sum_{i=1}^n p_i \cdot f_i(a \cdot t) \right\}_{p_1, \dots, p_n} = \left\{ \sum_{i=1}^n p_i \cdot f_i(t + b) \right\}_{p_1, \dots, p_n} = \left\{ \sum_{i=1}^n p_i \cdot f_i(t) \right\}_{p_1, \dots, p_n}.$$

- It turns out that under these conditions, all the basic functions are polynomials.
- So, all their linear combinations are polynomials.
- Thus, it is reasonable to approximate the actual history by a polynomial.



## 219. Two Simple Situations

- We will compare two simple situations:
  - a situation in which the level of living is consistently bad  $u_0 = u_{-1} = \dots = u_{-k} = \dots = c_1$  for small  $c_1$ ;
  - a situation in which the level of living used to be much better, but now somewhat decreased:

$$u_{-1} = u_{-2} = \dots = c_+ \text{ but } u_0 = c_- < c_+.$$

- In the first situation, of course, a reasonable extrapolation should lead to the exact same small value  $u_0 = c$ .
- Thus, the overall utility is equal to

$$u_0 + q \cdot u_1 + \dots = c \cdot (1 + q + q^2 + \dots) = \frac{c}{1 - q}.$$

- But what to expect in the second situation?
- Let us start our analysis with the simplest possible linear extrapolation.

## 220. Linear Extrapolation

- In this case, we make our future predictions based only on two utility values:  $u_0$  and  $u_{-1}$ .
- Since  $u_0 < u_{-1}$ , we get a linear decreasing function.
- Its values tend to  $-\infty$  as the time  $t$  increases.
- So, when  $q$  is close to 1, the corresponding value

$u_0 + q \cdot u_1 + \dots \approx u_0 + u_1 + \dots$  becomes negative.

- This explains why in the second situation, the revolution is much more probable.

## 221. What About More Realistic Approximations?

- One may think that the above explanation is caused by our oversimplification of the extrapolation model.
- Of course, linear extrapolation is a very crude and oversimplified idea.
- What happens if we use higher degree polynomials for extrapolation?
- Let us assume that for extrapolation, we use polynomials of order  $d$ .
- The corresponding family of polynomials have  $d+1$  parameters, so we can fit  $d+1$  values  $u_0, u_{-1}, \dots, u_{-d}$ .
- Let us find the polynomial  $P(t)$  of degree  $d$  that fits these values:  $P(0) = c_-, P(-1) = \dots = P(-d) = c_+$ .

## 222. Realistic Approximations (cont-d)

- For  $Q(t) \stackrel{\text{def}}{=} P(t) - c_+$ , we have  $Q(-d) = \dots = Q(-1) = 0$  and  $Q(0) = c_- - c_+$ .
- This polynomial of degree  $d$  has  $d$  roots  $t = -1, \dots, t = -d$ , so  $Q(t) = C \cdot (t + 1) \cdot (t + 2) \cdot \dots \cdot (t + d)$ , and

$$Q(t) = c_+ + (c_- - c_+) \cdot \frac{(t + 1) \cdot (t + 2) \cdot \dots \cdot (t + d)}{1 \cdot 2 \cdot \dots \cdot d}.$$

- Since  $c_- < c_+$ , this value is negative – and tends to  $-\infty$  as the time  $t$  increases.
- In comparison with the linear extrapolation case, it tends to  $-\infty$  even faster: as  $t^d$ .
- So, *the revolution phenomenon can be explained* no matter what degree of extrapolation we use.

## 223. Discussion

- We have explained the seemingly counterintuitive revolution phenomenon.
- Based on our analysis, we can make auxiliary conclusions (which also fit well with common sense).
- Revolutions only happen if people care about the future.
- If they don't, if  $q \approx 0$ , people are happy with their present-day level of living.
- The more into the past the people go in their analysis, the more probable it is that they will revolt.
- People who do not know their history are less prone to revolutions than people who do.

Part VIII

# How Mathematics and Computing Can Help Fight the Pandemic: Two Pedagogical Examples

## 224. First Example: Need for Social Distancing

- This problem is related to the pandemic-related need to observe a social distance of at least 2 meters (6 feet).
- Two persons are on two sides of a narrow-walkway street, waiting for the green light.
- They start walking from both sides simultaneously.
- For simplicity, let us assume that they walk with the same speed.
- If they follow the shortest distance path – i.e., a straight line  $AB$  – they will meet in the middle.
- This is not good; so one of them should move somewhat to the left, another should move somewhat to the right.
- At all moments of time, they should be at least 2 meters away from each other.
- What is the fastest way for them to do it?

## 225. Formulating This Problem in Precise Terms

- The situation is absolutely symmetric with respect to the reflection in the midpoint  $M$  of the segment  $AB$ .
- So, it is reasonable to require that:
  - the trajectory of the 2nd person can be obtained from the trajectory of the 1st person
  - by this reflection.
- Thus, at any given moment of time, the midpoint  $M$  is the midpoint between the two persons.
- In these terms:
  - the requirement that they are separated by  $\geq 2m$
  - means that each of them should always be at a distance at least 1 meter from the midpoint  $M$ .
- In other words, both trajectories should avoid the disk of radius 1 meter with a center at the midpoint  $M$ .

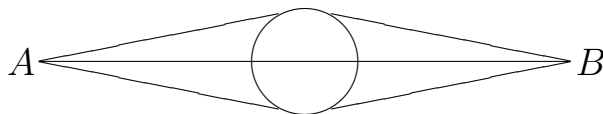


## 226. Formulating the Problem (cont-d)

- We want the fastest possible trajectory.
- The speed is assumed to be constant.
- So, they should follow the shortest possible trajectory.
- In other words, we need to find:
  - the shortest possible trajectory going from point  $A$  to point  $B$
  - that avoids the disk centered at the midpoint  $M$  of the segment  $AB$ .

## 227. Solution

- To get the shortest path, outside the disk, the trajectory should be straight.
- Where it touches the circle, it should be smooth.
- Thus, the solution is as follows:
  - first, we follow a straight line until it touches the circle as a tangent,
  - then, we follow the circle,
  - and finally, we follow the straight line again – which again starts as a tangent to the circle:



## 228. Second Example: Need for Fast Testing

- One of the challenges related to the COVID-19 pandemic is that:
  - this disease has an unusually long incubation period,
  - about 2 weeks.
- As a result, people with no symptoms may be carrying the virus and infecting others.
- As of now, the only way to prevent such infection is to perform massive testing of the population.
- The problem is that there is not enough test kits to test everyone.

## 229. What Was Proposed

- To solve this problem, researchers proposed the following idea:
  - instead of testing everyone individually,
  - why not combine material from a group of several people, and
  - test each combined sample by using a single test kit.
- If no viruses are detected in the combined sample, this means that all the people from the corresponding group are virus-free.
- So there is no need to test them again.
- After this, we need to individually test only folks from the groups that showed the presence of the virus.

## 230. Resulting Problem

- Suppose that we need to test a large population of  $N$  people.
- Based on the previous testing, we know the proportion  $p$  of those who have the virus.
- In accordance with the above idea, we divide  $N$  people into groups.
- The question is: what should be the size  $s$  of each group?
- If the size is too small, we are still using too many test kits.

## 231. Resulting Problem (cont-d)

- If the size is too big, then:
  - every group, with a high probability, has a sick person,
  - so we are not dismissing any people after such testing, and thus, we are not saving any kits at all.
- So what is the optimal size of the group?
- Of course, this is a simplified formulation.
- It does not take into account that:
  - for large group sizes  $s$ , when each individual testing material is diluted too much,
  - tests may not be able to detect infected individuals.

## 232. Let Us Formulate This Problem in Precise Terms

- If we divide  $N$  people into groups of  $s$  persons each, we thus get  $N/s$  groups.
- The probability that a person is virus-free is  $1 - p$ .
- Thus, the probability that all  $s$  people from a group are virus-free is  $(1 - p)^s$ .
- So, out of  $N/s$  groups, the number of virus-free groups is  $(1 - p)^s \cdot (N/s)$ .
- Each of these groups has  $s$  people.
- So the overall number of tested people can be obtained by multiplying the number of virus-free groups by  $s$ .
- This results in  $(1 - p)^s \cdot N$ .

### 233. Formulation and Solution

- For the remaining  $N - (1 - p)^s \cdot N$  folks, we need individual testing.
- So, the overall number of needed test kits is

$$N_t = \frac{N}{s} + N - (1 - p)^s \cdot N.$$

- We want to minimize the number of test kits.
- So, we want to find the group size  $s$  for which this number is the smallest possible.
- Equating the derivative to 0 and dividing both sides of this equation by  $N$ , we get:

$$-\frac{1}{s^2} - (1 - p)^s \cdot \ln(1 - p) = 0.$$



## 234. Solution (cont-d)

- For small  $p$ , we have  $(1-p)^s \approx 1$  and  $\ln(1-p) \approx -p$ , so  $-\frac{1}{s^2} + p \approx 0$ , and  $s \approx \frac{1}{\sqrt{p}}$ .
- For example:
  - for  $p = 1\%$ , we have  $s \approx 10$ ;
  - for  $p = 0.1\%$ , we get  $s \approx 30$ ; and
  - for  $p = 0.01\%$ , we get  $s \approx 100$ .
- The resulting number of tests  $N_t$  can also be approximately estimated.
- When the group size  $s$  is described by the approximate formula, we have  $\frac{N}{s} \approx \sqrt{p} \cdot N$ .
- If we take into account that  $(1-p)^s \approx 1 - p \cdot s$ , then

$$N - (1-p)^s \cdot N \approx p \cdot s \cdot N \approx \sqrt{p} \cdot N.$$

### 235. Solution (cont-d)

- Thus, we get  $N_t \approx \sqrt{p} \cdot N$ .
- For example:
  - for  $p = 1\%$ , we need 10 times fewer test kits than for individual testing;
  - for  $p = 0.1\%$ , we need 30 times fewer test kits; and
  - for  $p = 0.01\%$ , we need 100 times fewer test kits.

## 236. Bibliography

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## Part IX

# Which Distributions (or Families of Distributions) Best Represent

# Interval Uncertainty: Case of Permutation-Invariant Criteria

## 237. Interval Uncertainty Is Ubiquitous

- An engineering designs comes with numerical values of the corresponding quantities, be it:
  - the height of ceiling in civil engineering or
  - the resistance of a certain resistor in electrical engineering.
- Of course, in practice, it is not realistic to maintain the exact values of all these quantities.
- We can only maintain them with some tolerance.
- As a result, the engineers:
  - not only produce the desired (“nominal”) value  $x$  of the corresponding quantity,
  - they also provide positive and negative tolerances  $\varepsilon_+ > 0$  and  $\varepsilon_- > 0$ .

## 238. Interval Uncertainty Is Ubiquitous (cont-d)

- The actual value must be in the interval  $\mathbf{x} = [\underline{x}, \overline{x}]$ , where  $\underline{x} \stackrel{\text{def}}{=} x - \varepsilon_-$  and  $\overline{x} \stackrel{\text{def}}{=} x + \varepsilon_+$ .
- All the manufacturers need to do is to follow these interval recommendations.
- There is no special restriction on probabilities of different values within these intervals.
- These probabilities depends on the manufacturer.
- Even for the same manufacturer, they may change when the manufacturing process changes.

## 239. Data Processing Under Interval Uncertainty Is Often Difficult

- Interval uncertainty is ubiquitous.
- So, many researchers have considered different data processing problems under this uncertainty.
- This research area is known as *interval computations*.
- The problem is that the corresponding computational problems are often very complex.
- They are much more complex than solving similar problems under *probabilistic* uncertainty:
  - when we know the probabilities of different values within the corresponding intervals,
  - we can use Monte-Carlo simulations to gauge the uncertainty of data processing results.

## 240. Interval Data Processing Is Difficult (cont-d)

- A similar problem for interval uncertainty:
  - is NP-hard already for the simplest nonlinear case
  - when the whole data processing means computing the value of a quadratic function.
- It is even NP-hard to find the range of variance when inputs are known with interval uncertainty.
- This complexity is easy to understand.
- Interval uncertainty means that we may have different probability distributions on the given interval.
- So, to get guaranteed estimates, we need, in effect, to consider all possible distributions.
- And this leads to very time-consuming computations.
- For some problems, this time can be sped up, but in general, the problems remain difficult.



## 241. It Is Desirable to Have a Family of Distributions Representing Interval Uncertainty

- Interval computation problems are NP-hard.
- In practical terms, this means that the corresponding computations will take forever.
- So, we cannot consider *all* possible distributions on the interval.
- A natural idea is to consider *some* typical distributions.
- This can be a finite-dimensional family of distributions.
- This can be even a finite set of distributions – or even a single distribution.
- For example, in measurements, practitioners often use uniform distributions on the corresponding interval.
- This selection is even incorporated in some international standards for processing measurement results.

## 242. Family of Distributions (cont-d)

- Of course, we need to be very careful which family we choose.
- By limiting the class of possible distributions, we introduce an artificial “knowledge”.
- Thus, we modify the data processing results.
- So, we should select the family depending on what characteristic we want to estimate.
- We need to beware that:
  - a family that works perfectly well for one characteristic
  - may produce a completely misleading result when applied to some other desired characteristic.
- Examples of such misleading results are well known.

## 243. Continuous Vs. Discrete Distributions

- Usually, in statistics and in measurement theory:
  - when we say that the actual value  $x$  belongs to the interval  $[a, b]$ ,
  - we assume that  $x$  can take any real value between  $a$  and  $b$ .
- However, in practice:
  - even with the best possible measuring instruments,
  - we can only measure the value of the physical quantity  $x$  with some uncertainty  $h$ .
- Thus, from the practical viewpoint, it does not make any sense to distinguish between  $a$  and  $a + h$ .
- Even with the best measuring instruments, we will not be able to detect this difference.

## 244. Continuous Vs. Discrete (cont-d)

- From the practical viewpoint, it makes sense to divide the interval  $[a, b]$  into small subintervals

$$[a, a + h], [a + h, a + 2h], \dots$$

- Within each of them the values of  $x$  are practically indistinguishable.
- It is sufficient to find the probabilities  $p_1, p_2, \dots, p_n$  that the actual value  $x$  is in one of the subintervals:
  - the probability  $p_1$  that  $x$  is in the first small subinterval  $[a, a + h]$ ;
  - the probability  $p_2$  that  $x$  is in the first small subinterval  $[a + h, a + 2h]$ ; etc.
- These probabilities should, of course, add up to 1:

$$\sum_{i=1}^n p_i = 1.$$

## 245. Continuous Vs. Discrete (cont-d)

- In the ideal case, we get more and more accurate measuring instruments – i.e.,  $h \rightarrow 0$ .
- Then, the corresponding discrete probability distributions will tend to continuous ones.
- So, from this viewpoint:
  - selecting a probability distribution means selecting a tuple of values  $p = (p_1, \dots, p_n)$ , and
  - selecting a family of probability distributions means selecting a family of such tuples.

## 246. Example: Estimating Maximum Entropy

- Whenever we have uncertainty, a natural idea is to provide a numerical estimate for this uncertainty.
- It is known that one of the natural measures of uncertainty is Shannon's entropy  $-\sum_{i=1}^n p_i \cdot \log_2(p_i)$ .
- In the case of interval uncertainty, we can have several different tuples.
- In general, for different tuples, entropy is different.
- As a measure of uncertainty of the situation, it is reasonable to take the largest possible value.
- Indeed, Shannon's entropy can be defined as:
  - the average number of binary (“yes”-“no”) questions
  - that are needed to uniquely determine the situation.

## 247. Maximum Entropy (cont-d)

- The larger this number, the larger the initial uncertainty.
- Thus, it is natural to take the largest number of such questions as a characteristic of interval uncertainty.
- For this characteristic, we want to select a distribution:
  - whose entropy is equal to
  - the largest possible entropy of all possible probability distributions on the interval.
- Selecting such a “most uncertain” distribution is known as the *Maximum Entropy approach*.
- This approach has been successfully used in many practical applications.

## 248. Maximum Entropy (cont-d)

- It is well known that:
  - out of all possible tuples with  $\sum_{i=1}^n p_i = 1$ ,
  - the entropy is the largest possible when all the probabilities are equal to each other, i.e., when

$$p_1 = \dots = p_n = 1/n.$$

- In the limit  $h \rightarrow 0$ , such distributions tend to the uniform distribution on the interval  $[a, b]$ .
- This is one of the reasons why uniform distributions are recommended in some measurement standards.



## 249. Modification of This Example

- In addition to Shannon's entropy, there are other measures of uncertainty.
- They are usually called *generalized entropy*.
- For example, in many applications, practitioners use the quantity  $-\sum_{i=1}^n p_i^\alpha$  for some  $\alpha \in (0, 1)$ .
- It is known that when  $\alpha \rightarrow 0$ , this quantity, in some reasonable sense, tends to Shannon's entropy.
- To be more precise:
  - the tuple at which the generalized entropy attains its maximum under different condition
  - tends to the tuple at which Shannon's entropy attains its maximum.
- The maximum of this characteristic is also attained when all the probabilities  $p_i$  are equal to each other.

## 250. Other Examples and Idea

- A recent paper analyzed how to estimate sensitivity of Bayesian networks under interval uncertainty.
- It also turned out that;
  - if we limit ourselves to a single distribution,
  - then the most adequate result also appears if we select a uniform distribution.
- The same uniform distribution appears in many different situations, under different optimality criteria.
- This makes us think that there must be a general reason for this distribution.
- In this talk, we indeed show that there is such a reason.

## 251. Beyond the Uniform Distribution

- For other characteristics, other possible distributions provide a better estimate. For example:
  - if we want to estimate the *smallest* possible value of the entropy,
  - then the corresponding optimal value 0 is attained for several different distributions.
- Specifically, there are  $n$  such distributions corresponding to different values  $i_0 = 1, \dots, n$ .
- In each of these distributions, we have  $p_{i_0} = 1$  and  $p_i = 0$  for all  $i \neq i_0$ .
- In the continuous case  $h \rightarrow 0$ :
  - these probability distributions correspond to point-wise probability distributions
  - in which a certain value  $x_0$  appears with probability 1.

## 252. Beyond the Uniform Distribution (cont-d)

- Similar distributions appear for several other optimality criteria.
- For example, when we minimize generalized entropy.
- How can we explain that these distributions appear as solutions to different optimization problems?
- Similar to the uniform case, there should also be a general explanation.
- A simple general explanation will indeed be provided in this talk.

## 253. Let Us Use Symmetries

- In general, our knowledge is based on *symmetries*, i.e., on the fact that some situations are similar.
- Indeed, if all the world's situations were completely different, we would not be able to make any predictions.
- Luckily, real-life situations have many features in common.
- So we can use the experience of previous situations to predict future ones.
- For example, when a person drops a pen, it starts falling down with the acceleration of  $9.81 \text{ m/sec}^2$ .
- If this person moves to a different location, he or she will get the exact same result.
- This means that the corresponding physics is invariant with respect to shifts in space.

## 254. Let Us Use Symmetries (cont-d)

- Similarly, if the person repeats this experiment in a year, the result will be the same.
- This means that the corresponding physics is invariant with respect to shifts in time.
- Alternatively, if the person turns around a little bit, the result will still be the same.
- This means that the underlying physics is also invariant with respect to rotations, etc.
- This is a very simple example, but such symmetries are invariances are actively used in modern physics.

## 255. Let Us Use Symmetries (cont-d)

- Moreover, many previously proposed fundamental physical theories can be derived from symmetries:
  - Maxwell's equations that describe electrodynamics,
  - Schroedinger's equations that describe quantum phenomena,
  - Einstein's General Relativity equation that describe gravity.
- Symmetries also help to explain many empirical phenomena in computing.
- From this viewpoint:
  - a natural way to look for what the two examples have in common
  - is to look for invariances that they have in common.

## 256. Permutations – Natural Symmetries in the Entropy Example

- We have  $n$  probabilities  $p_1, \dots, p_n$ .
- What can we do with them that would preserve the entropy?
- The easiest possible transformations is when we do not change the values themselves, just swap them.
- Bingo! Under such swap, the value of the entropy does not change.
- Interestingly, the above-described generalized entropy is also permutation-invariant.
- Thus, we are ready to present our general results.



## 257. Definitions and Results

- We say that a function  $f(p_1, \dots, p_n)$  is *permutation-invariant* if for every permutation, we have

$$f(p_1, \dots, p_n) = f(p_{\pi(1)}, \dots, p_{\pi(n)}).$$

- By a *permutation-invariant optimization problem*, we mean a problem of optimizing:
  - a permutation-invariant function  $f(p_1, \dots, p_n)$
  - under constraints of the type  $g_i(p_1, \dots, p_n) = a_i$  or  $h_j(p_1, \dots, p_n) \geq b_j$
  - for permutation-invariant functions  $g_i$  and  $h_j$ .

- **Proposition.** *If a permutation-invariant optimization problem has only one solution, then for this solution:*

$$p_1 = \dots = p_n.$$

- This explains why we get the uniform distribution in several cases (maximum entropy etc.)

## 258. Proof

- We will prove this result by contradiction.
- Suppose that the values  $p_i$  are not all equal.
- This means that there exist  $i$  and  $j$  for which  $p_i \neq p_j$ .
- Let us swap  $p_i$  and  $p_j$ , and denote the corresponding values by  $p'_i$ , i.e.:
  - we have  $p'_i = p_j$ ,
  - we have  $p'_j = p_i$ , and
  - we have  $p'_k = p_k$  for all other  $k$ .
- The values  $p_i$  satisfy all the constraints.
- All the constraints are permutation-invariant.
- So, the new values  $p'_i$  also satisfy all the constraints.
- Since the objective function is permutation-invariant, we have  $f(p_1, \dots, p_n) = f(p'_1, \dots, p'_n)$ .

## 259. Proof (cont-d)

- Since the values  $(p_1, \dots, p_n)$  were optimal, the values  $(p'_1, \dots, p'_n) \neq (p_1, \dots, p_n)$  are thus also optimal.
- This contradicts to the assumption that the original problem has only one solution.
- This contradiction proves for the optimal tuple  $(p_1, \dots, p_n)$  that all the values  $p_i$  are indeed equal to each other.
- The proposition is proven.

## 260. Discussion

- What if the optimal solution is not unique?
- We can have a case when we have a small finite number of solutions.
- We can also have a case when we have a 1-parametric family of solutions – depending on one parameter.
- In our discretized formulation, each parameter has  $n$  values, so this means that we have  $n$  possible solutions.
- Similarly, a 2-parametric family means that we have  $n^2$  possible solutions, etc.
- We say that a problem has a *small finite number of solutions* if it has  $< n$  solutions.
- We say that a problem has a *d-parametric family of solutions* if it has  $\leq n^d$  solutions.

## 261. Second Result

- **Proposition.**

- *If a permutation-invariant optimization problem has a small finite number of solutions,*
  - *then it has only one solution.*
- Due to Proposition 1, in this case, the only solution is the uniform distribution  $p_1 = \dots = p_n$ .

## 262. Proof

- Since  $\sum p_i = 1$ :
  - there is only one possible solution for which

$$p_1 = \dots = p_n :$$

- the solution for which

$$p_1 = \dots = p_n = 1/n.$$

- Thus, if the problem has more than one solution, some values  $p_i$  are different from others.
- In particular, some values are different from  $p_1$ .
- Let  $S$  denote the set of all  $j$  for which  $p_j = p_1$ .
- Let  $m$  denote the number of elements in this set.
- Since some values  $p_i$  are different from  $p_1$ , we have

$$1 \leq m \leq n - 1.$$

### 263. Proof (cont-d)

- Due to permutation-invariance, each permutation of this solution is also a solution.
- For each  $m$ -size subset of  $\{1, \dots, n\}$ , we can have a permutation that transforms  $S$  into this set.
- Thus, it produces a new solution to the original problem.
- There are  $\binom{n}{m}$  such subsets.
- For  $0 < m < n$ , the smallest value  $n$  of  $\binom{n}{m}$  is attained when  $m = 1$  or  $m = n - 1$ .
- Thus, if there is more than one solution, we have at least  $n$  different solutions.
- Since we assumed that we have fewer than  $n$  solutions, this means that we have only one. Q.E.D.

## 264. One More Result

- **Proposition.** *If a permutation-invariant optimization problem has a 1-parametric family of solutions, then:*
  - *this family of solutions is characterized by a real number  $c \leq 1/(n-1)$ , for which*
  - *all these solutions have the following form:  $p_i = c$  for  $i \neq i_0$  and  $p_{i_0} = 1 - (n-1) \cdot c$ .*
- In particular, for  $c = 0$ :
  - we get the above-mentioned 1-parametric family of distributions for which
  - Shannon's entropy (or generalized entropy) attain the smallest possible value.



## 265. Proof

- We have shown that:
  - if in one of the solutions, for some value  $p_i$  we have  $m$  different indices  $j$  with this value,
  - then we will have at least  $\binom{n}{m}$  different solutions.
- For all  $m$  from 2 to  $n - 2$ , this number is at least as large as  $\binom{n}{2} = \frac{n \cdot (n - 1)}{2}$  and is, thus, larger than  $n$ .
- Since overall, we only have  $n$  solutions, this means that it is not possible to have  $2 \leq m \leq n - 2$ .
- So, the only possible values of  $m$  are 1 and  $n - 1$ .

## 266. Proof (cont-d)

- If there was no group with  $n - 1$  values:
  - this would mean that all the groups must have  $m = 1$ ,
  - i.e., consist of only one value.
- In other words, in this case, all  $n$  values  $p_i$  would be different.
- In this case, each of  $n!$  permutations would lead to a different solution.
- So we would have  $n! > n$  solutions, but there are only  $n$  solutions.
- Thus, this case is also impossible.
- So, we do have a group of  $n - 1$  values with the same  $p_i$ .
- Then we get exactly one of the solutions described in the formulation.

## 267. Conclusions

- Traditionally, in engineering, uncertainty is described by a probability distribution.
- In practice, we rarely know the exact distribution.
- In many practical situations:
  - the only information we know about a quantity
  - is the interval of possible values of this quantity.
- And we have no information about the probability of different values within this interval.
- Under such interval uncertainty, we cannot exclude any mathematically possible probability distribution; so:
  - to estimate the range of possible values of the desired uncertainty characteristic,
  - we must, in effect, consider all possible distributions.

## 268. Conclusions (cont-d)

- Not surprisingly, for many characteristics, the corresponding computational problem becomes NP-hard.
- For some characteristics, we can provide a reasonable estimate for their desired range if:
  - instead of all possible distributions,
  - we consider only distributions from some finite-dimensional family.
- For example:
  - to estimate the largest possible value of Shannon's entropy (or of its generalizations),
  - it is sufficient to consider only the uniform distribution.

## 269. Conclusions (cont-d)

- Similarly:
  - to estimate the smallest possible value of Shannon's entropy or of its generalizations,
  - it is sufficient to consider point-wise distributions.
- Different optimality criteria lead to the same distribution – or to the same family of distributions.
- This made us think that there should be a general reason for the appearance of these families.
- In this talk, we show that indeed:
  - the appearance of these distributions and these families can be explained
  - by the fact that all the corresponding optimization problems are permutation-invariant.

## 270. Conclusions (cont-d)

- Thus, in the future, if a reader encounters a permutation-invariant optimization problem:
  - for which it is known that there is a unique solution
  - or that there is only a 1-parametric family of solutions,
  - then there is no need to actually solve the corresponding problem.
- In such situations, it is possible to simply use our general symmetry-based results.
- Thus, we can find a distribution (or a family of distributions) that:
  - for the corresponding characteristic,
  - best represents interval uncertainty.

## Part X

# Expert Knowledge Makes Predictions More Accurate:

# Theoretical Explanation of an Empirical Observation

## 271. Empirical Observation That Needs Explaining

- It is known that the use of expert knowledge makes predictions more accurate.
- For example, computer-based meteorological forecasts are regularly corrected by experts.
- A typical improvement is that the accuracy consistently improves by 10%.
- How can we explain this?



## 272. Towards an Explanation

- Use of expert knowledge means, in effect, that we combine:
  - an estimate produced by a computer model and
  - an expert estimate.
- Let  $\sigma_m$  and  $\sigma_e$  denote the standard deviations, correspondingly, of the model and of the expert estimate.
- In effect, the only information that we have about comparing the two accuracies is that
  - expert estimates are usually less accurate
  - than model results:

$$\sigma_m < \sigma_e.$$

- So, if we fix  $\sigma_e$ , then the only thing we know about  $\sigma_m$  is that  $\sigma_m$  is somewhere between 0 and  $\sigma_e$ .

## 273. Towards an Explanation (cont-d)

- We have no reason to assume that some values from the interval  $[0, \sigma_e]$  are more probable than others.
- Thus, it makes sense to assume that all these values are equally probable.
- So, we have a uniform distribution on this interval.
- For this uniform distribution, the average value of  $\sigma_m$  is equal to  $0.5 \cdot \sigma_e$ .
- Thus, we have  $\sigma_e = 2 \cdot \sigma_m$ .
- In general:
  - if we combine two estimates  $x_m$  and  $x_e$  with accuracies  $\sigma_m$  and  $\sigma_e$ ,
  - then the combined estimate  $x_c$  is obtained by minimizing the sum 
$$\frac{(x_m - x_c)^2}{\sigma_m^2} + \frac{(x_e - x_c)^2}{\sigma_e^2}.$$

## 274. Towards an Explanation (cont-d)

- The resulting estimate is  $x_c = \frac{x_m \cdot \sigma_m^{-2} + x_e \cdot \sigma_e^{-2}}{\sigma_m^{-2} + \sigma_e^{-2}}$ , with accuracy  $\sigma_c^2 = \frac{1}{\sigma_m^{-2} + \sigma_e^{-2}}$ .
- For  $\sigma_e = 2\sigma_m$ , we have  $\sigma_e^{-2} = 0.25 \cdot \sigma_m^{-2}$ , thus  $\sigma_c^2 = \sigma_m^2 \cdot \frac{1}{1 + 0.25} = \sigma_m^2 \cdot \frac{1}{1.25} = 0.8 \cdot \sigma_m^2$ , thus  $\sigma_c \approx 0.9 \cdot \sigma_m$ .
- So we indeed get a 10% increase in the resulting prediction.

## 275. Reference

- N. Silver, *The Signal and the Noise: Why So Many Decisions Fail – but Some Don't*, Penguin Press, New York, 2012.